

BME & PN, A Case Study - Signalling Cascade

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Definition :

A place/transition Petri net is a quadruple

$$\mathcal{PN} = (P, T, f, m_0), \text{ where}$$

- P, T - finite, non-empty, disjoint sets (**places, transitions**)
- $f : ((P \times T) \cup (T \times P)) \rightarrow \mathbb{N}_0$ (**weighted directed arcs**)
- $m_0 : P \rightarrow \mathbb{N}_0$ (**initial marking**)

Interleaving Semantics : reachability graph / CTL, LTL

Definition :

A biochemically interpreted stochastic Petri net is a quintuple $\mathcal{SPN}_{Bio} = (P, T, f, v, m_0)$, where

- P, T - finite, non-empty, disjoint sets (**places, transitions**)
- $f : ((P \times T) \cup (T \times P)) \rightarrow \mathbb{N}_0$ (**weighted directed arcs**)
- $m_0 : P \rightarrow \mathbb{N}_0$ (**initial marking**)
- $v : T \rightarrow H$ (**stochastic firing rate functions**) with
 - $H := \bigcup_{t \in T} \left\{ h_t \mid h_t : \mathbb{N}_0^{|P|} \rightarrow \mathbb{R}^+ \right\}$
 - $v(t) = h_t$ for all transitions $t \in T$

Semantics : Continuous Time Markov Chain / CSL, PLTLc

Definition :

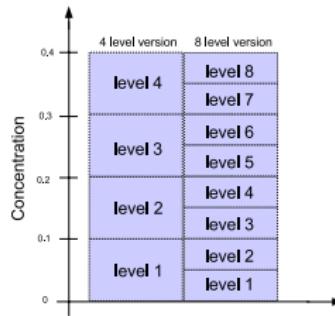
A biochemically interpreted continuous Petri net is a quintuple $\mathcal{CPN}_{Bio} = (P, T, f, v, m_0)$, where

- P, T - finite, non-empty, disjoint sets (**places, transitions**)
- $f : ((P \times T) \cup (T \times P)) \rightarrow \mathbb{R}_0^+$ (**weighted directed arcs**)
- $m_0 : P \rightarrow \mathbb{R}_0^+$ (**initial marking**)
- $v : T \rightarrow H$ (**continuous firing rate functions**) with
 - $H := \bigcup_{t \in T} \{h_t \mid h_t : \mathbb{R}^{|P|} \rightarrow \mathbb{R}^+\}$
 - $v(t) = h_t$ for all transitions $t \in T$

Semantics : ODEs / LTLc

Interpretation of tokens :

- *tokens = molecules, moles*
- *tokens = concentration levels*



Specialised stochastic firing rate function, two examples :

- *molecules semantics*

$$h_t := \textcolor{red}{c_t} \cdot \prod_{p \in \bullet t} \binom{m(p)}{f(p, t)} \quad (1)$$

- *concentration levels semantics*

$$h_t := \textcolor{red}{k_t} \cdot N \cdot \prod_{p \in \bullet t} \left(\frac{m(p)}{N} \right) \quad (2)$$

Various logics each with different expressivity :

- *Branching-time logics* consider all branching time lines
 - Computational Tree Logic (CTL)
 - Continuous Stochastic Logic (CSL)

"There is a possibility that I will stay hungry forever."

"There is a possibility that eventually I am no longer hungry."

- *Linear-time logics* consider separately all single time lines
 - Linear-time Temporal Logics (LTL, LTLc, PLTLc)

"I am hungry."

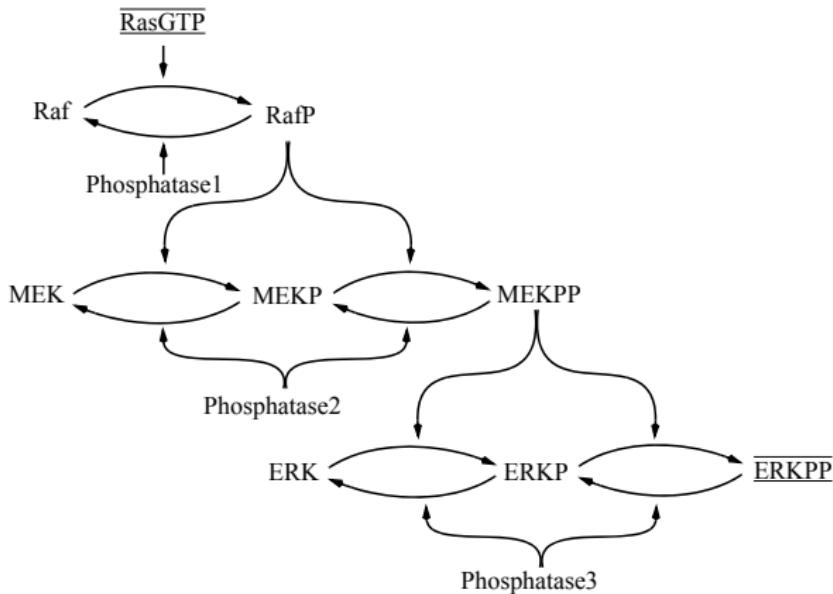
"I am always hungry."

"I will eventually be hungry."

"I will be hungry until I eat something."

Running Case Study

- ... a typical signalling cascade



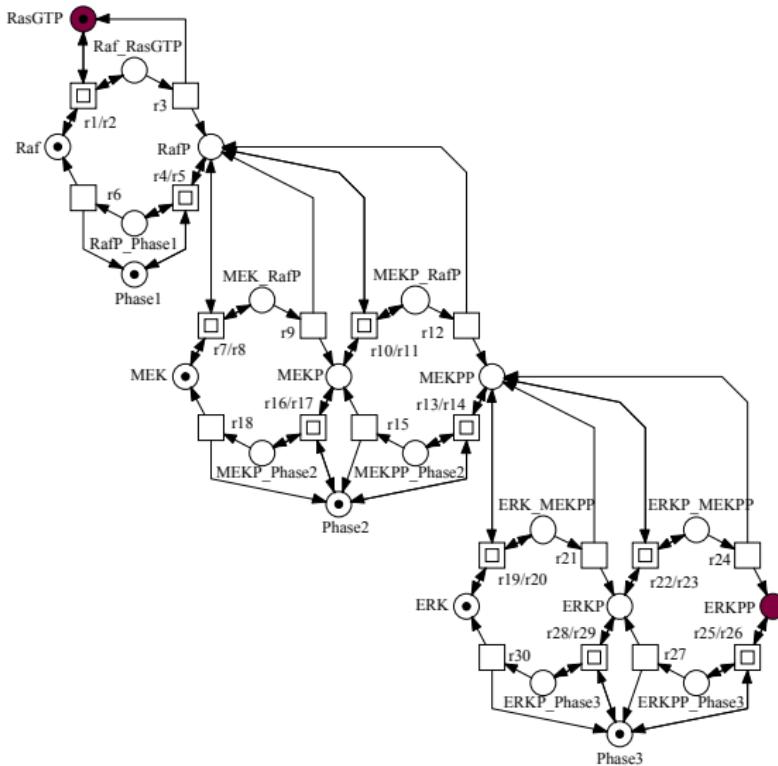
modelled in [Levchenko et al. 2000] like this ...

Running Case Study - Origin

[Levchenko et al. 2000], *Supplemental Material : ODEs*

$$\begin{aligned} \frac{dRaf}{dt} &= k_2 * Raf_RasGTP + k_6 * RafP_Phase1 - k_1 * Raf * RasGTP \\ \frac{dRasGTP}{dt} &= k_2 * Raf_RasGTP + k_3 * Raf_RasGTP - k_1 * Raf * RasGTP \\ \frac{dRaf_RasGTP}{dt} &= k_1 * Raf * RasGTP - k_2 * Raf_RasGTP - k_3 * Raf_RasGTP \\ \frac{dRafP}{dt} &= k_3 * Raf_RasGTP + k_{12} * MEKP_RafP + k_9 * MEK_RafP + \\ &\quad k_5 * RafP_Phase1 + k_8 * MEK_RafP + k_{11} * MEKP_RafP - \\ &\quad k_7 * RafP * MEK - k_{10} * MEKP * RafP - k_4 * Phase1 * RafP \\ \frac{dRafP_Phase1}{dt} &= k_4 * Phase1 * RafP - k_5 * RafP_Phase1 - k_6 * RafP_Phase1 \\ \frac{dMEK_RafP}{dt} &= k_7 * RafP * MEK - k_8 * MEK_RafP - k_9 * MEK_RafP \\ \frac{dMEKP_RafP}{dt} &= k_{10} * MEKP * RafP - k_{11} * MEKP_RafP - k_{12} * MEKP_RafP \\ \frac{dMEKP_Phase2}{dt} &= k_{16} * Phase2 * MEKP - k_{18} * MEKP_Phase2 - k_{17} * MEKP_Phase2 \\ \frac{dMEKPP_Phase2}{dt} &= k_{13} * MEKPP * Phase2 - k_{15} * MEKPP_Phase2 - k_{14} * MEKPP_Phase2 \\ \frac{dERK}{dt} &= k_{20} * ERK_MEKPP + k_{30} * ERKP_Phase3 - k_{19} * MEKPP * ERK \\ \frac{dERK_MEKPP}{dt} &= k_{19} * MEKPP * ERK - k_{20} * ERK_MEKPP - k_{21} * ERK_MEKPP \\ \frac{dERKP_MEKPP}{dt} &= k_{22} * MEKPP * ERKP - k_{24} * ERKP_MEKPP - k_{23} * ERKP_MEKPP \\ \text{etcetera} &= \dots \end{aligned}$$

Running Case Study



- initial marking construction

P-invariants

- subnetwork identification

- P-invariants : token preserving modules (*mass conservation*)
- T-invariants : state repeating modules (*elementary modes*)

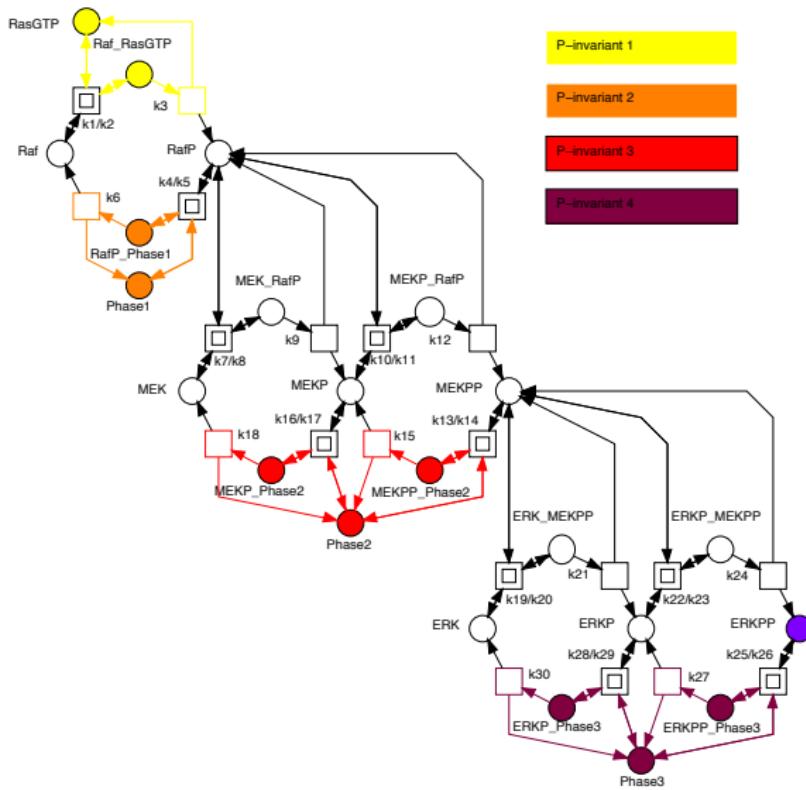
- general behavioural properties

- *boundedness* : every place gets finite token number only
- *liveness* : every transition may happen forever
- *reversibility* : every state may be reached forever

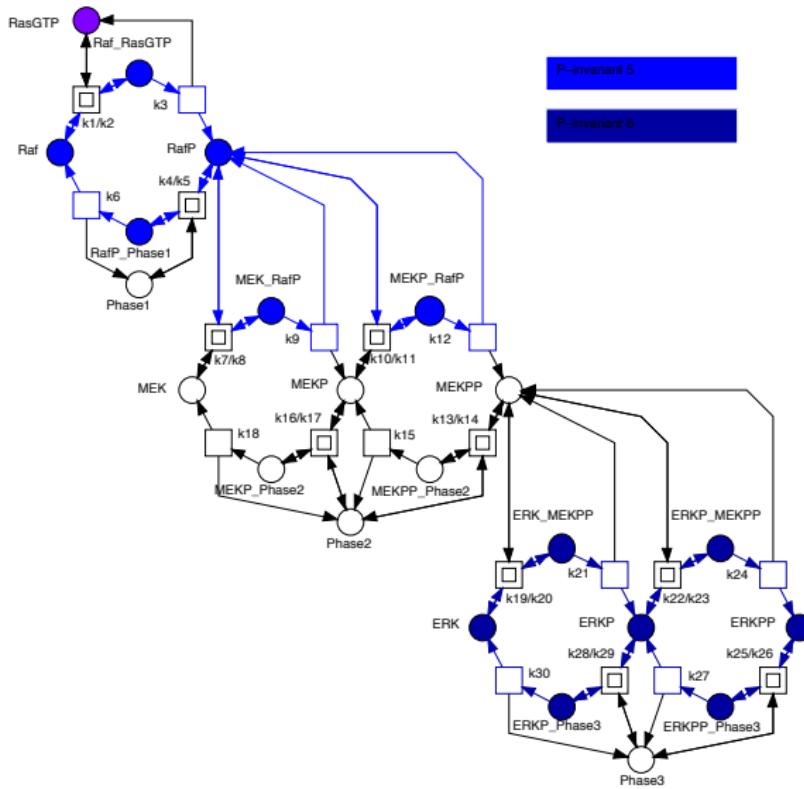
- special behavioural properties

CTL / LTL model checking

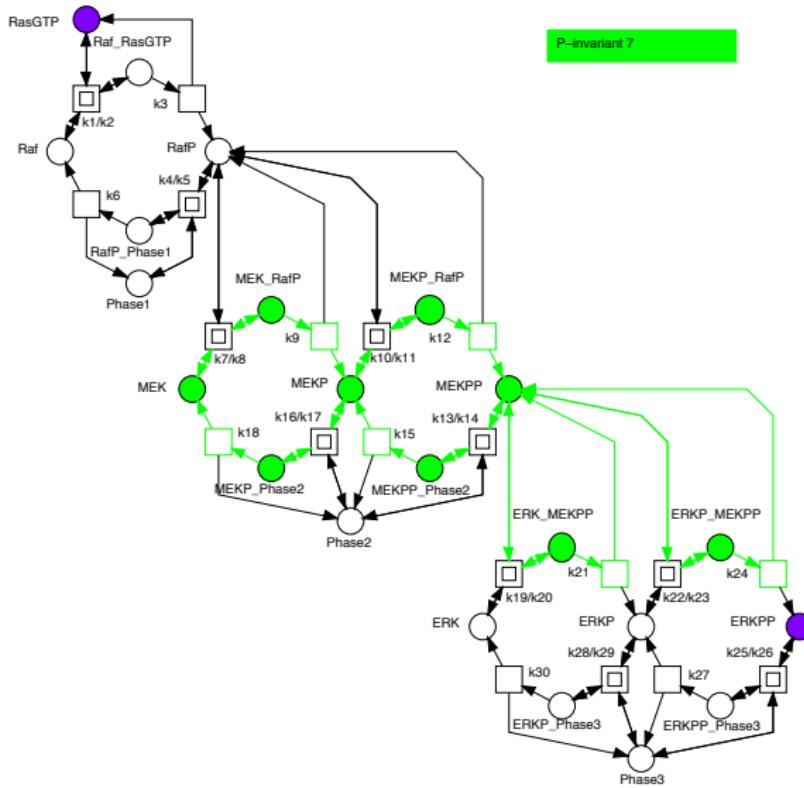
Running Case Study - P-invariants



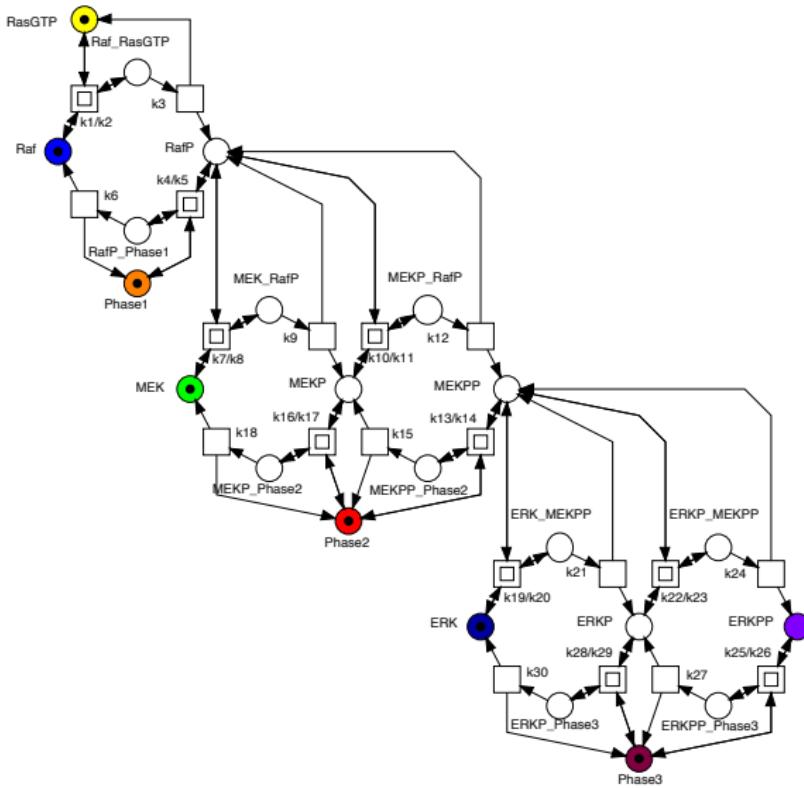
Running Case Study - P-invariants



Running Case Study - P-invariants



Running Case Study - initial marking



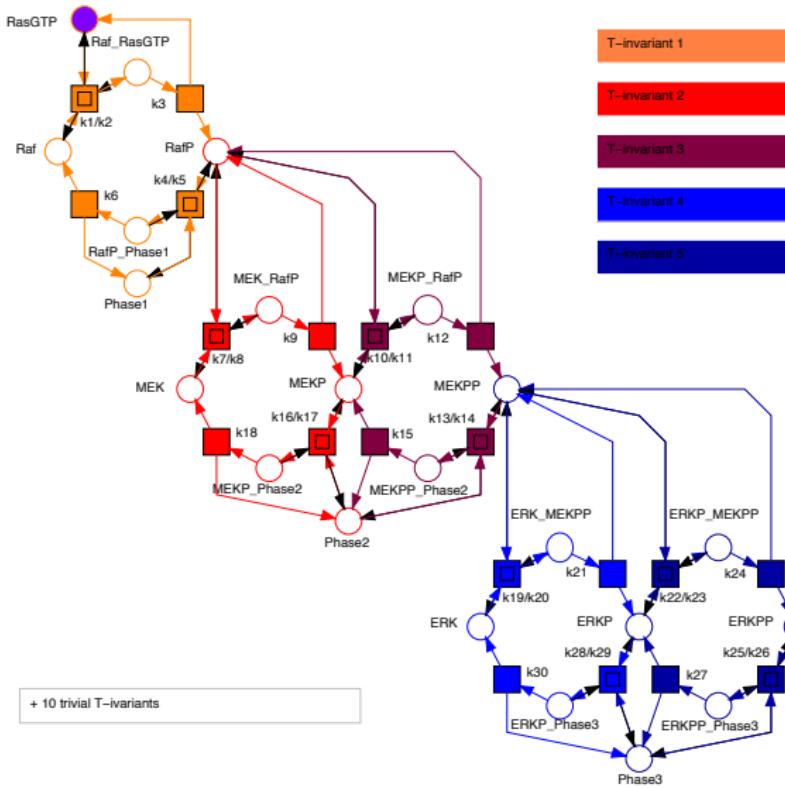
Running Case Study - general properties

- *state space*

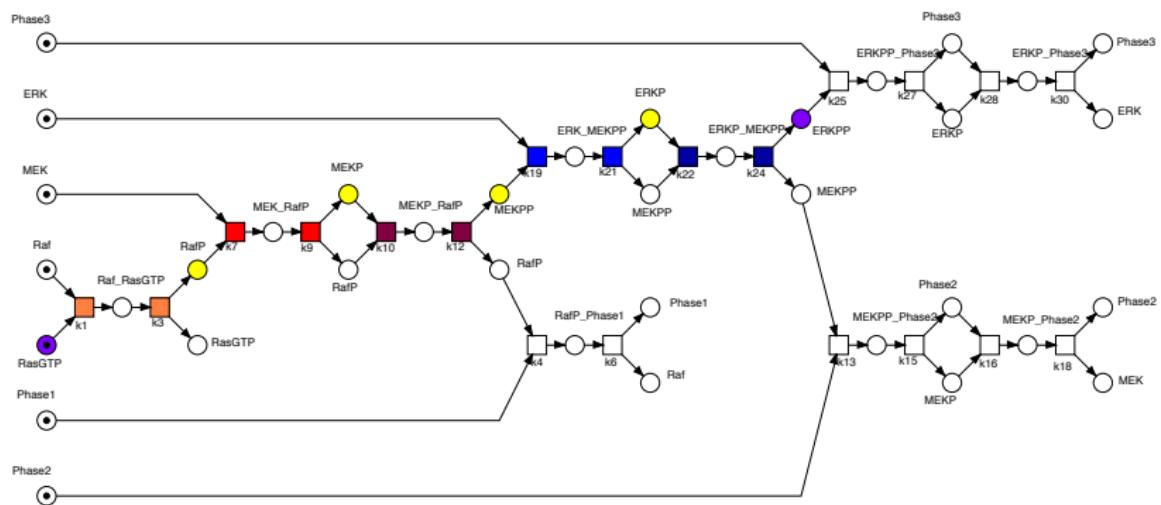
levels	reachability graph number of states	IDD data structure number of nodes
1	118	52
4	$2.4 \cdot 10^4$	115
8	$6.1 \cdot 10^6$	269
80	$5.6 \cdot 10^{18}$	13,472
120	$1.7 \cdot 10^{21}$	29,347

- Covered by P-invariants (CPI) \Rightarrow **bounded**
- Siphon-Trap Property (STP) holds \Rightarrow **no dead states**
- reachability graph
 - strongly connected \Rightarrow **reversible**
 - contains every transition (reaction) \Rightarrow **live**

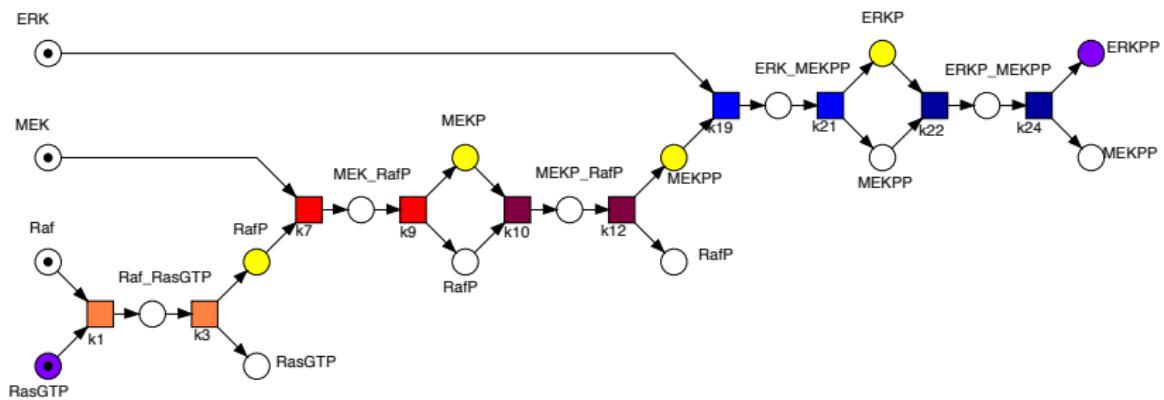
Running Case Study - T-invariants



Running Case Study - partial order run of I/O T-invariant



Running Case Study - partial order run of I/O T-invariant

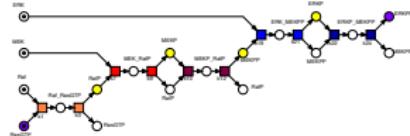


There is a path ...

- $EX\phi$
 - if there is a state reachable by one step where ϕ holds.
- $EF\phi$
 - if there is a path where ϕ holds finally, i.e., at some point.
- $EG\phi$
 - if there is a path where ϕ holds globally, i.e., forever.
- $E(\phi_1 U \phi_2)$
 - if there is a path where ϕ_1 holds until ϕ_2 holds.

For all path ...

- $AX\phi$
 - if ϕ holds for all states which are reachable by one step.
- $AF\phi$
 - if ϕ holds finally (at some point) for all paths.
- $AG\phi$
 - if ϕ holds globally (i.e. for ever) for all paths.
- $A(\phi_1 U \phi_2)$
 - if ϕ_1 holds until ϕ_2 holds for all paths.



property Q1 :

The signal sequence predicted by the partial order run of the I/O T-invariant is the only possible one;
i.e., starting at the initial state, it is necessary to pass through RafP, MEKP, MEKPP and ERKP in order to reach ERKPP.

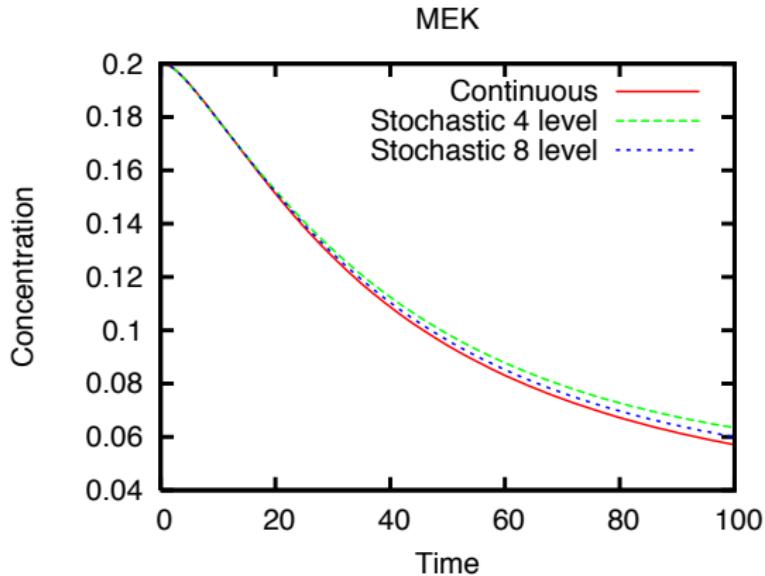
$$\neg [E(\neg \text{RafP} \cup \text{MEKP}) \vee \\ E(\neg \text{MEKP} \cup \text{MEKPP}) \vee \\ E(\neg \text{MEKPP} \cup \text{ERKP}) \vee \\ E(\neg \text{ERKP} \cup \text{ERKPP})]$$

- *isomorphy of reachability graph and CTMC,*
thus all qualitative properties still valid
- *How many levels needed for quantitative evaluation ?*
 - state space(1 levels) = 118 (Boolean interpretation)
 - state space(4 levels) = 24,065
 - state space(8 levels) = 6,110,643
- *equivalence check*

$$C_{RafP}(t) = \frac{0.1}{s} \cdot \underbrace{\sum_{i=1}^{4s} (i \cdot P(L_{RafP}(t) = i))}_{\text{expected value of } L_{RafP}(t)}$$

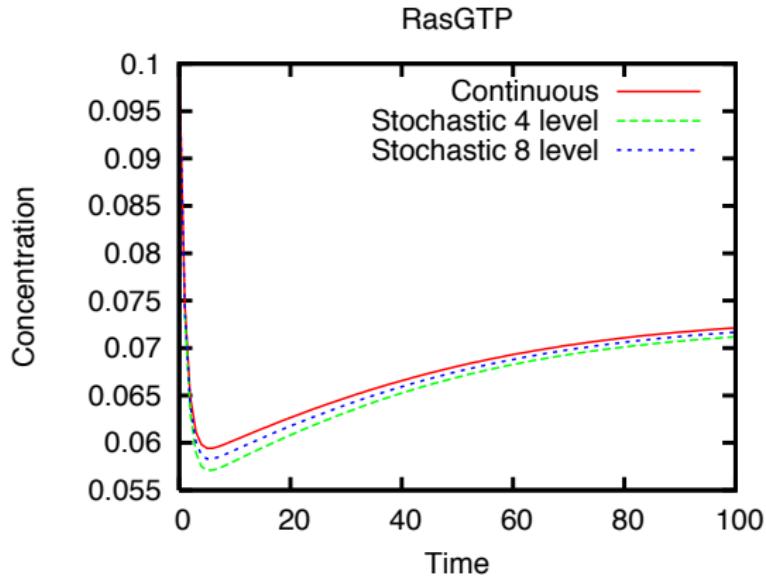
Stochastic Model Checking - Preparation

- equivalence check, results, e.g. for MEK :



Stochastic Model Checking - Preparation

- equivalence check, results, e.g. for RasGTP :



Replaces the path quantifiers (E, A) in CTL by the probability operator $P_{\leq x}$, where $\leq x$ specifies the probability x of the formula.

- $P_{=?}[X\phi]$
 - prob there is a state reachable by one step where ϕ holds.
- $P_{=?}[F\phi]$
 - prob there is a path where ϕ holds finally, i.e., at some point.
- $P_{=?}[G\phi]$
 - prob there is a path where ϕ holds globally, i.e., forever.
- $P_{=?}[\phi_1 U \phi_2]$
 - prob there is a path where ϕ_1 holds until ϕ_2 holds.

Syntactic sugar

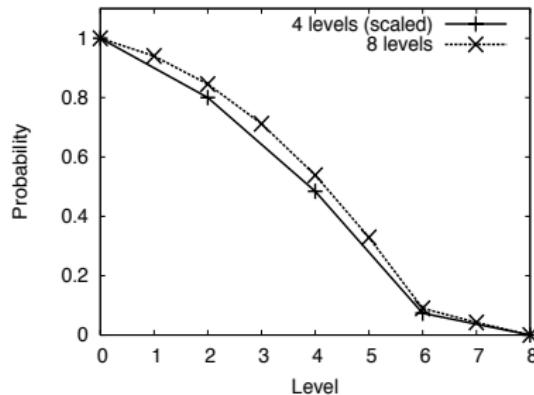
- $\phi_1\{\phi_2\}$ - ϕ_1 happens from the first time ϕ_2 happens, where no temporal operators in ϕ_2 .

Stochastic Model Checking (CSL)

property S1 :

What is the probability of the concentration of RafP increasing, when starting in a state where the level is already at L ?

$$P_{=?} [(\text{RafP} = L) \mathbf{U}^{<=100} (\text{RafP} > L) \{ \text{RafP} = L \}]$$

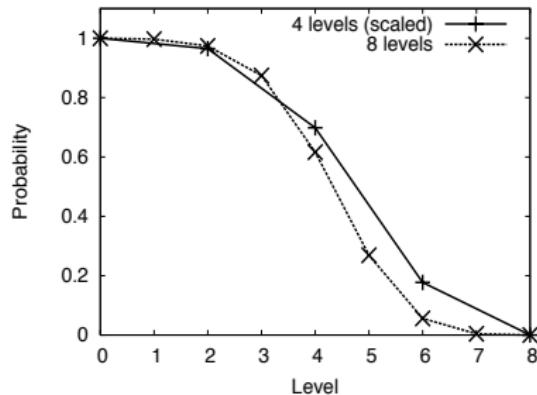


Stochastic Model Checking (CSL)

property S2 :

What is the probability that RafP is the first species to react ?

$$\mathbf{P}_{=?} [((\text{MEKPP} = 0) \wedge (\text{ERKPP} = 0)) \mathbf{U}^{<=100} (\text{RafP} > \text{L}) \\ \{ (\text{MEKPP} = 0) \wedge (\text{ERKPP} = 0) \wedge (\text{RafP} = 0) \}]$$



Example figures for MC2 model checking of property S1 at varying number of levels/molecules.

Levels	MC Time ^a	Simulation Output Size
4	10 s ^b	750 KB
8	15 s ^b	1.5 MB
40	1.5 minutes ^b	7.5 MB
400	1 minute ^c	80 MB
4,000	30 minutes ^c	900 MB

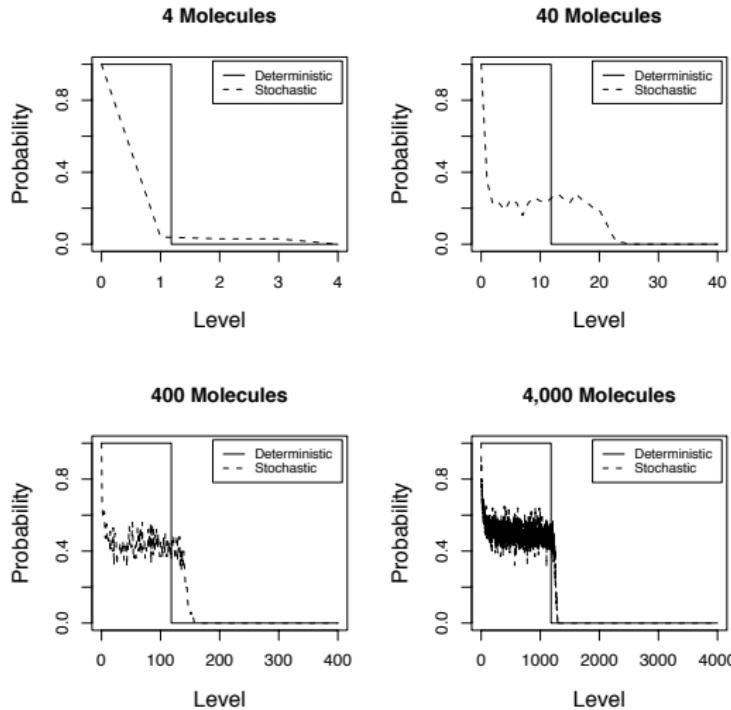
^a Both Gillespie simulation and MC2 checking.

^b Computation on a standard workstation.

^c Distributed computation on a computer cluster comprising 45 Sun X2200 servers each with 2 dual core processors (180 CPU cores).

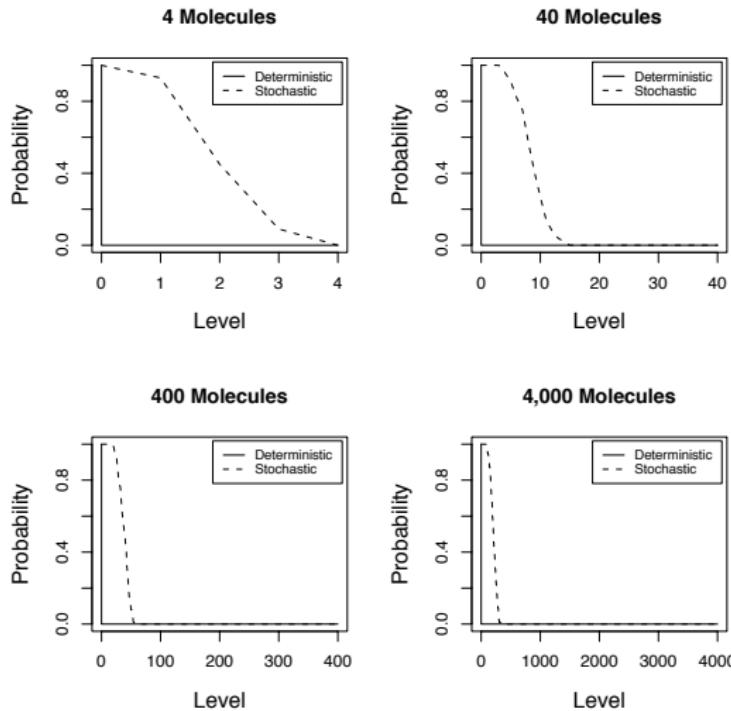
Stochastic Model Checking - Simulative Approach

- *S1 at varying number of molecules shows progression towards deterministic behaviour as number of molecules increases.*



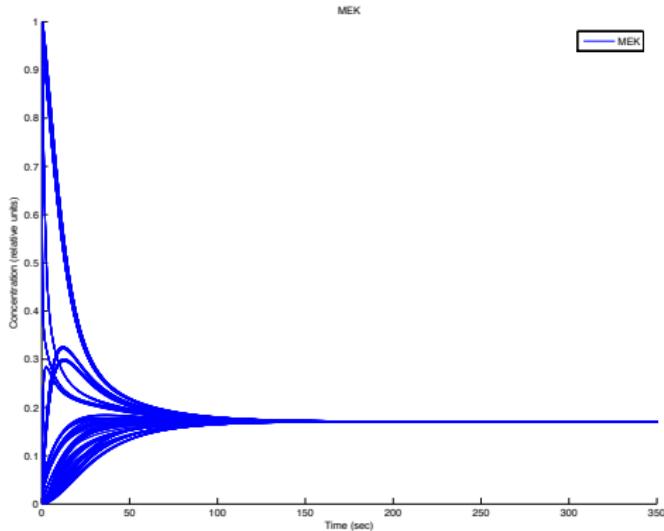
Stochastic Model Checking - Simulative Approach

- *S2 at varying number of molecules shows progression towards deterministic behaviour as number of molecules increases.*



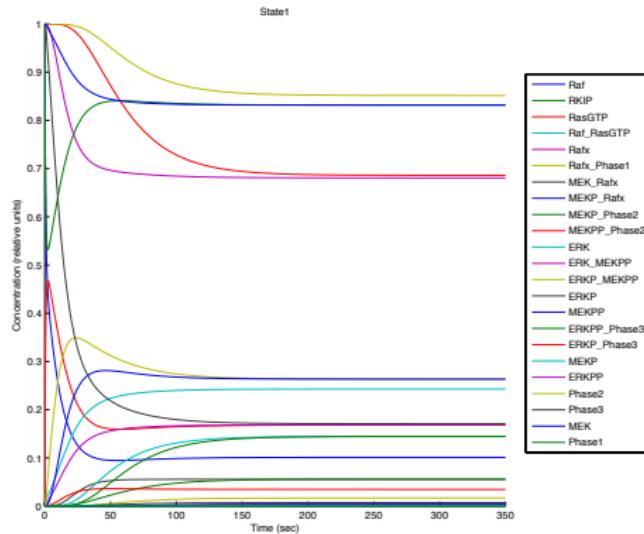
Continuous Model Checking - Preparation

- steady state analysis, results for all 118 'good' states, e.g. for MEK :



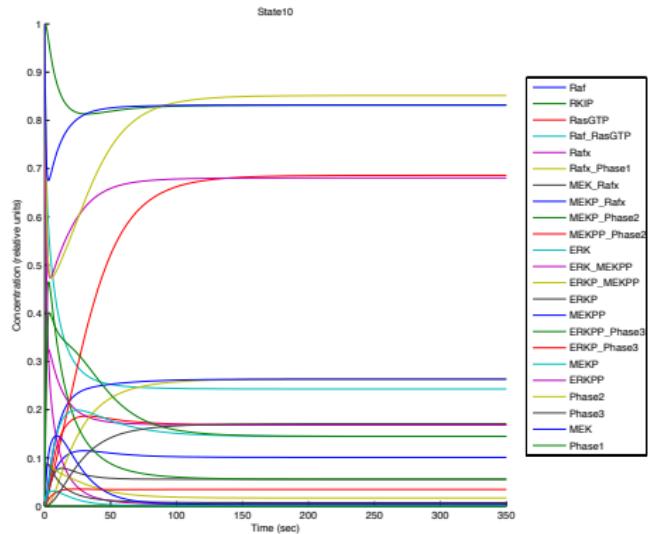
Continuous Model Checking - Preparation

- steady state analysis for state 1 :



Continuous Model Checking - Preparation

- steady state analysis for state 10 :



For all single path ...

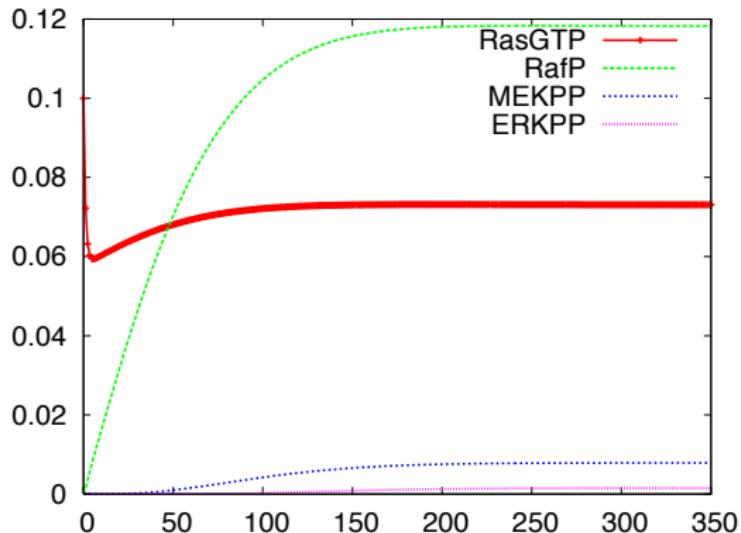
- $X\phi$
 - ϕ happens in the next time point.
- $F\phi$
 - ϕ happens at some time.
- $G\phi$
 - ϕ always happens.
- $A(\phi_1 U \phi_2)$
 - ϕ_1 happens until ϕ_2 happens.

Syntactic sugar

- $\phi_1\{\phi_2\}$ - ϕ_1 happens from the first time ϕ_2 happens, where no temporal operators in ϕ_2 .

Continuous Model Checking (LTLc)

- *transient analysis for RasGTP, RafP, MEKPP, ERKPP :*



property C1 :

The concentration of RafP rises to a significant level, while the concentrations of MEKPP and ERKPP remain close to zero ;
i.e. *RafP is really the first species to react.*

$$((\text{MEKPP} < 0.001) \wedge (\text{ERKPP} < 0.0002)) \text{ U } (\text{RafP} > 0.06)$$

property C2 :

if the concentration of RafP is at a significant concentration level and that of ERKPP is close to zero, then both species remain in these states until the concentration of MEKPP becomes significant ; i.e. *MEKPP is the second species to react.*

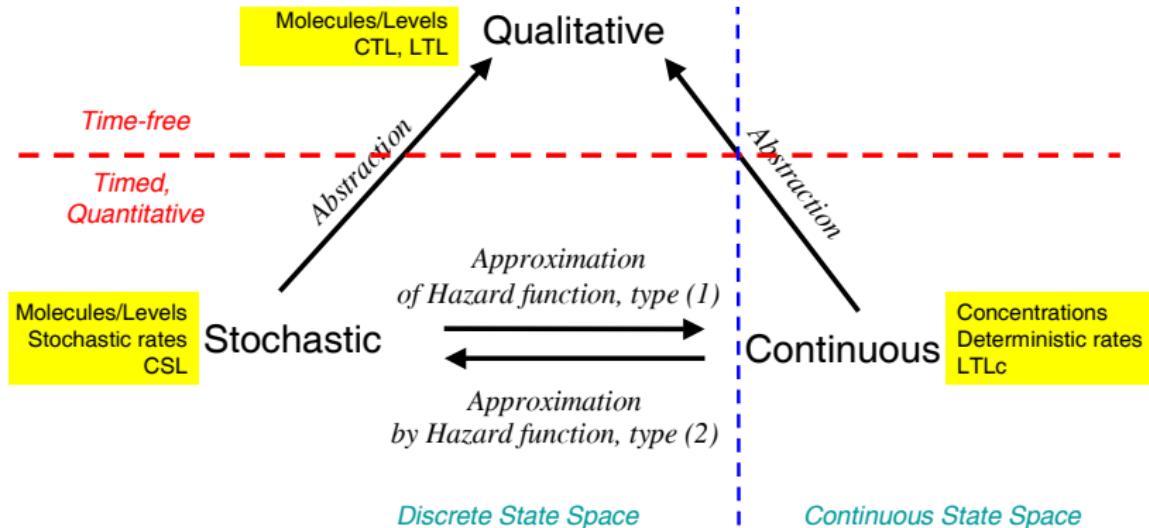
$$\begin{aligned} & ((\text{RafP} > 0.06) \wedge (\text{ERKPP} < 0.0002)) \Rightarrow \\ & ((\text{RafP} > 0.06) \wedge (\text{ERKPP} < 0.0002)) \mathbf{U} (\text{MEKPP} > 0.004) \end{aligned}$$

property C3 :

if the concentrations of RafP and MEKPP are significant, they remain so, until the concentration of ERKPP becomes significant ;
i.e. *ERKPP is the third species to react.*

$$\begin{aligned} & ((\text{RafP} > 0.06) \wedge (\text{MEKPP} > 0.004)) \Rightarrow \\ & \quad ((\text{RafP} > 0.06) \wedge (\text{MEKPP} > 0.004)) \mathbf{U} (\text{ERKPP} > 0.0005) \end{aligned}$$

Framework



- *validation criterion 1*
 - all expected structural properties hold
 - all expected general behavioural properties hold
- *validation criterion 2*
 - CPI
 - no minimal P-invariant without biological interpretation (?)
- *validation criterion 3*
 - CTI
 - no minimal T-invariant without biological interpretation
 - no known biological behaviour without corresponding T-invariant
- *validation criterion 4*
 - all expected special behavioural properties hold
 - temporal-logic properties yield TRUE

- *model construction, animation, simulation*
 - Snoopy
- *qualitative analysis*
 - Charlie
 - Marcie
- *stochastic analysis*
 - analytical model checking : Marcie/CSL (*previously PRISM*)
 - simulative model checking : Marcie/PLTLc, MC2/PLTLc
- *continuous analysis*
 - model checking : MC2(LTLc) (*previously BioCham*)

Finally

This is an updated version of

M Heiner, D Gilbert, R Donaldson :
Petri Nets for Systems and Synthetic Biology ;
SFM 2008, Bertinoro, Springer LNCS 5016, pp. 215-264, 2008.

all data files and analysis results available at
www-dssz.informatik.tu-cottbus.de/examples/levchenko