

The Petri Net Model of the Sucrose-to-Starch Breakdown in the potato tuber



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Outline

- Introduction
- Sucrose-to-starch breakdown in the potato tuber
- The Petri net model
- Qualitative analysis
- Simulation of the net
- Conclusions

Introduction

Cooperation: Björn Junker, Max Planck Institute for Molecular Plant Physiology,
Golm

Situation before starting kinetic modelling:

incomplete kinetic data

literature search

try-and-error-technique to find the steady state

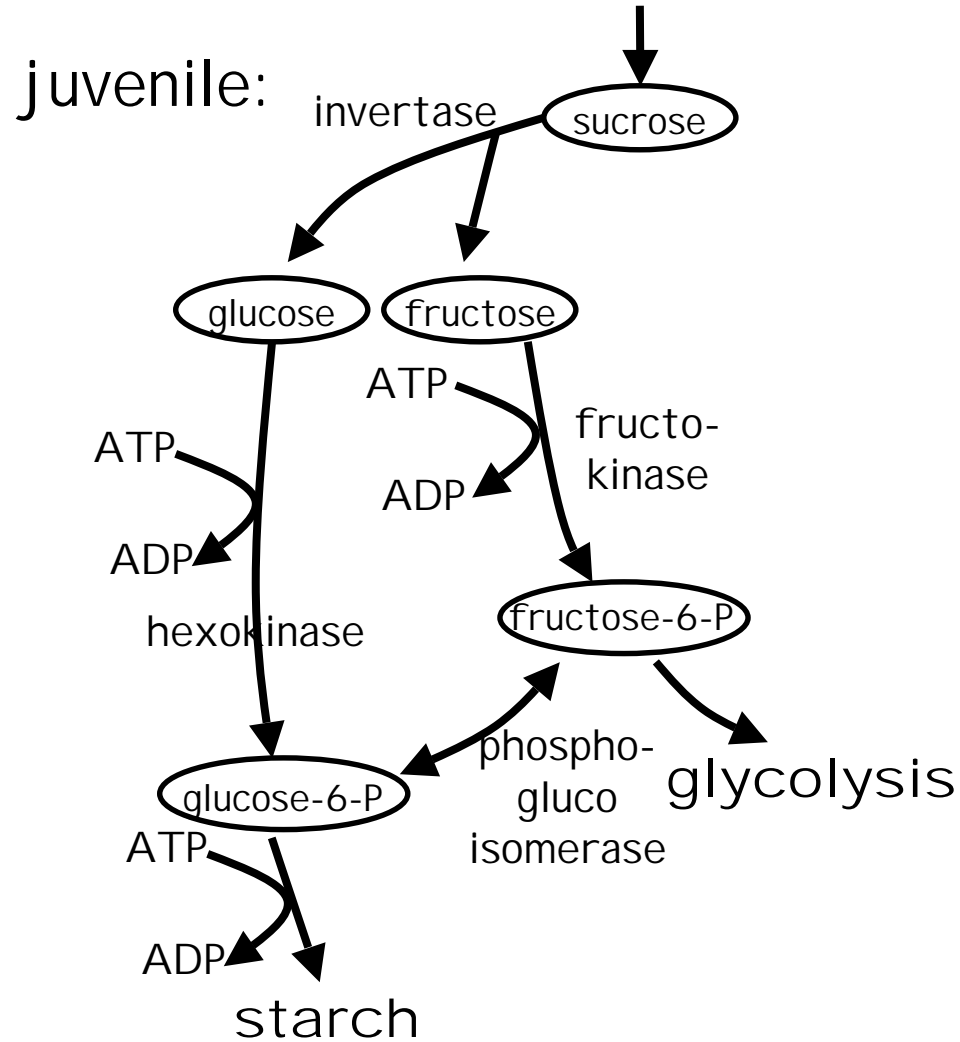
using GEPASI Mendes, *Comp.Appl.Biosci.* (1993)

Qualitative modelling as the first step

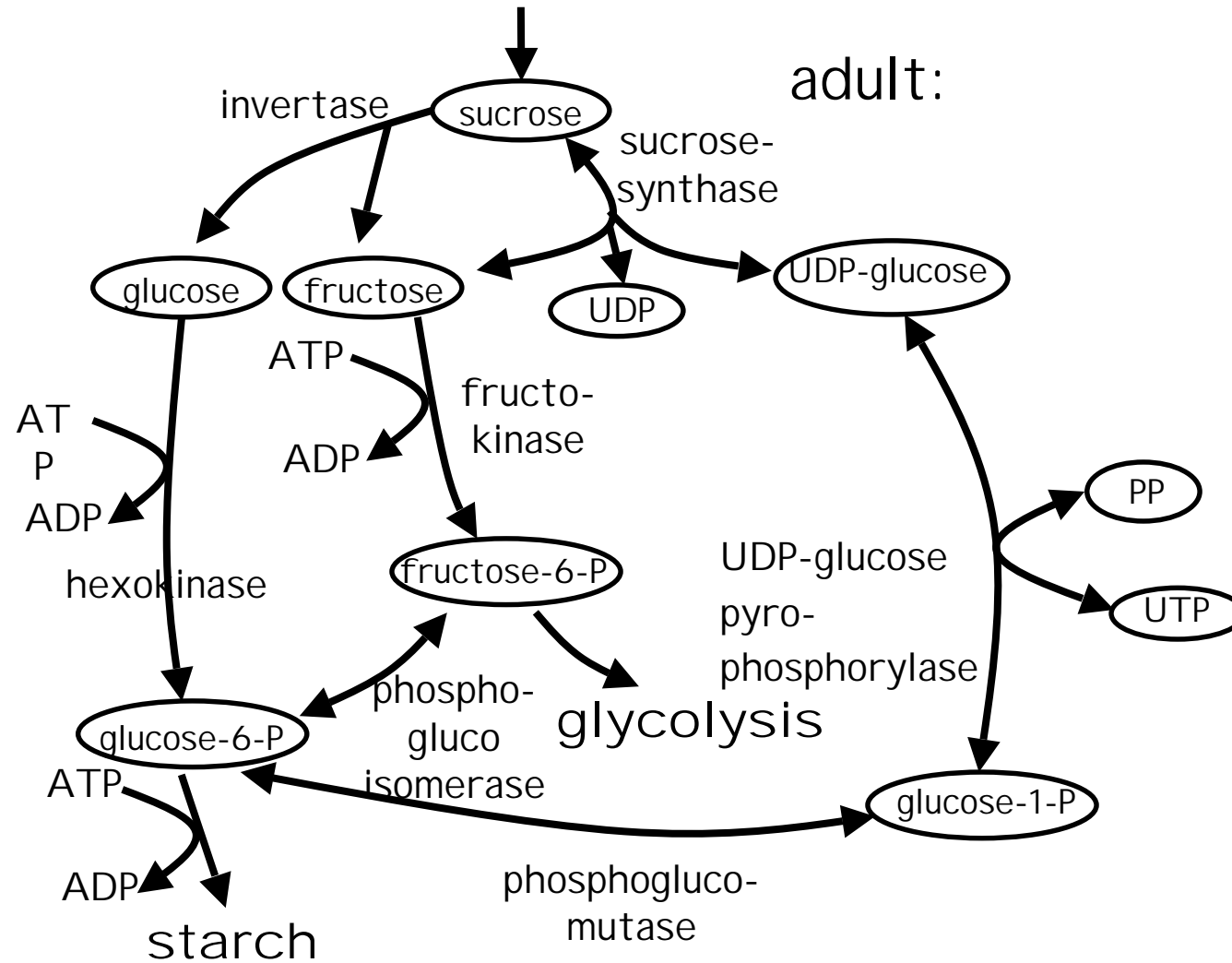
Basic dynamic properties: liveness, reversibility, boundedness,
dead states, deadlocks, traps

Basic structure properties: invariants, robustness, alternative pathways,

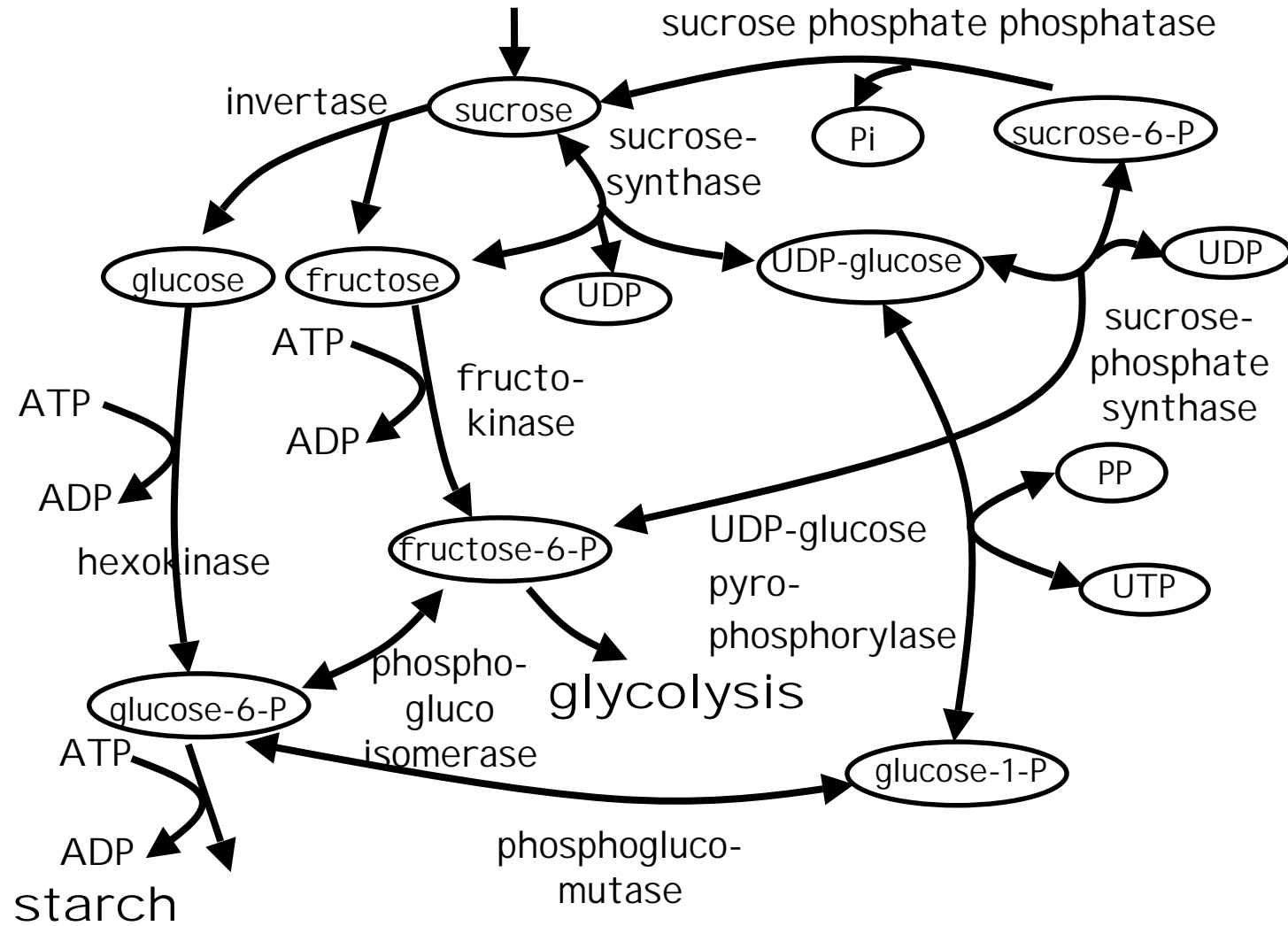
ucrose-to-starch-pathway in potato tuber



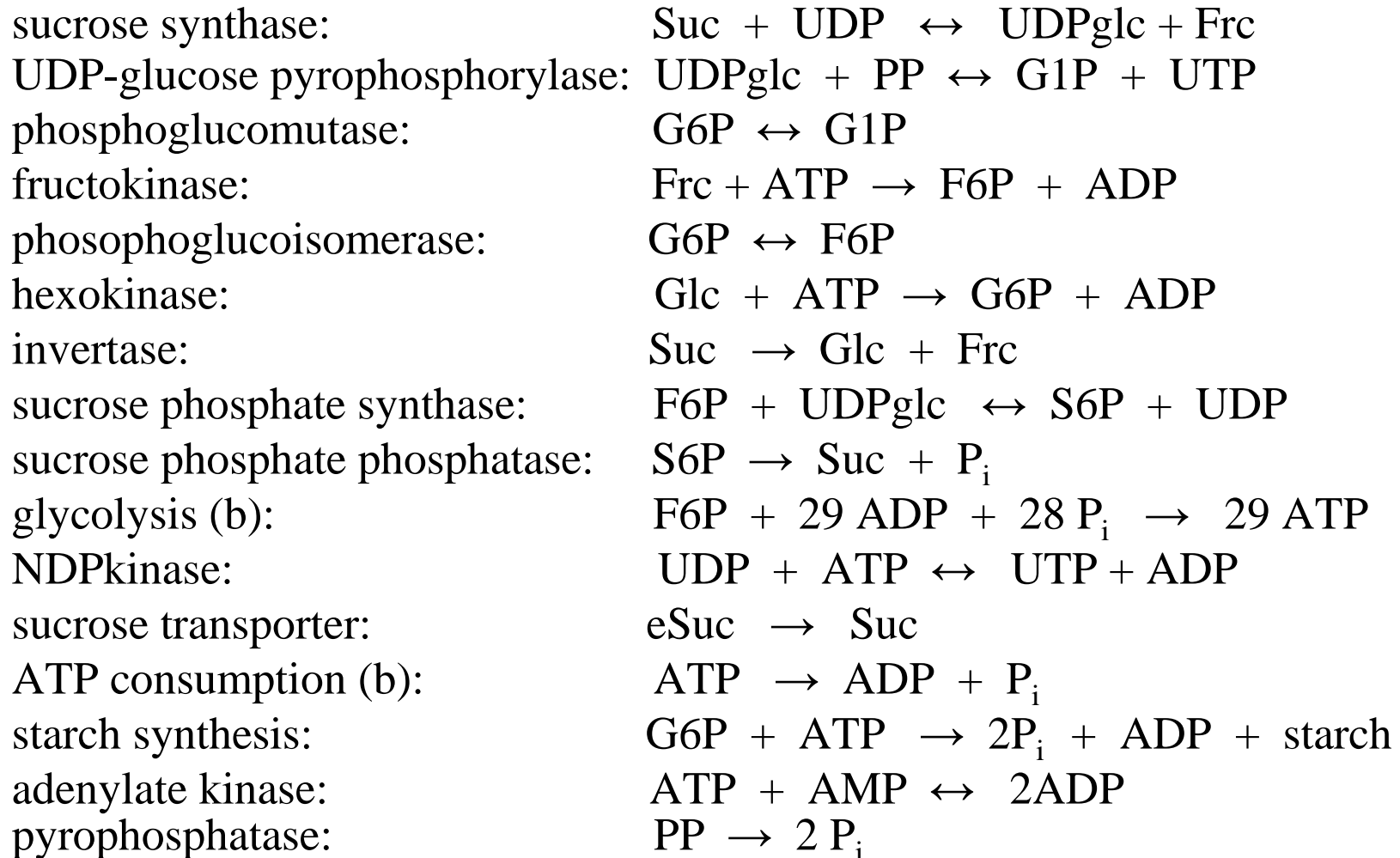
ucrose-to-starch-pathway in potato tuber



ucrose-to-starch-pathway in potato tuber



ucrose-to-starch-pathway in potato tuber



etri net basics

Nodes :
(vertices)

places

transitions

passive elements
conditions
states
chemical compounds
metabolites

active elements
events
actions
chemical reactions
conversions of metabolites
catalysed by enzymes



event



etri net basics

Nodes :
(vertices)

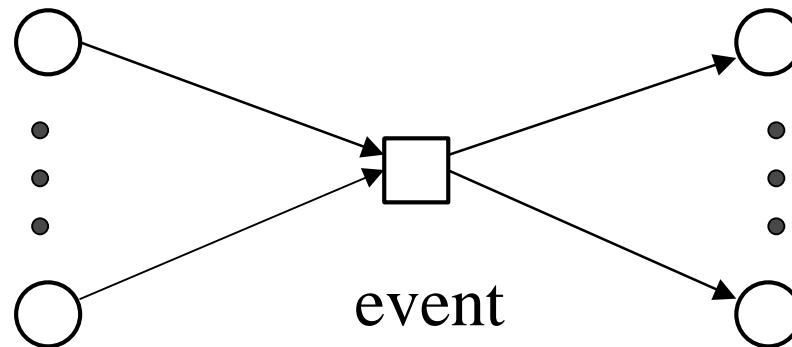
places

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catalysed by enzymes

Arcs:
(edges)



pre-conditions
pre-places

post-conditions
post-places

etri net basics

Nodes :
(vertices)

places

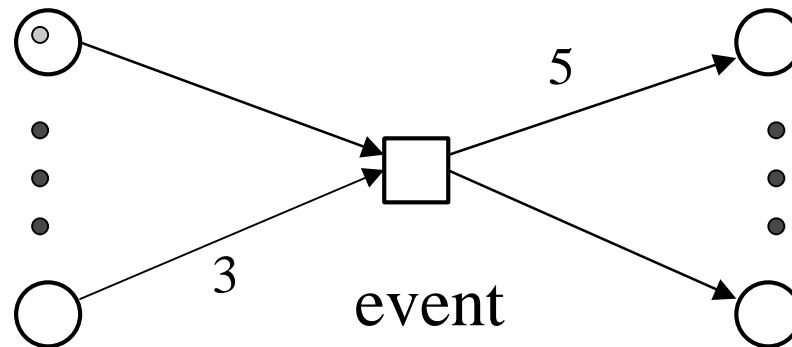
transitions

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Arcs:
(edges)

Tokens



pre-conditions
pre-places

post-conditions
post-places

etri net basics

Tokens: movable objects in discrete units,
e.g. units of substances (mole)

○ condition is not fulfilled

◦ condition is (one time) fulfilled

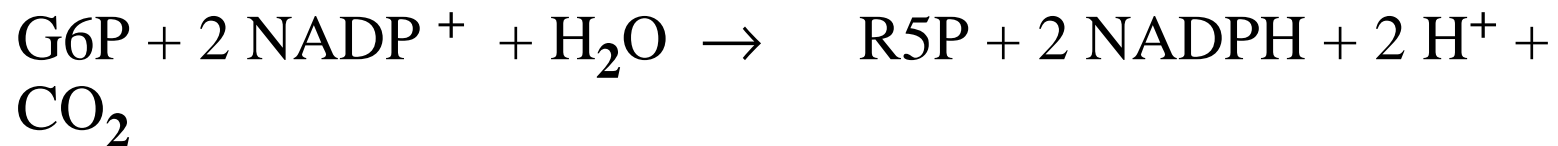
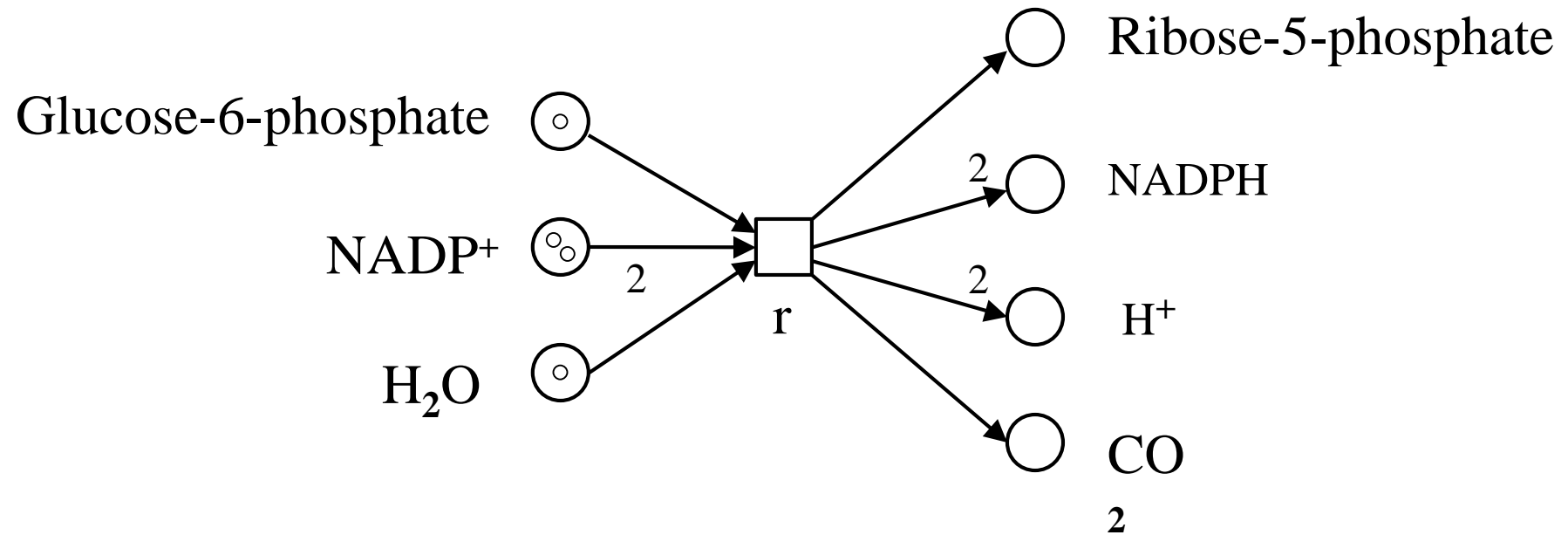
Ⓝ condition is n times fulfilled

Marking: system state, token distribution, initial marking

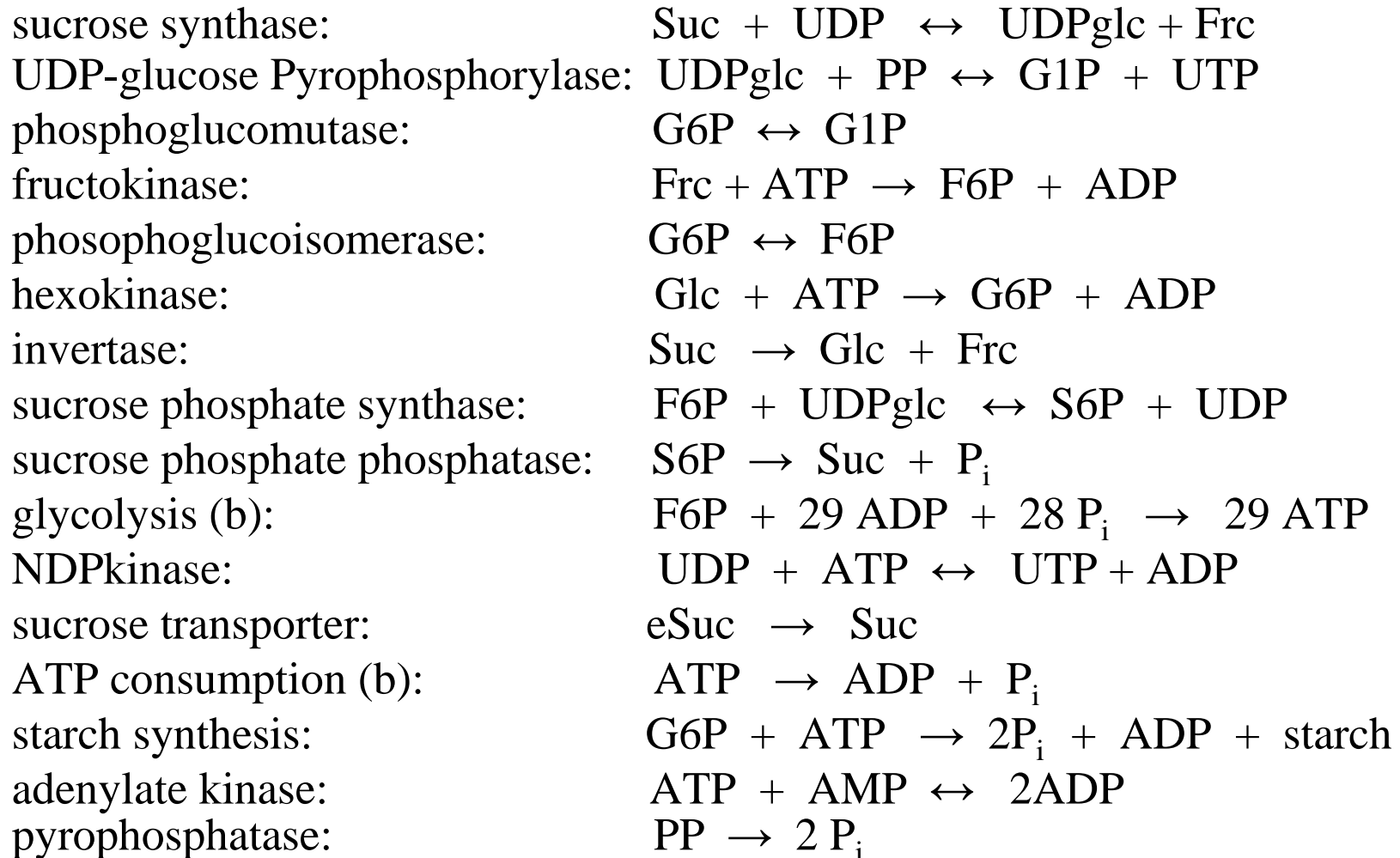
Token flow: occurring of an event (firing of a transition)

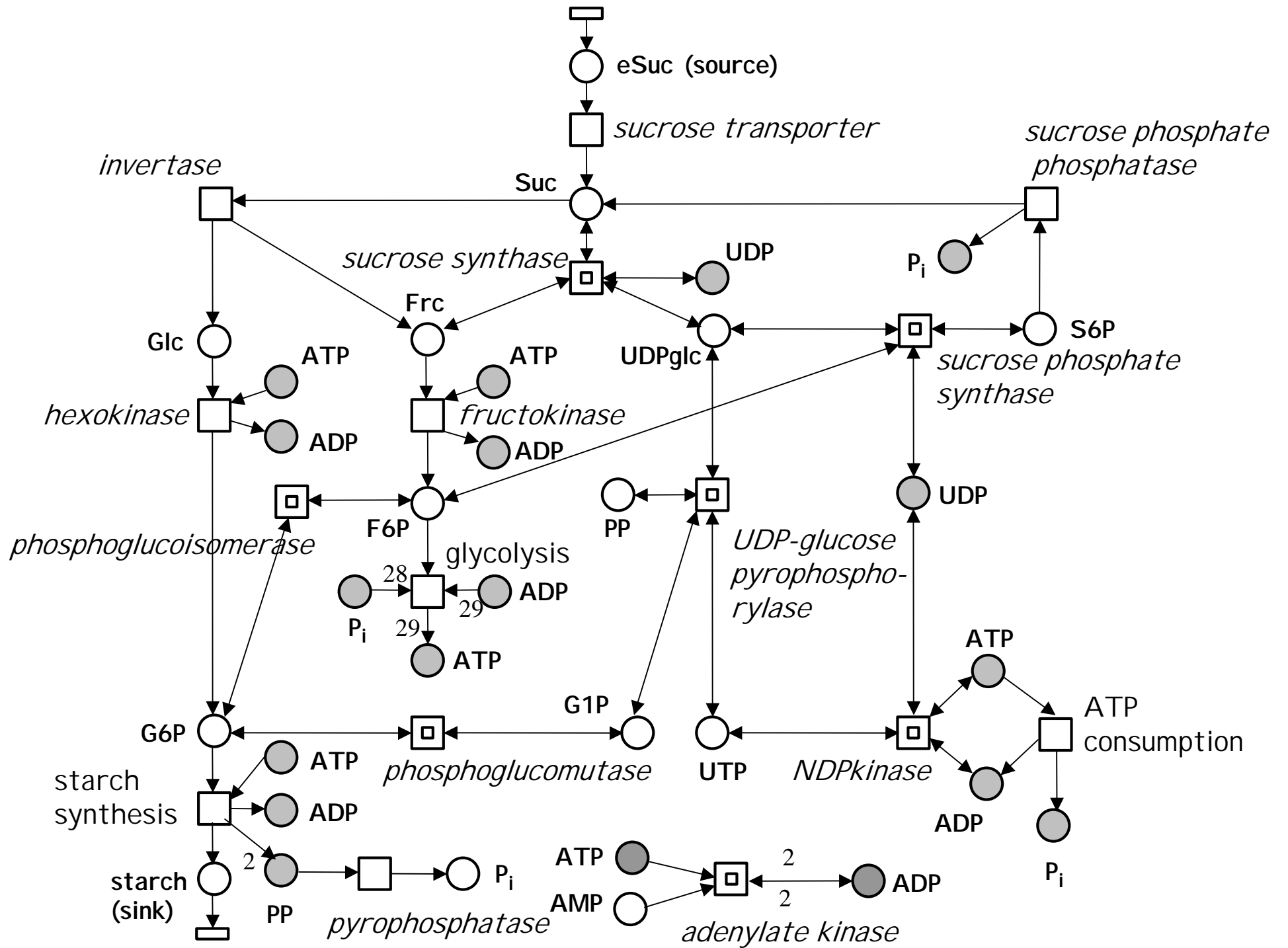
etri net basics

Example: Pentose Phosphate Pathway - sum reaction



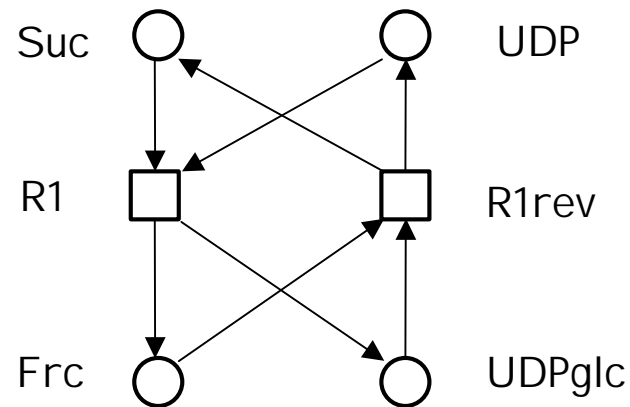
ucrose-to-starch-pathway in potato tuber



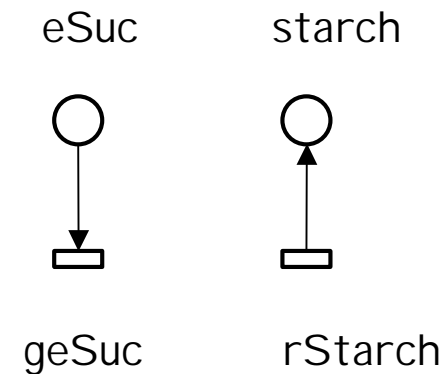


ucrose-to-starch-pathway in potato tuber

A hierarchical node:



Interface to the environment



Tools:

Editing: Ped Heiner BTU Cottbus

Animation: PedVisor <http://www.informatik.tu-cottbus.de/~wwwdssz/>

Qualitative analysis: INA Starke HU Berlin

[http://www.informatik.hu-](http://www.informatik.hu-berlin.de/~starke/ina.html)

[berlin.de/~starke/ina.html](http://www.informatik.hu-berlin.de/~starke/ina.html)

odel validation

- (1) Dynamical (behavioural) properties
- (2) Reachability analysis
- (3) Structural analysis
- (4) Invariant analysis
- (5) Model checking

ynamic (behavioural) properties

Liveness and Reversibility

- a net is live, if all its transitions are live in the initial marking
- a net is reversible, if the initial marking can be reached from each possible state
- How often can a transition fire? (0-times, n-times, ∞ times)
- infinite systems behaviour, search for dead transitions
- prediction of system deadlocks

ynamic (behavioural) properties

Boundedness

- a net is bounded, if there exists a positive integer number k , which represents a maximal number of tokens on each place in all states
- What is the maximal token number for a place?
(0, 1, k , ∞) boundedness (k-bounded)
- for bounded nets special algorithms exist

reachability analysis

How many and which system states could be reached ? (0, 1, k, ∞)

- the reachability graph represents all possible states
- computational problem for large and dense biological networks
- for unbounded networks: computation of the coverability graph
- Is a certain system state again and again reachable?
progressiveness
- Is a certain system state never reachable? safety

structural analysis

- aims at discovering net structures to derive conclusions on dynamic properties

Elementary properties:

ordinary: the multiplicity of every arc is equal one

homogeneous: for any place all outgoing edges have the same multiplicity

pure: there is no transition, for which a pre-place is also a post-place (loop-free)

conservative: for each place the sum of input arc weights is equal to the sum of output arc weights – a conservative net is bounded

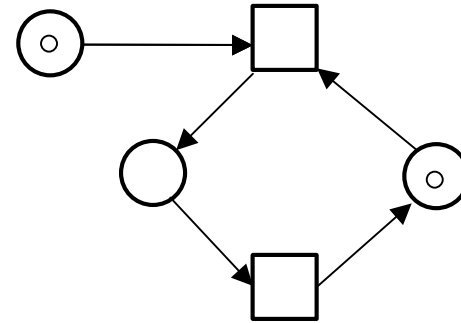
static conflict-free: there are no transitions with a common pre-place

connected, strongly connected: in graph-theoretical sense

structural analysis

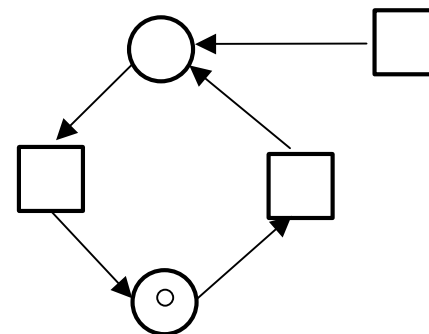
structural deadlock:

a set of places that delivers its tokens until a state is reached, where the place set is empty and there is no possibility to get a new token



trap:

the opposite situation that tokens cannot be removed from a place set (accumulation of substances)



invariant analysis

- properties, which are conserved during the working of the system
- independent of the initial marking
- only the net structure is relevant for their calculation

Are there invariant structures, which are independent from firing of the system?

Place-invariants (P-invariants)
Transition-invariants (T-invariants)

P-invariant analysis

Interpretation

P-invariants

- set of places, whose weighted sum of tokens is constant
- covered by P-invariants: sufficient condition for boundedness
liveness
- set of metabolites, whose total net concentrations remain unchanged
ADP, ATP
NADP⁺, NADPH

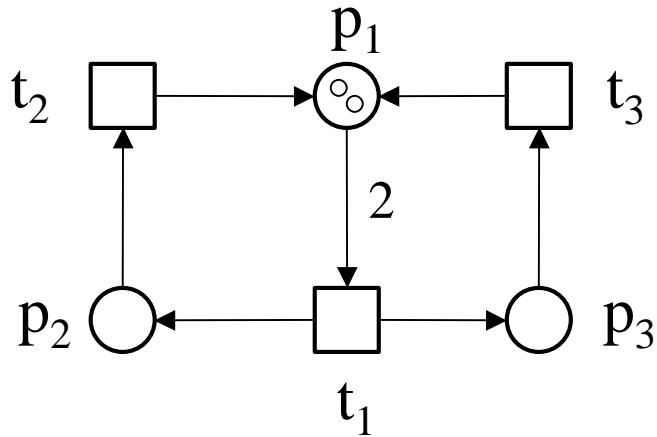
T-invariants

- set of transitions, whose firing reproduces a given marking
- covered by T-invariants: necessary condition for
- minimal set of enzymes which could operate at steady state
- indicate the presence of cyclic firing sequences

Elementary modes

Schuster, Hilgetag, Schuster (1993)

invariant analysis



incidence matrix $C = P \times T$

$$C = \begin{matrix} & t_1 & t_2 & t_3 \\ \begin{matrix} p_1 \\ p_2 \\ p_3 \end{matrix} & \begin{pmatrix} -2 & 1 & 1 \\ 1 & -1 & 0 \\ 1 & 0 & -1 \end{pmatrix} \end{matrix}$$

place (P-) invariant:

$$x C = 0$$

$$-2x_1 + x_2 + x_3 = 0$$

$$x_1 - x_2 = 0$$

$$x_1 - x_3 = 0$$

transition (T-) invariant:

$$C y = 0$$

$$-2y_1 + y_2 + y_3 = 0$$

$$y_1 - y_2 = 0$$

$$y_1 - y_3 = 0$$

nvariant analysis

Minimal semi-positive solutions are of interest with

- all components of the solution vector are ≥ 0
- basis of the semi-positive solution space such that none solution is contained in another solution, Lautenbach (1973)

The calculation

- of all integer solutions is in P
- of all semi-positive solutions is in P
- of all semi-positive integer solutions is NP-complete, Schrijver (1999)

Extreme Pathways Schilling et al. (2000)

- minimal basis of semi-positive integer solutions
- subset of T-invariants – biological interpretation?

Qualitative analysis using INA

Elementary properties

The net is not statically conflict-free.

The net has transitions without pre-place.

The net is not covered by semipositive P-invariants.

The net is not structurally bounded.

The net is not safe.

The net has transitions without post-place.

The net is not ordinary.

At least the following transitions are live: 0.SucTrans, 1.Inv, 18.geSuc,

At least the following places are simultaneously unbounded: 0.Suc, 1.eSuc, 2.Glc, 3.Frc,

The net is marked.

The net is not homogenous.

The net has no nonempty clean trap.

The net has no places without post-transition.

The net is connected.

The net is pure.

The net is not strongly connected.

The net is not bounded.

The net is not live and safe.

Transition 18.geSuc has no pre-place.

Transition 21.rStarch has no post-place.

The net is not conservative.

The net is not marked with exactly one token

The net has not a non-blocking multiplicity

The net has no places without pre-transition

Maximal in/out-degree: 6

ORD	HOM	NBM	PUR	CSV	SCF	CON	SC	Ft0	tF0	Fp0	pF0	MG	SM	FC	EFC	ES
N	N	N	Y	N	N	Y	N	Y	Y	N	N	N	N	N	N	N

Qualitative analysis using INA

Structural properties

DTP	CPI	CTI	B	SB	REV	DSt	BSt	DTr	DCF	L	LV	L&S
?	N	Y	N	N	?	?	?	?	?	?	?	N

- liveness could not be decided because the net is unbounded and the reachability graph cannot be calculated
- the coverability graph has more than 4 million states

smaller bounded version: more than 10^{10} states of the reachability graph

Invariant analysis

The net is not covered by P-invariants.

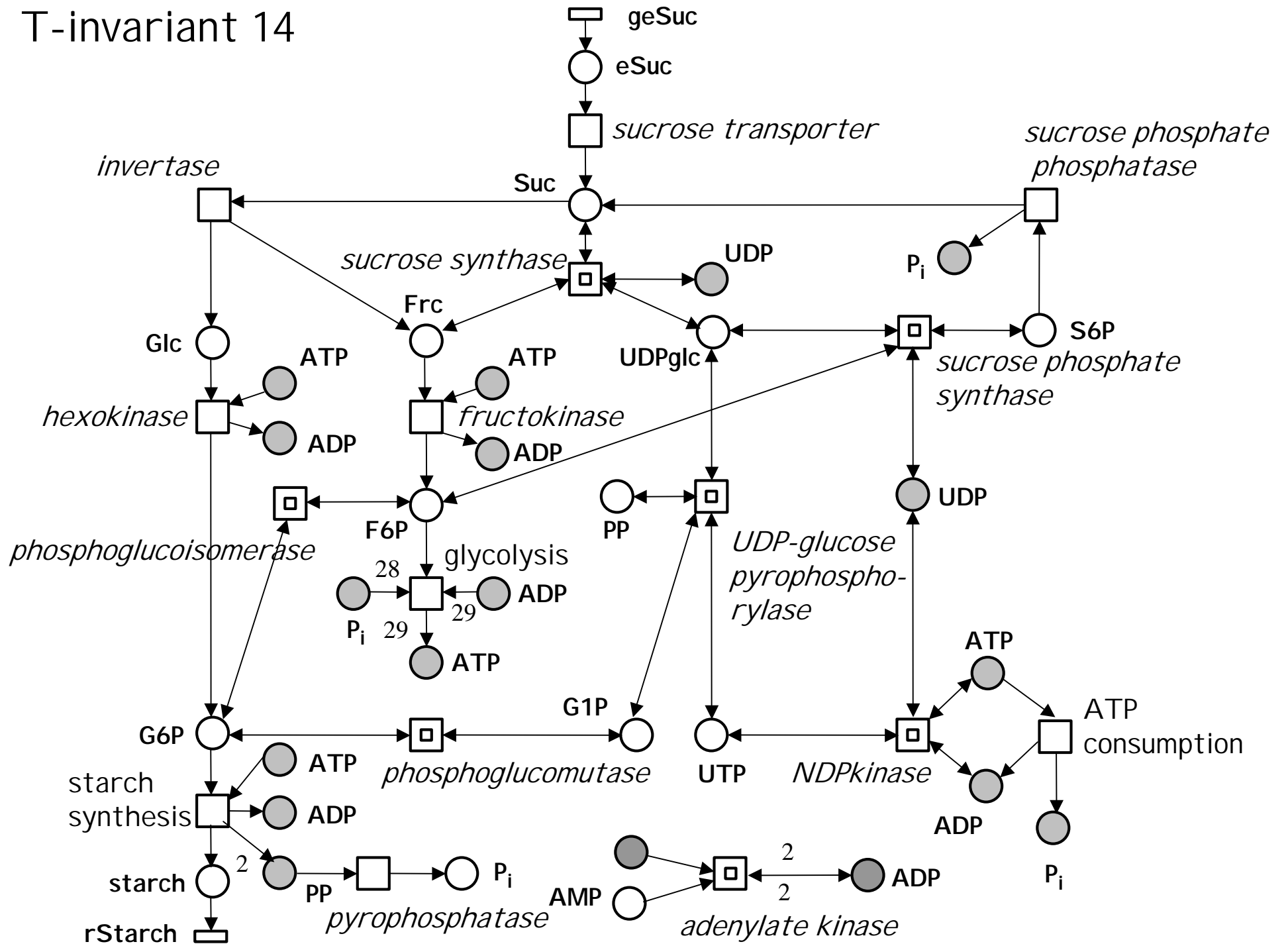
Following P-invariants were calculated:

1. UDPglc, UTP, UDP
2. ATP, AMP, ADP
3. G6P, F6P, G1P, UTP, ATP(2), ADP, S6P, P_i, PP(2)

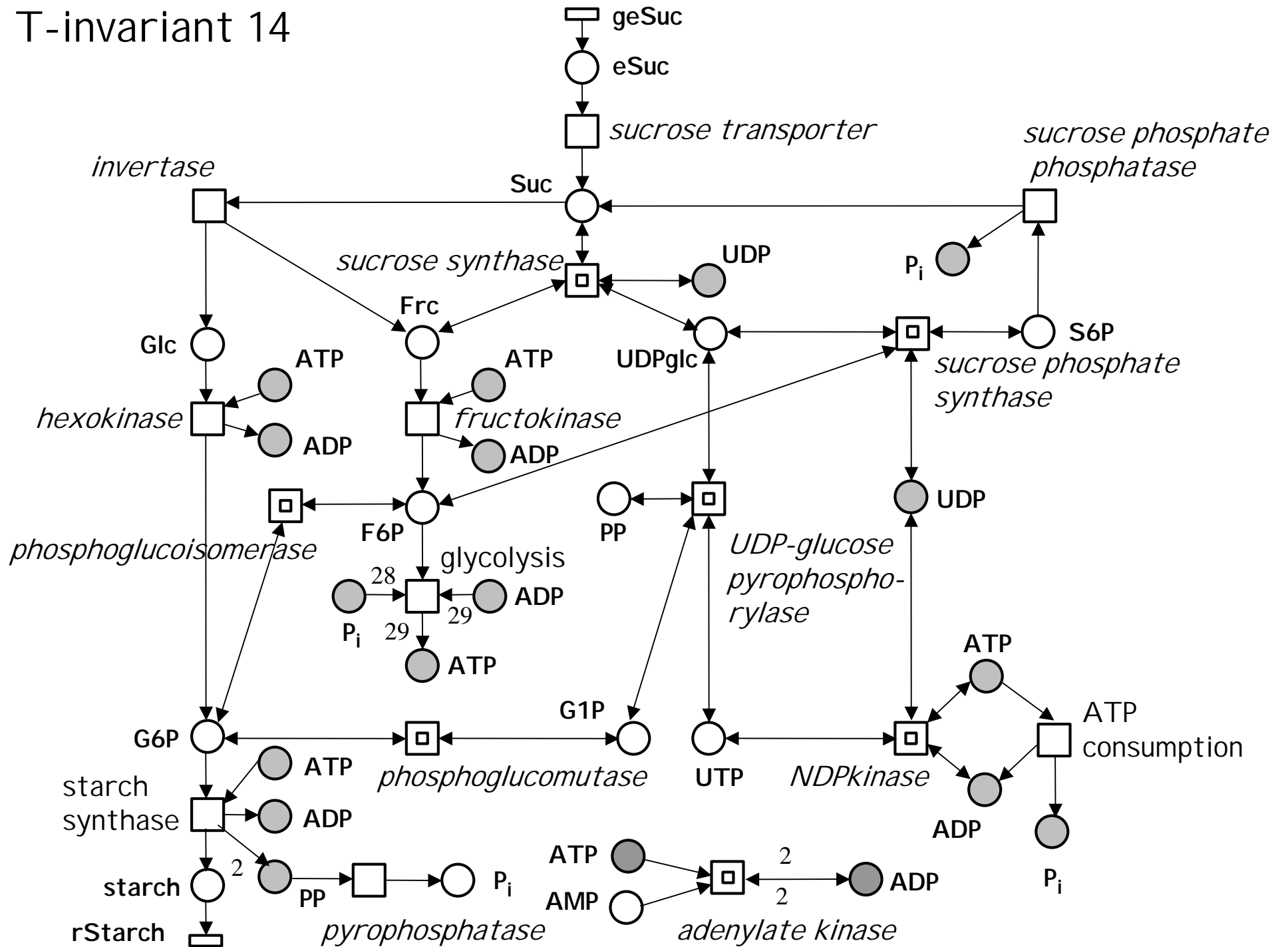
The net is covered by 19 T-invariants

- 7 trivials:
1. SPS, SPS_rev,
 2. UGPase, UGPASE_rev,
 3. SuSy_SuSy_rev,
 4. PGM, PGM_rev,
 5. NDPkin, NDPkin_rev,
 6. AdK, AdK_rev,
 7. PGI, PGI_rev

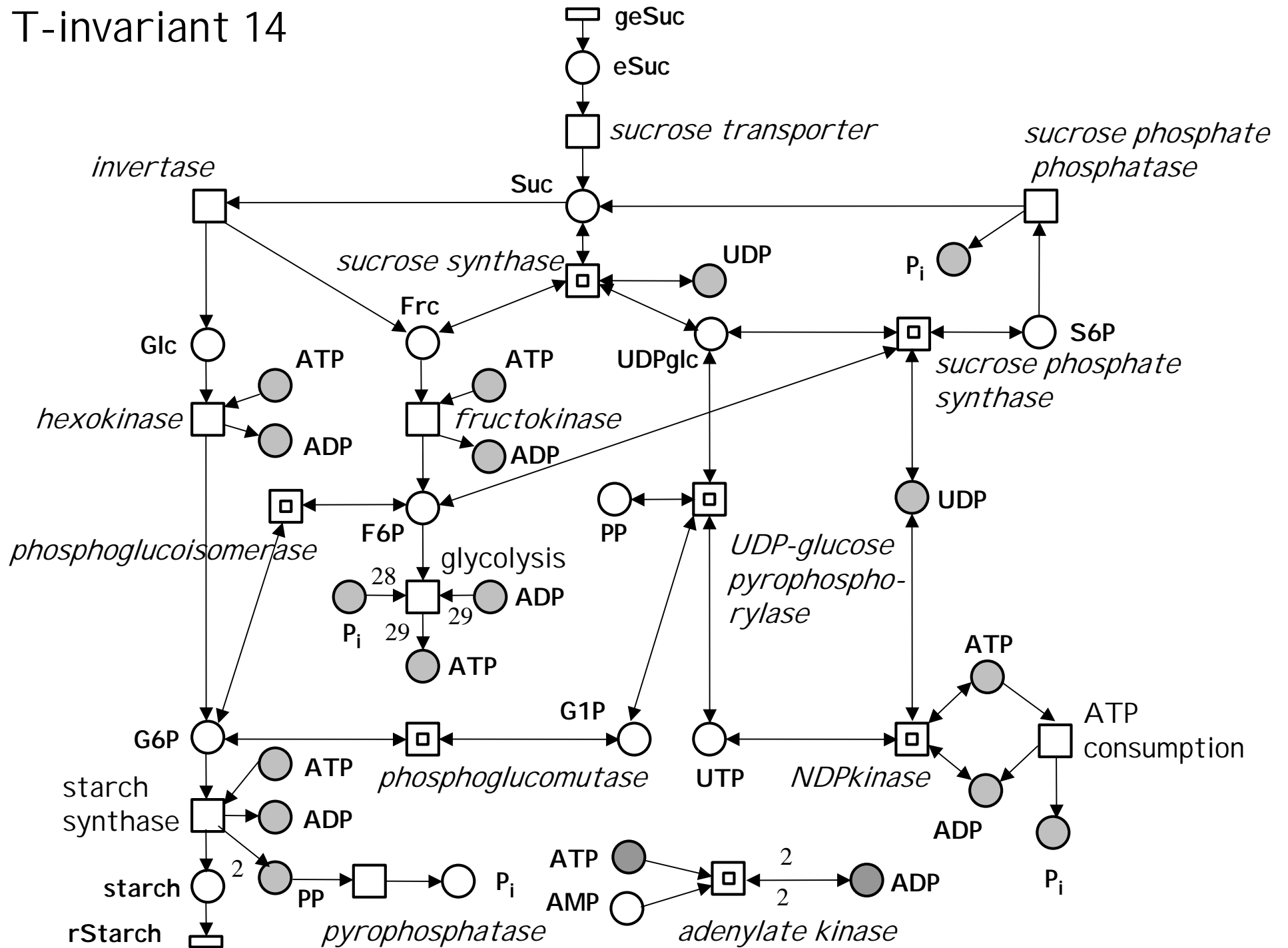
T-invariant 14



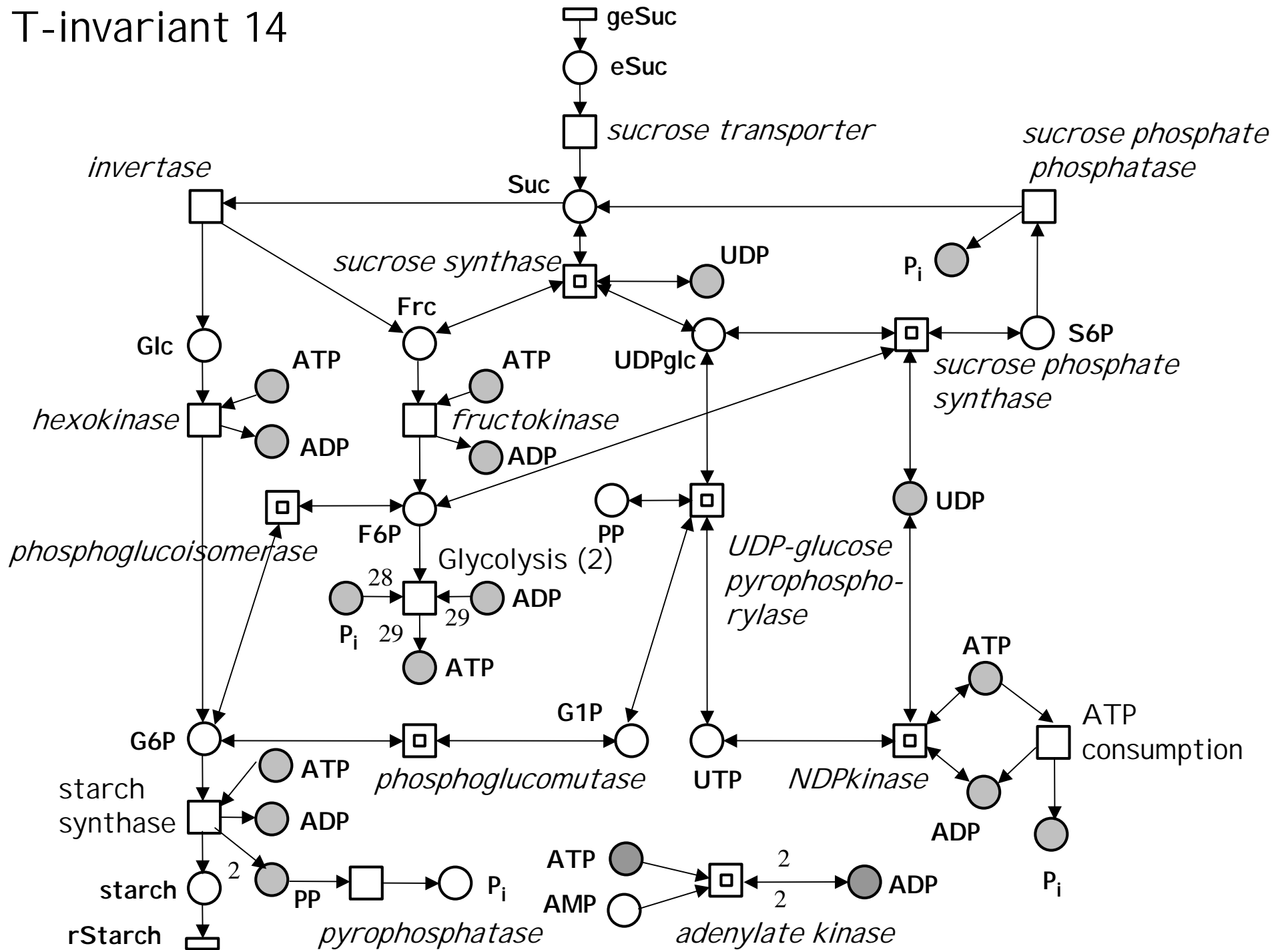
T-invariant 14



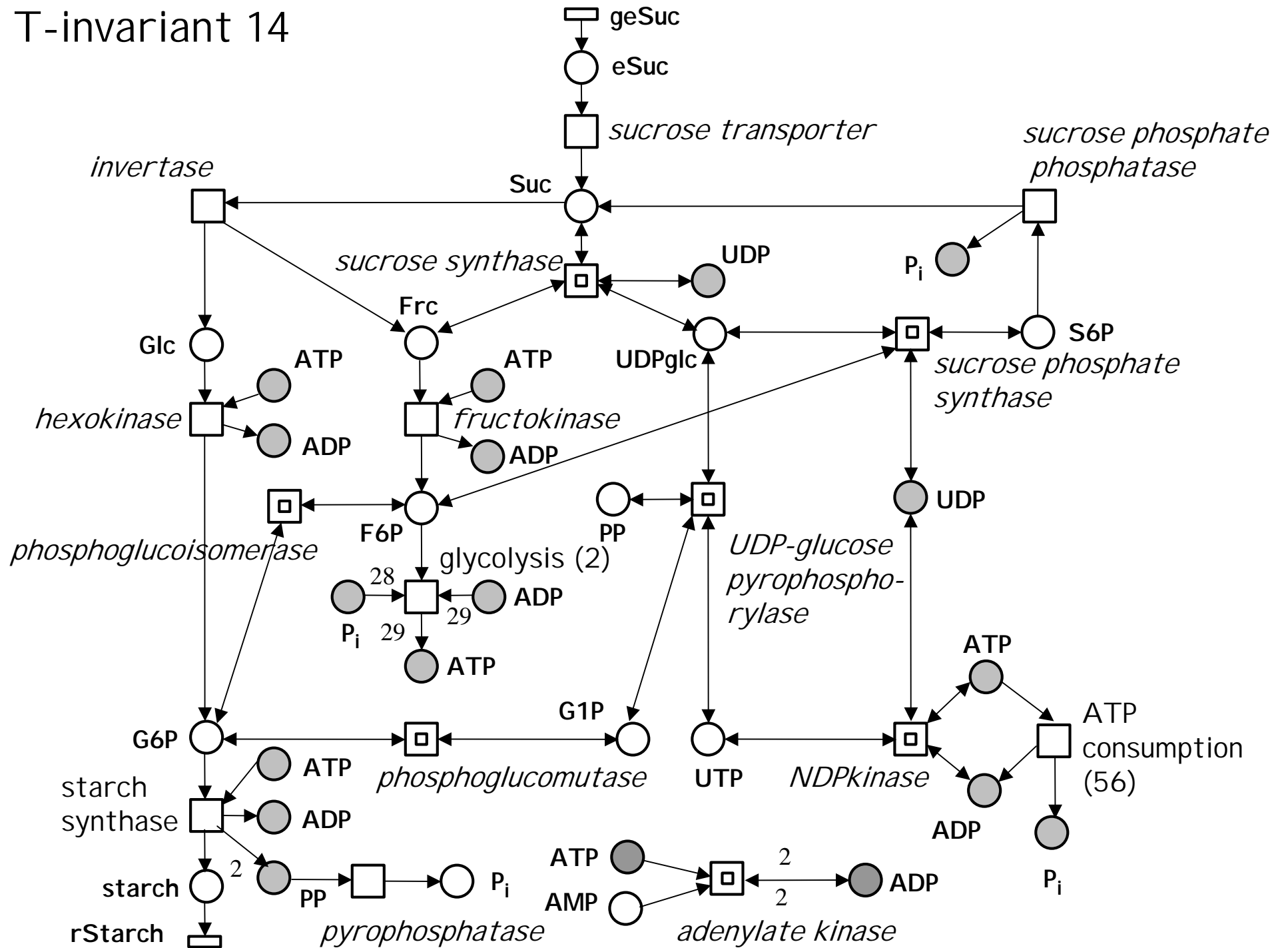
T-invariant 14



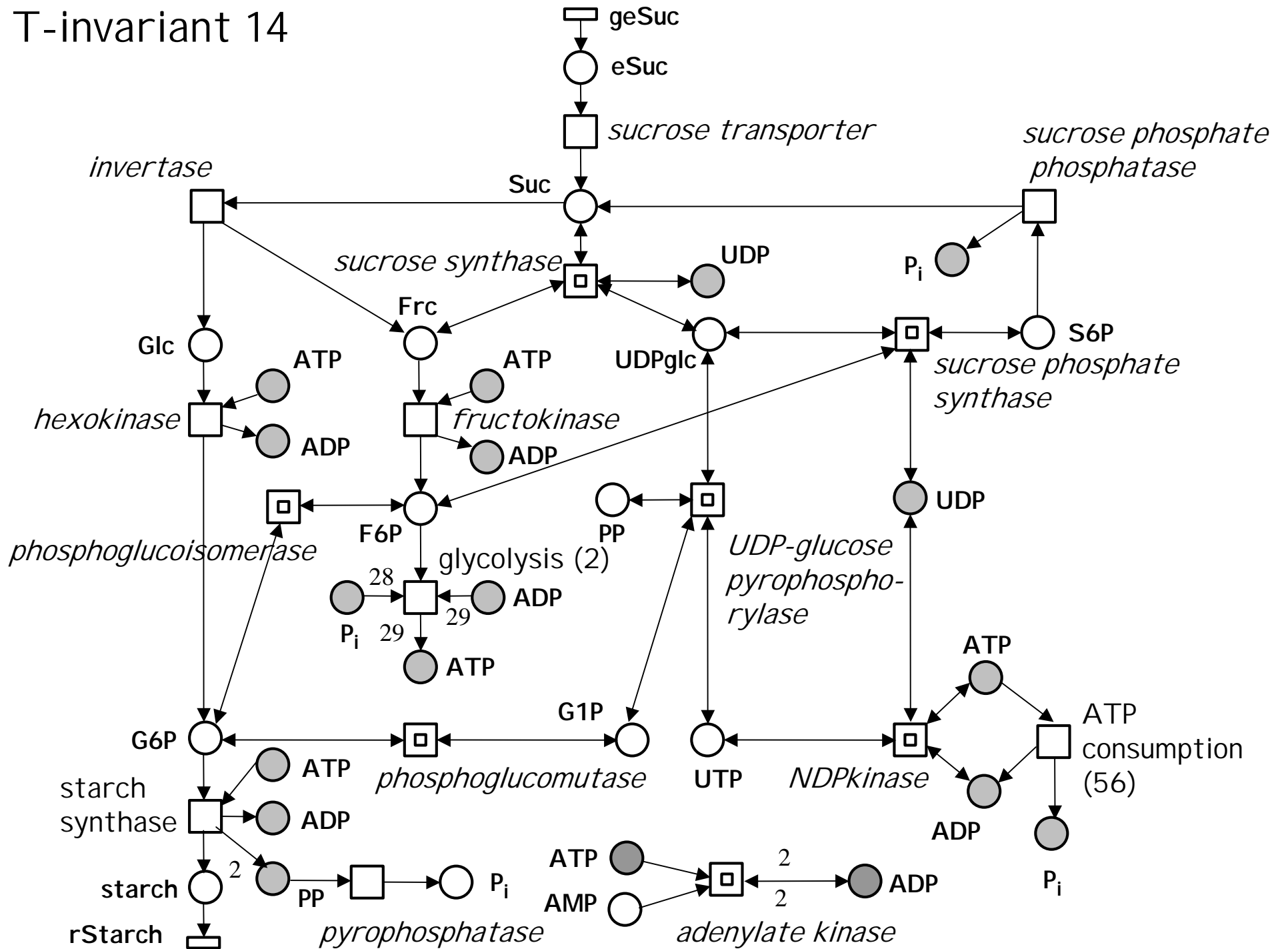
T-invariant 14



T-invariant 14



T-invariant 14



T-Invariant analysis

21	0.gGlu	:	1,
	1.R1	:	1,
	2.gADP	:	2,
	3.gPi	:	2,
	4.gNAD	:	2,
	5.PGIsomerase:		1,
	7.Aldolase1	:	1,
	10.TPIsomerase:		1,
	11.GAPDH	:	2,
	13.PGKinase	:	2,
	15.PGMutase1	:	2,
	17.Enolase	:	2,
	19.rNADH	:	2,
	26.rATP	:	2,
	27.rPyr	:	2,
	28.PFKinase	:	1,
	29.PyrKinase1	:	1

T-Invariant analysis

21	0.gGlu	:	1,	Substrates: Glucose, 2ADP, 2P _i , 2NAD ⁺
	1.R1	:	1,	
	2.gADP	:	2,	
	3.gPi	:	2,	
	4.gNAD	:	2,	
	5.PGIsomerase:		1,	
	7.Aldolase1	:	1,	
	10.TPIsomerase:		1,	
	11.GAPDH	:	2,	
	13.PGKinase	:	2,	
	15.PGMutase1	:	2,	
	17.Enolase	:	2,	
	19.rNADH	:	2,	
	26.rATP	:	2,	
	27.rPyr	:	2,	
	28.PFKinase	:	1,	
	29.PyrKinase1	:	1	

T-Invariant analysis

21	0.gGlu	:	1,	
	1.R1	:	1,	Substrates:
	2.gADP	:	2,	Glucose, 2ADP, 2P _i ,
	3.gPi	:	2,	2NAD ⁺
	4.gNAD	:	2,	
	5.PGIsomerase:		1,	
	7.Aldolase1	:	1,	Products:
	10.TPIsomerase:		1,	2Pyruvate, 2ATP, 2NADH
	11.GAPDH	:	2,	
	13.PGKinase	:	2,	
	15.PGMutase1	:	2,	
	17.Enolase	:	2,	
	19.rNADH	:	2,	
	26.rATP	:	2,	
	27.rPyr	:	2,	
	28.PFKinase	:	1,	
	29.PyrKinase1	:	1	

T-Invariant analysis

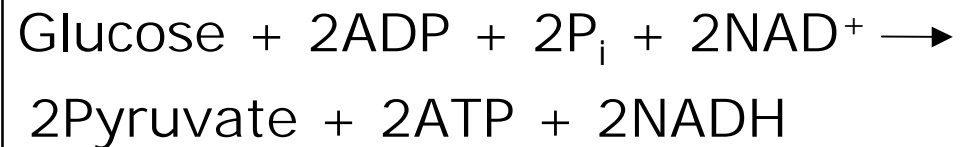
21 | 0.gGlu : 1,
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| 2.gADP : 2,
| 3.gPi : 2,
| 4.gNAD : 2,
| 5.PGIsomerase: 1,
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| 13.PGKinase : 2,
| 15.PGMutase1 : 2,
| 17.Enolase : 2,
| 19.rNADH : 2,
| 26.rATP : 2,
| 27.rPyr : 2,
| 28.PFKinase : 1,
| 29.PyrKinase1 : 1

Substrates:

Glucose, 2ADP, 2P_i,
2NAD⁺

Products:

2Pyruvate, 2ATP, 2NADH



T-Invariant analysis

Invariant number	sucrose cleavage		hexoses go into		ATP cons	ATP used for cycling		
	SuSy	Inv	Glyc	StaSy		Inv SuSy_rev	Inv SPS, SPP	SuSy SPS, SPP
8	x		x	x				x
9	x		x	x	x			
10	x		x	x				
11		x	x	x		x		
12		x	x	x			x	
13		x	x	x				x
14		x	x	x	x			
15		x	x	x				
16		x	x			x		
17		x	x				x	
18		x	x					x
19		x	x		x			

Robustness

Robustness: sensitivity of the system against parameter (fragility) changes (altered enzyme activity, mutations)
(Voit, 2000)

Stelling et al., Nature (2002): linear correlation between robustness and the number of elementary modes (T-invariants)

Our suggestion: - enzyme distribution over T-invariants
- number of alternative paths

Potato net: - fructokinase occurs in all T-invariants
- there is no enzyme that occurs in only one T-invariant

Conclusions & Outlook

Petri nets provide

- (1) a unique description of biological networks
- (2) methods for qualitative analysis to check models by the calculation of system properties.

- (3) The complexity of biological systems make it necessary to extend Petri net methods.
- (4) Automatic interpretation of T-invariants is necessary.