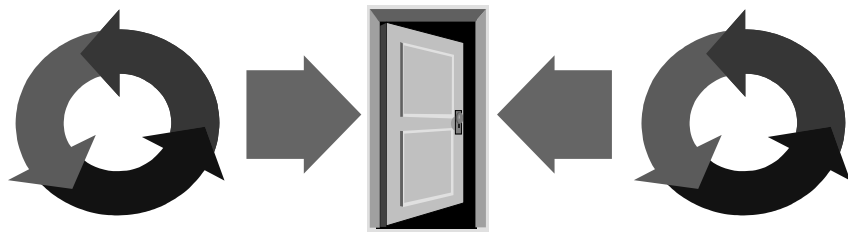


## Monitors to Implement Semaphores



Concepts: monitors:

- encapsulated data + access procedures
- mutual exclusion of access procedure
- > single access procedure active in monitor
- + condition synchronization

Practice: private data and synchronized methods (exclusion).  
wait(), notify() and notifyAll() for condition synch.  
single thread active in the monitor at a time

## Monitors, Basic Principles

**Active** entities (that initiate actions) -> **threads**.

**Passive** entities (that respond to actions) -> **monitors**.

Mutual exclusion of access procedures

For each monitor exists an exclusion lock.

To **enter** the monitor,

a thread acquires the mutual exclusion lock

To **exit** the monitor,

a thread releases the lock, and therefore the monitor,  
for other threads.

## Semaphores

Semaphores are widely used for dealing with inter-process synchronization in operating systems. Semaphore  $s$  is an integer variable that can take only non-negative values.

The only operations permitted on  $s$  are  $up(s)$  and  $down(s)$ . Blocked processes are held in a FIFO queue.

```
down(s): if  $s > 0$ 
            then decrement  $s$ 
            else block execution of calling process
            endif

up(s):    if processes blocked on  $s$ 
            then awaken one of them
            else increment  $s$ 
            endif
```

## Semaphores

*How to model semaphores ?*  
using Petri nets

*How to implement semaphores ?*  
using Java's condition synchronization by  
wait(), notifyAll  
notify()

## Condition Synchronization in Java

Java provides a thread **wait queue** per monitor object with the following methods:

**public final void notify()**

Wakes up a single thread that is waiting on this object's queue.

**public final void notifyAll()**

Wakes up all threads that are waiting on this object's queue.

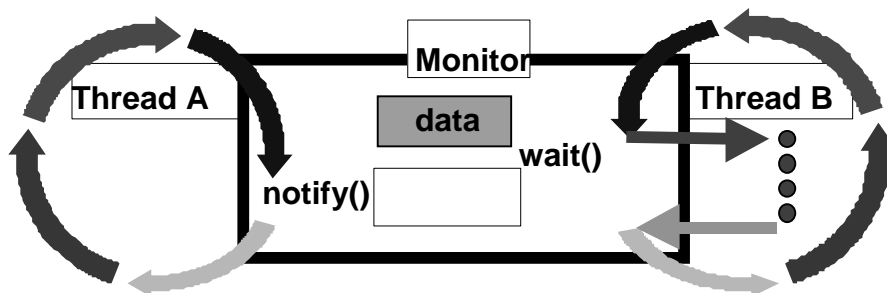
**public final void wait()**

**throws InterruptedException**

Waits to be notified by another thread. The waiting thread releases the synchronization lock associated with the monitor. When notified, the thread must wait to reacquire the monitor before resuming execution.

## condition synchronization in Java

**Wait()** - causes the thread to exit the monitor, permitting other threads to enter the monitor.



## Semaphores in Java

Semaphores are passive objects, therefore implemented as **monitors**.

*(In practice, semaphores are a low-level mechanism often used in implementing the higher-level monitor construct.)*

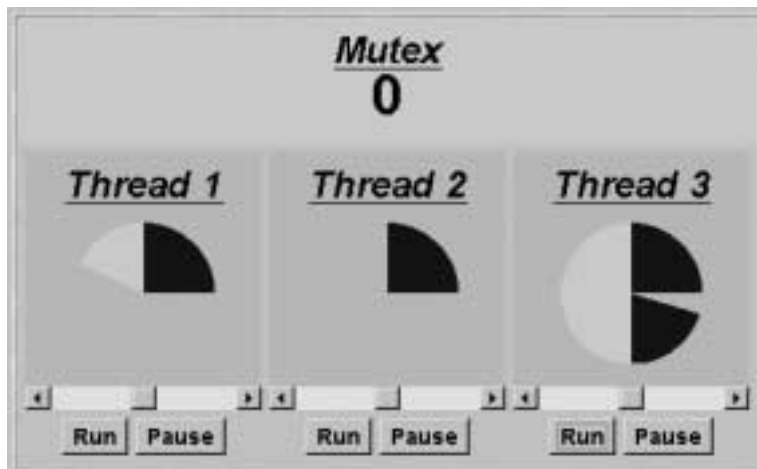
```
public class Semaphore {
    private int value;

    public Semaphore (int initial)
    {value = initial;}

    synchronized public void up() {
        ++value;
        notify();
    }

    synchronized public void down()
    throws InterruptedException {
        while (value == 0) wait();
        --value;
    }
} // Semaphore
```

## SEMADEMO display



current  
semaphore  
value

thread 1 is  
executing  
critical  
actions.

thread 2 is  
blocked  
waiting.

thread 3 is  
executing  
non-critical  
actions.

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Concurrency: monitors & condition synchronization

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## SEMADEMO program - MutexLoop

```
class MutexLoop implements Runnable {  
    Semaphore mutex;  
  
    MutexLoop (Semaphore sema) {mutex=sema;}  
  
    public void run() {  
        try {  
            while(true) {  
                while(!ThreadPanel.rotate());  
                mutex.down();           // get mutual exclusion  
                while(ThreadPanel.rotate()); //critical actions  
                mutex.up();             //release mutual exclusion  
            }  
        } catch (InterruptedException e){}  
    }  
}
```

Threads and  
semaphore are  
created by the applet  
**start()** method.

**ThreadPanel.rotate()** returns **false**  
while executing non-critical actions (dark color) and **true**  
otherwise.

Concurrency: monitors & condition synchronization

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## Monitors, Summary

**Active** entities (that initiate actions) are implemented as **threads**.

**Passive** entities (that respond to actions) are implemented as **monitors**.

Each guarded action in a monitor is implemented as a **synchronized** method which uses a while loop and **wait()** to implement the guard. The while loop condition is the negation of the guard condition.

Changes in the state of the monitor are signaled to waiting threads using **notify()** or **notifyAll()**.

Concurrency: monitors & condition synchronization

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## Summary

### Concepts

monitors: encapsulated data + access procedure

mutual exclusion + condition synchronization

### Practice

private data and synchronized methods in Java

**wait()**, **notify()** and **notifyAll()** for condition  
synchronization

single thread active in the monitor at a time

Concurrency: monitors & condition synchronization

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