

VALIDATION & VERIFICATION OF SYSTEM DESIGN SPECIFICATIONS

**Monika Heiner
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- ❑ **V&V needs formal semantics**
 - > *model-based system design*

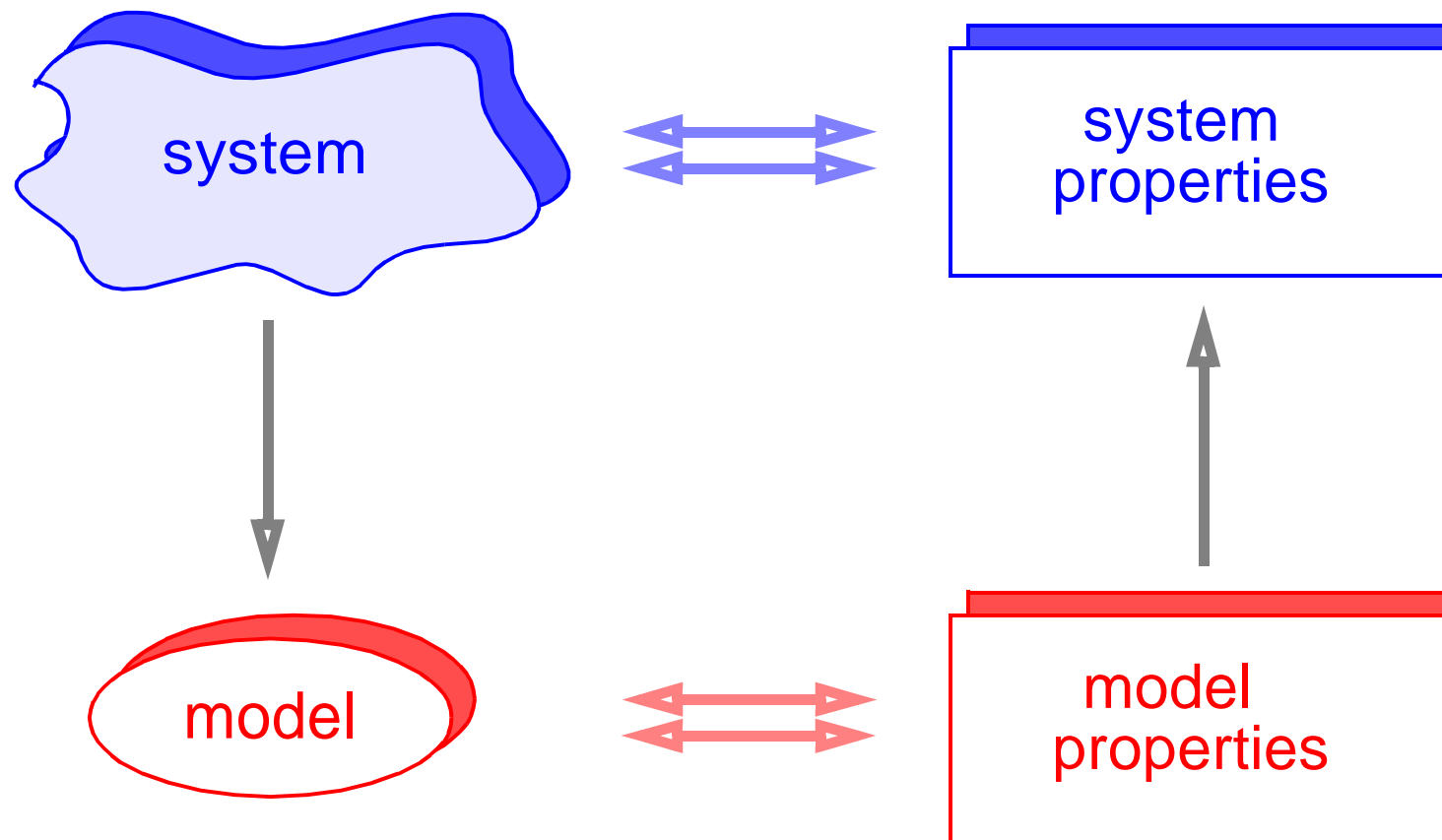
- ❑ **possible formal modelling techniques**
 - ... many ...*

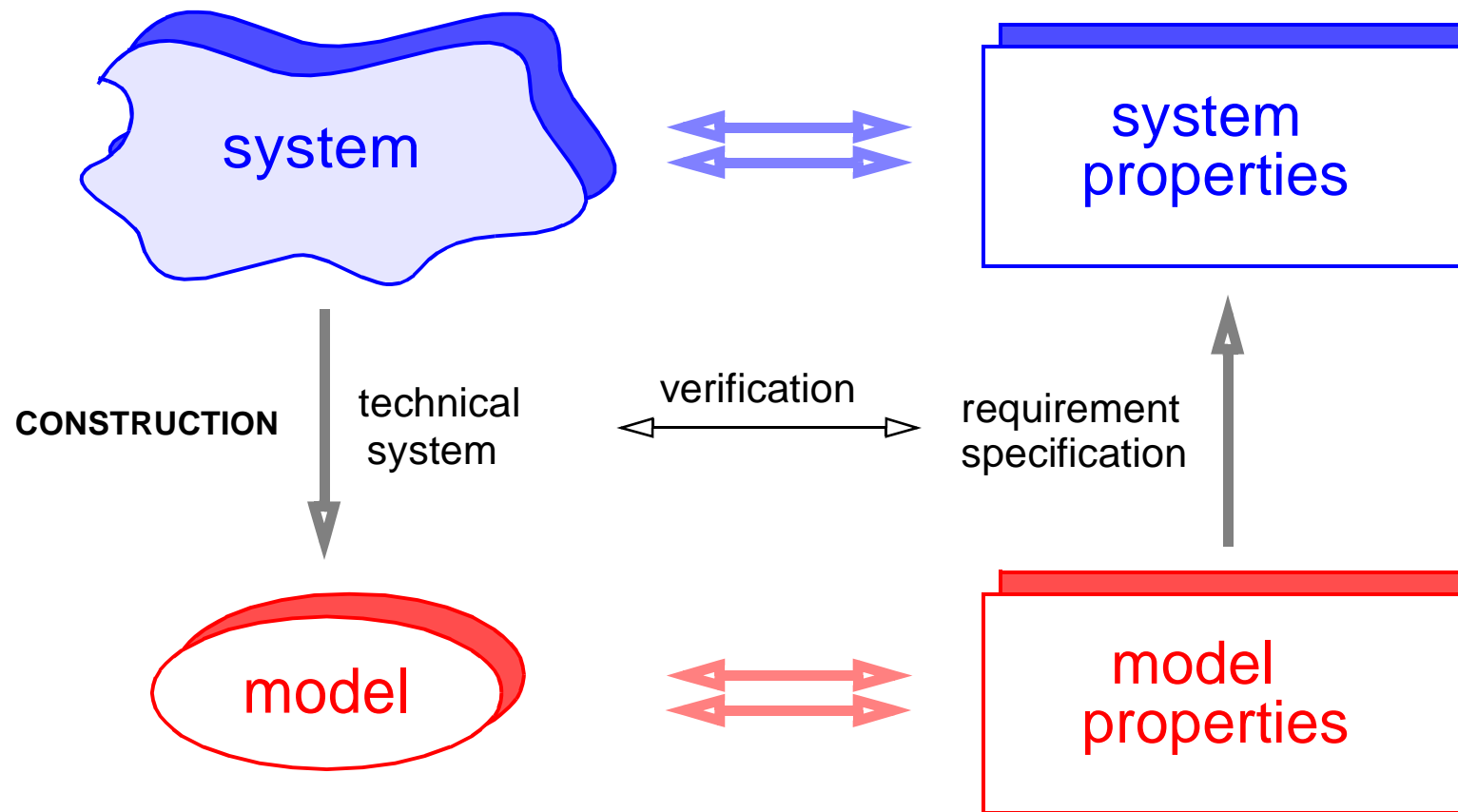
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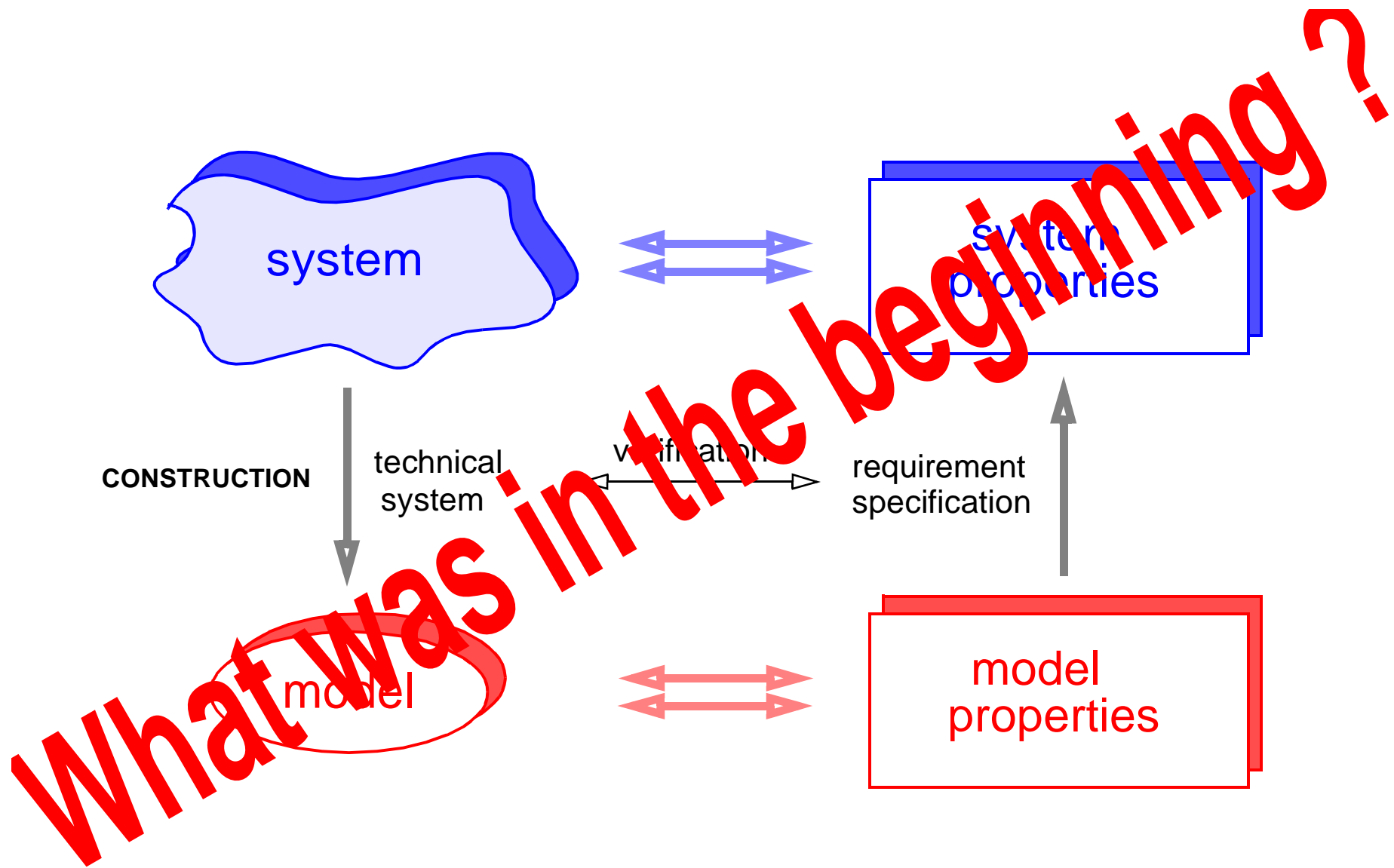
- ❑ **possible formal modelling techniques**
 - *many . . .*
 - > *Petri nets*

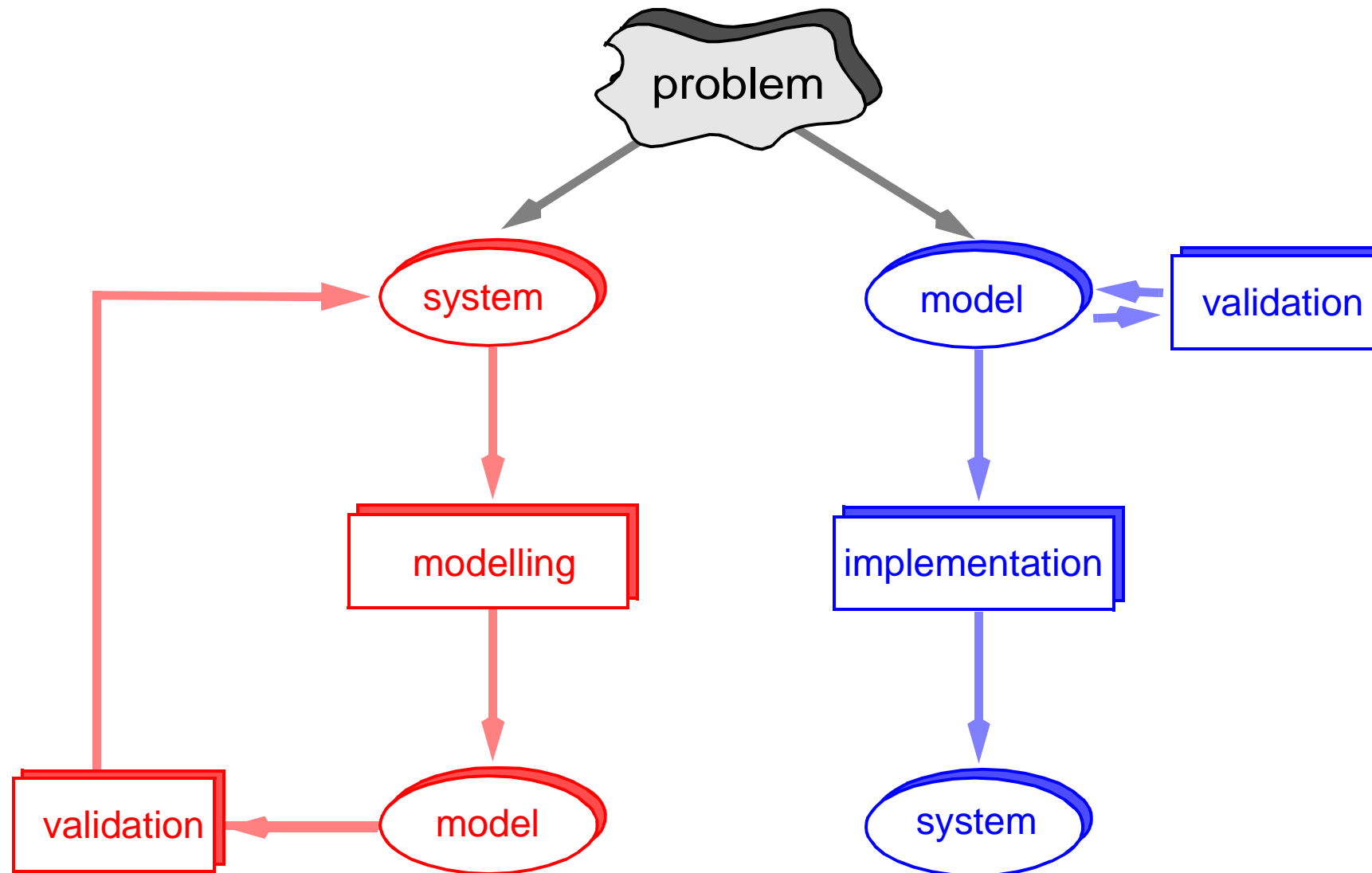
- ❑ **introduction into Petri nets**
 - > *modelling and analysis*
 - > *typical Petri net properties*
 - > *typical Petri net analysis techniques*

- ❑ **applications**
 - *many . . .*
 - > *error-correcting Petri nets*







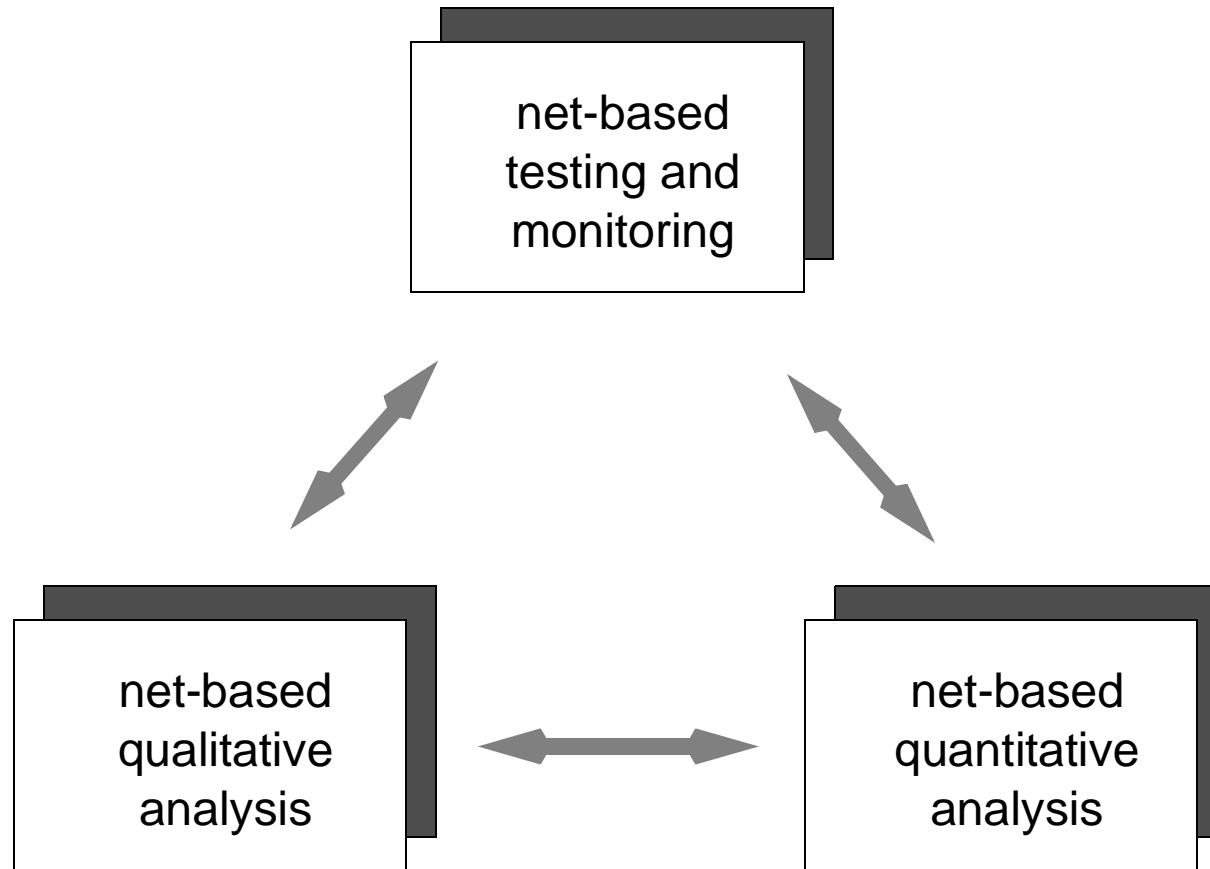


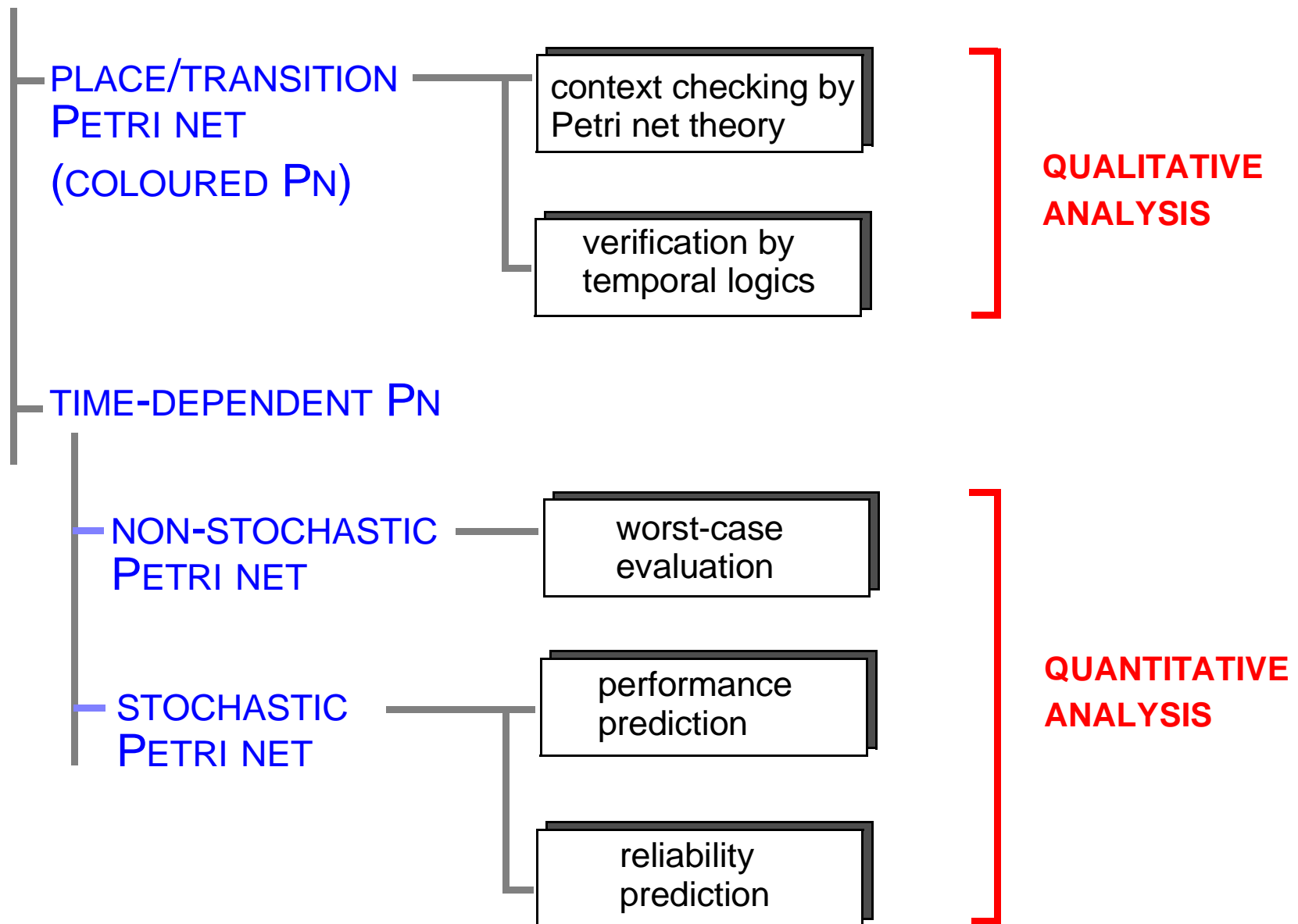
- ❑ **a suitable intermediate representation for**
 - > *different (specification/programming) languages,*
 - > *different phases of software development cycle,*
 - > *different validation methods;*

- ❑ **modelling power**
 - > *partial order (true concurrency) semantics*
 - > *applicable on any abstraction level*
 - > *specification of limited resources possible*

- ❑ **analyzing power**
 - > *combination of static and dynamic analysis techniques*
 - > *rich choice of methods, algorithms, tools*

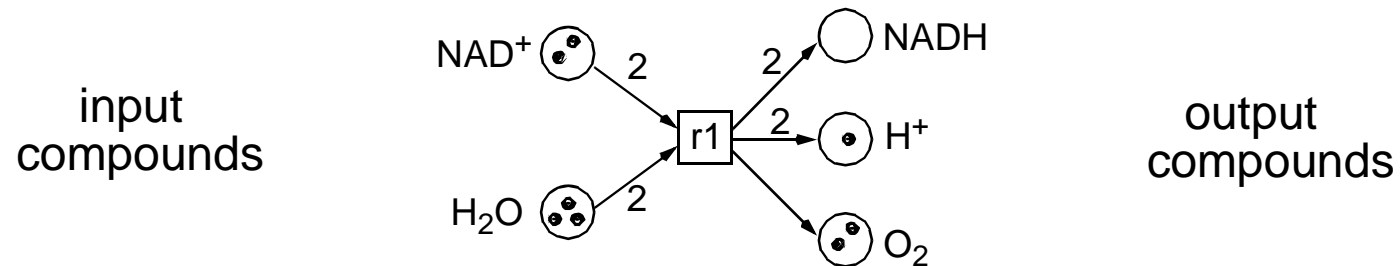
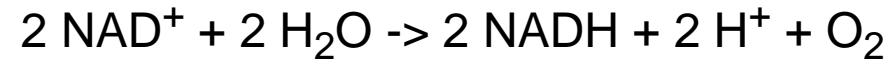
- ❑ **BUT: modelling power <-> analyzing power**



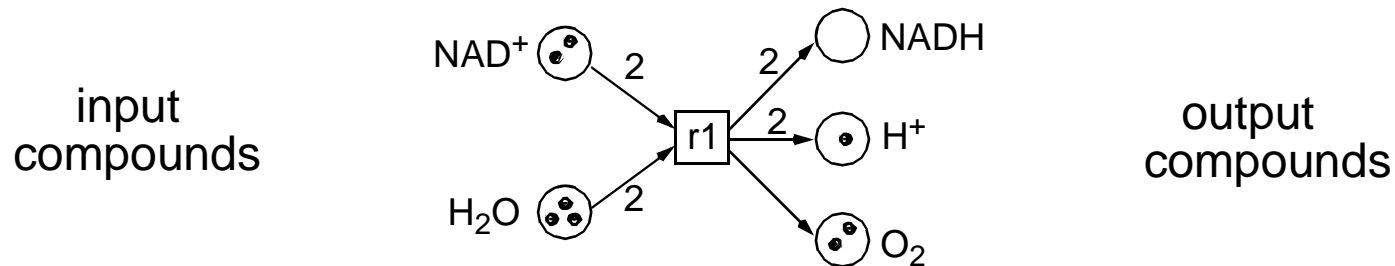
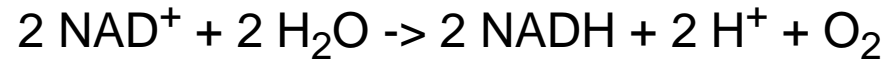


Petri nets, a crash course

□ atomic actions -> Petri net **transitions** -> chemical reactions

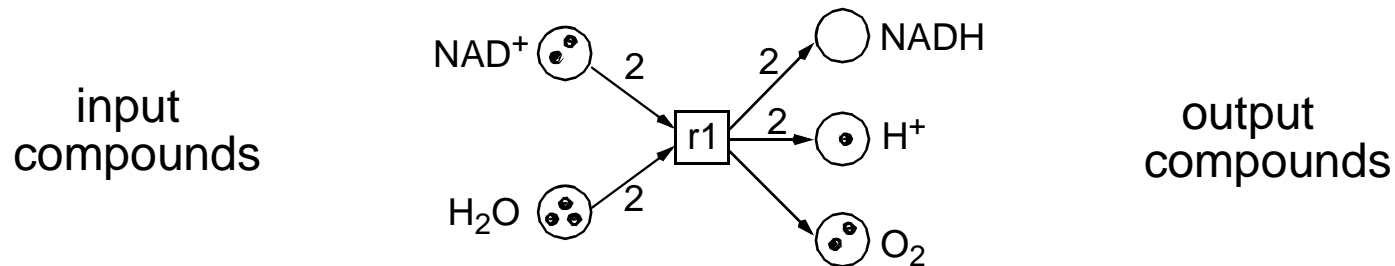
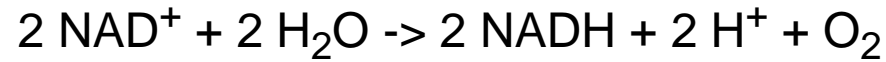


□ atomic actions -> Petri net **transitions** -> chemical reactions



□ local conditions -> Petri net **places** -> chemical compounds

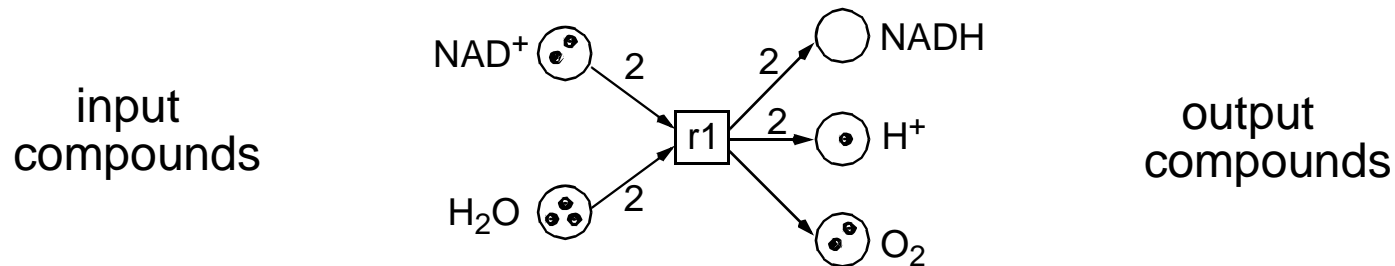
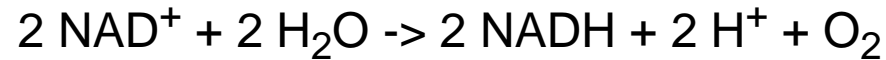
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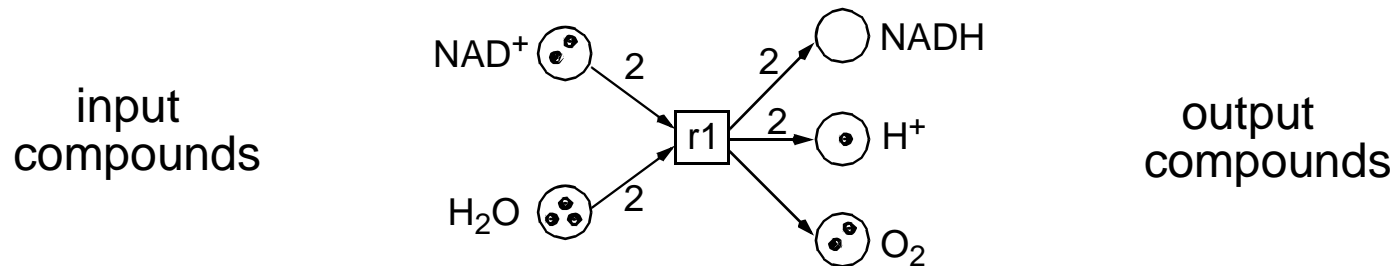
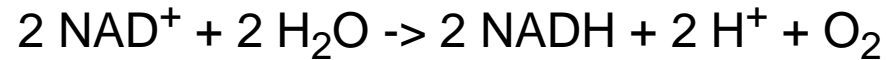


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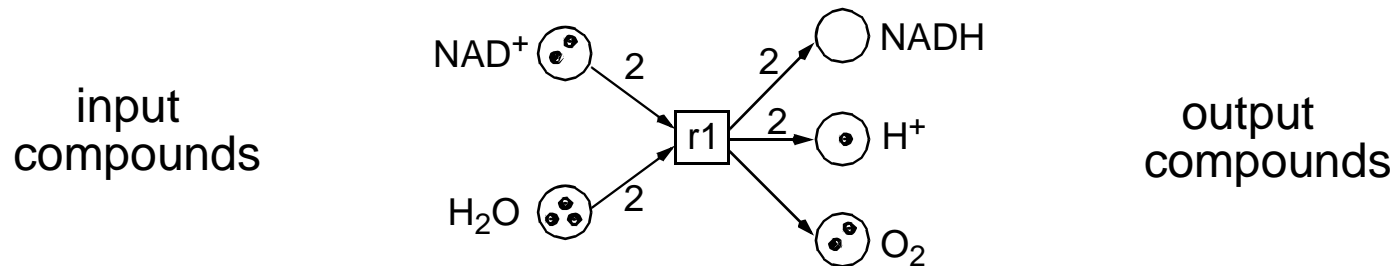
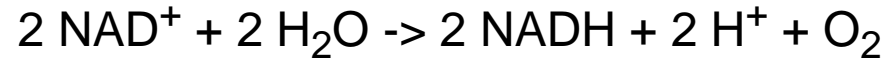
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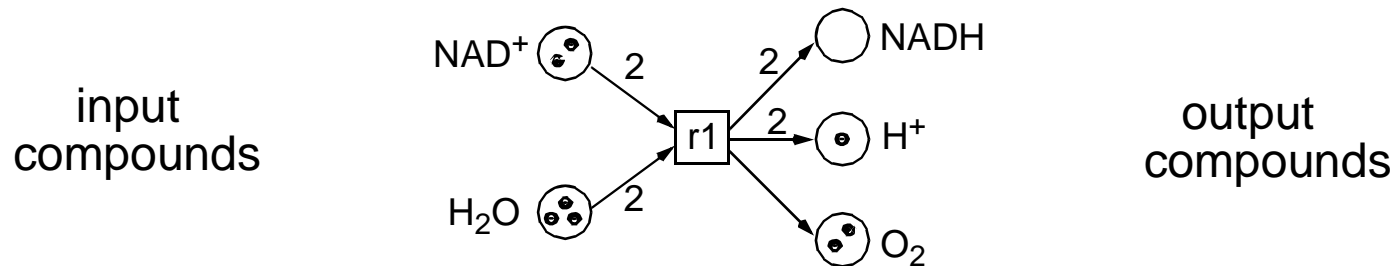
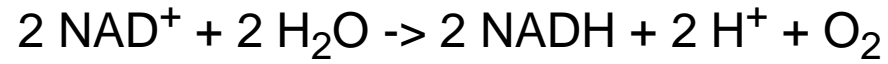
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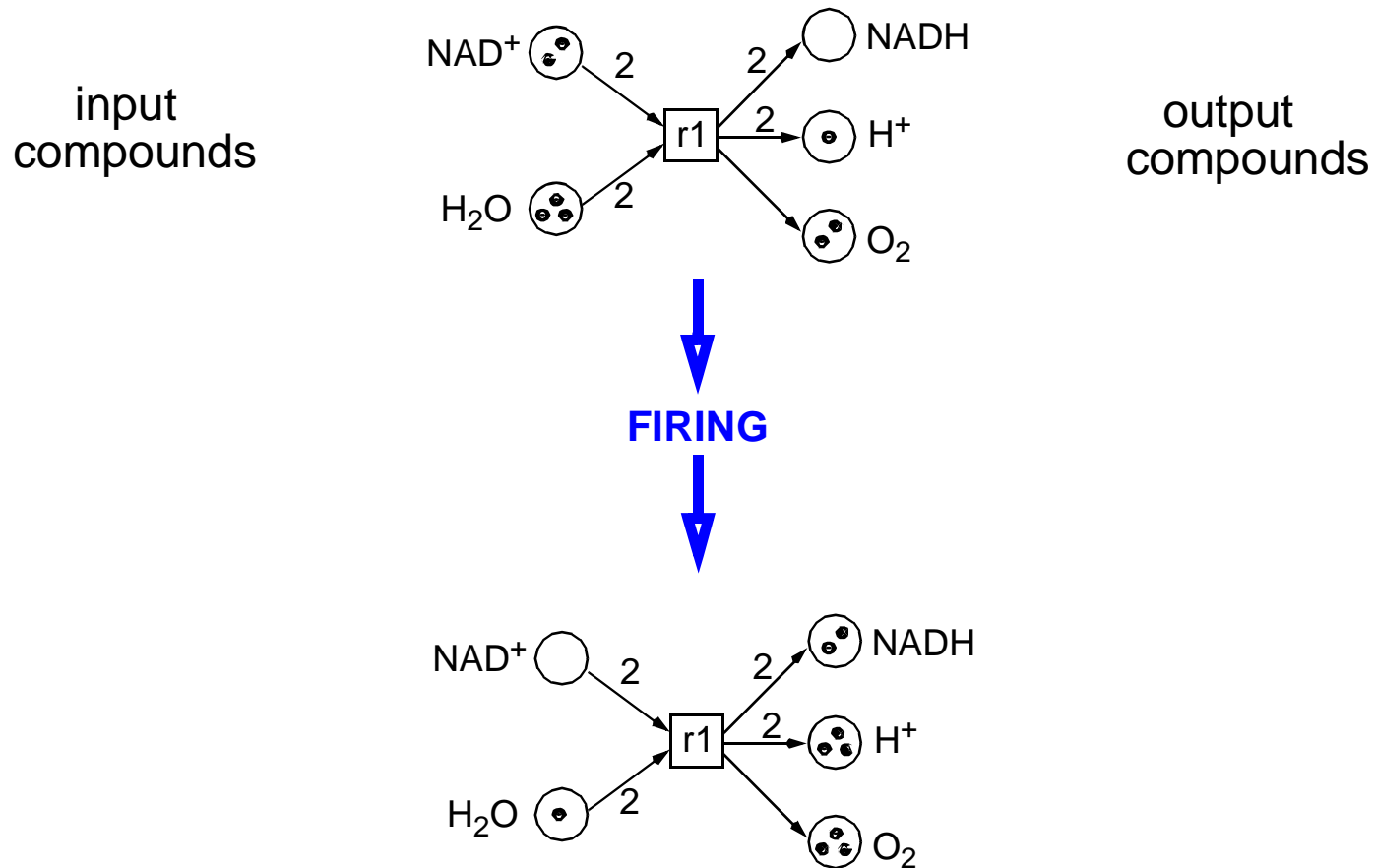
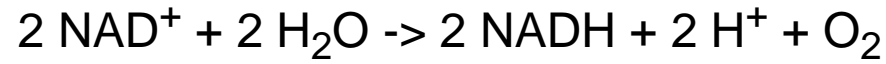
□ system state -> **marking** -> compounds distribution

□ $\text{PN} = (\text{P}, \text{T}, \text{F}, \text{m}_0)$, $\text{F}: (\text{P} \times \text{T}) \cup (\text{T} \times \text{P}) \rightarrow \mathbb{N}_0$, $\text{m}_0: \text{P} \rightarrow \mathbb{N}_0$

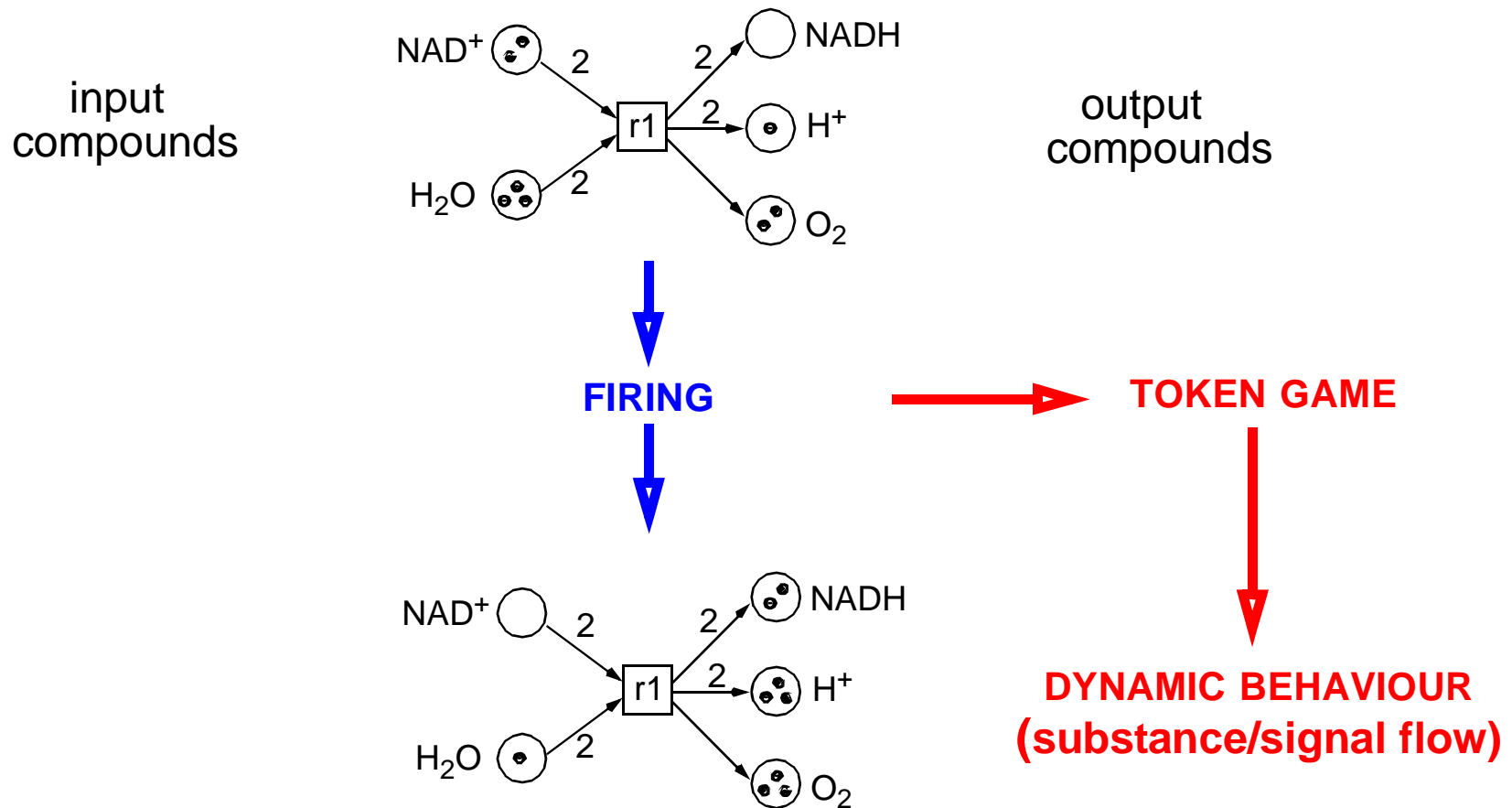
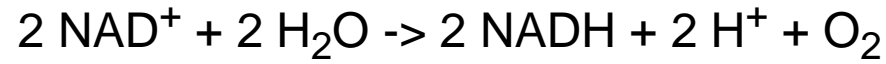
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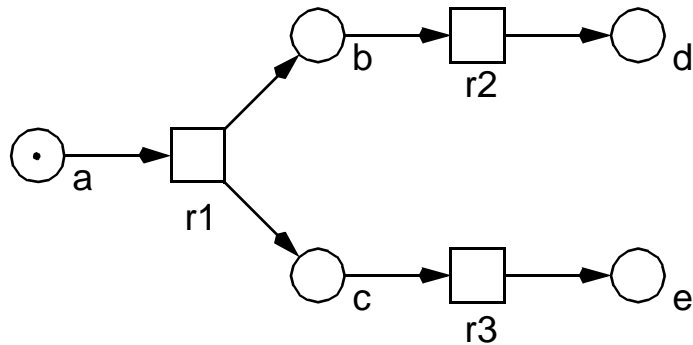


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- **order between r1 - r2 and r1 - r3**

-> *causality* $x < y [x-y]$

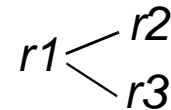
-> *dependency*

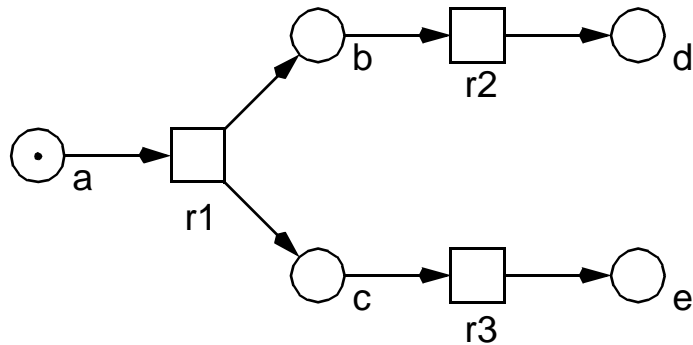
- **no order between r2 , r3**

-> *concurrency* $x || y$

-> *independency*

- **partial order run**





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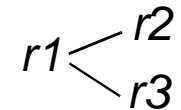
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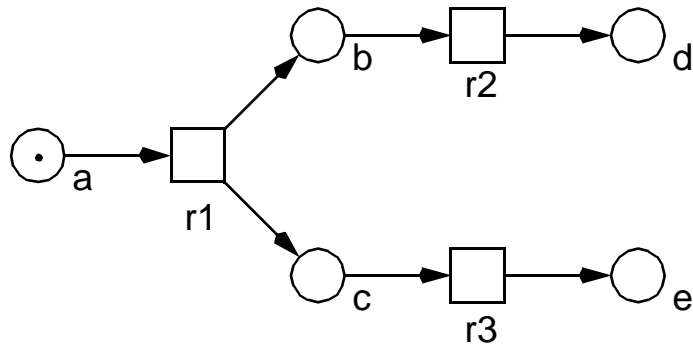
□ **partial order run**



-> **PARTIAL ORDER SEMANTICS**

“true concurrency semantics”

all partially ordered runs



□ possible interleaving runs

-> $r1 - r2 - r3$

-> $r1 - r3 - r2$

□ totally ordered runs

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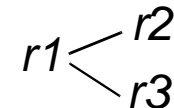
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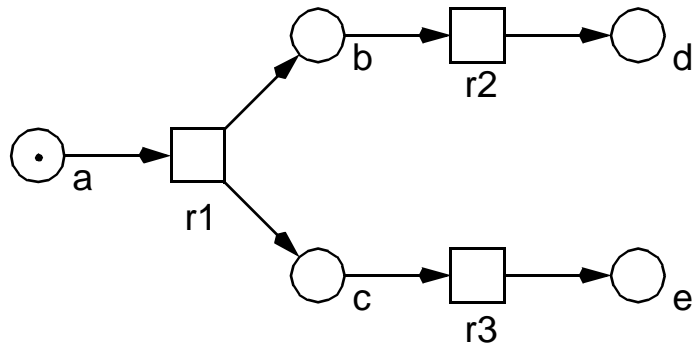
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□ possible interleaving runs

-> $r1 - r2 - r3$

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□ totally ordered runs

-> INTERLEAVING SEMANTICS

all totally ordered runs

□ order between $r1 - r2$ and $r1 - r3$

-> causality $x < y [x-y]$

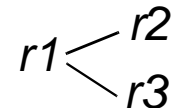
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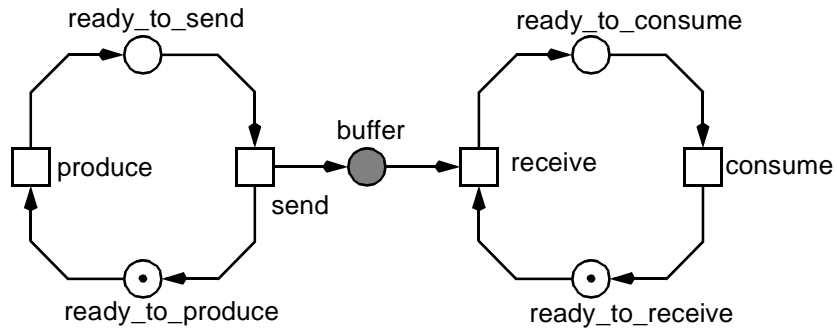


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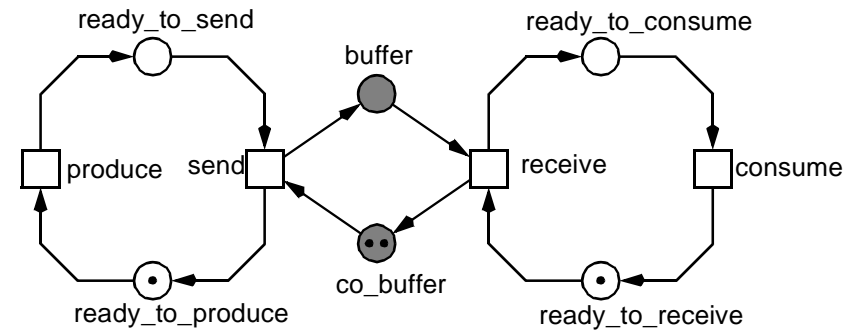
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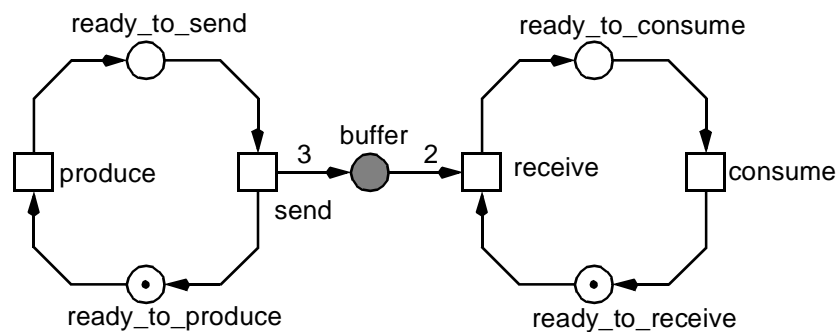
SYSTEM WITHOUT ARC WEIGHTS



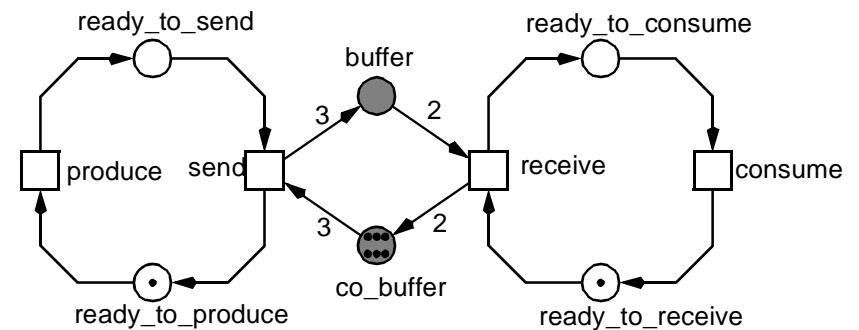
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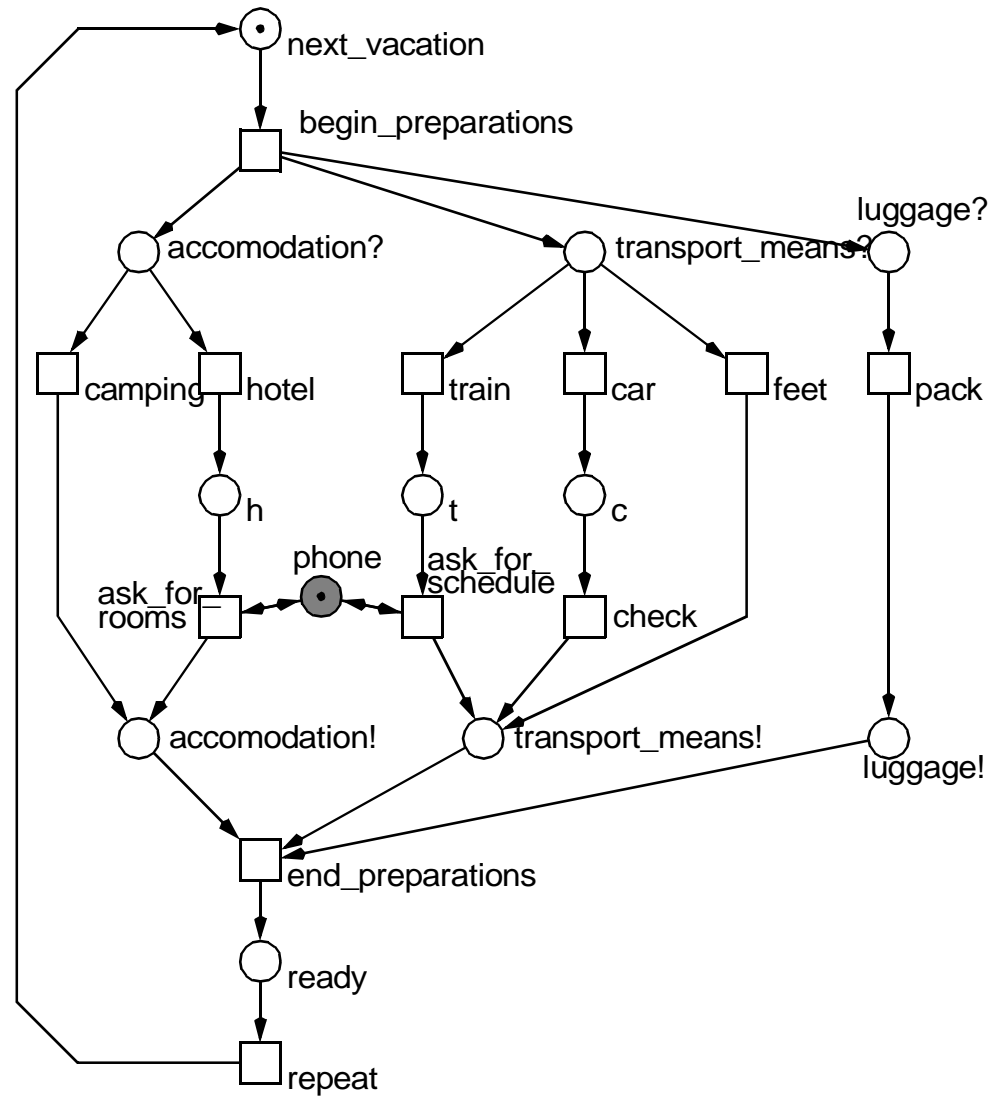
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PETRI NET, EXAMPLE 2



□ How many tokens can reside at most in a given place ?

-> (0, 1, *k*, ∞)

-> **BOUNDEDNESS**

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-> *never*

-> **UNREACHABLE -> SAFETY PROPERTIES**

-> *n-times*

-> **REPRODUCIBLE**

-> *always reachable again*

-> **REVERSIBLE (HOME STATE)**

-> **reversible initial state**

-> **REVERSIBILITY**

- ❑ **How many tokens can reside at most in a given place ?**
 - > $(0, 1, k, \infty)$
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- ❑ **Are there behaviourally invariant subnet structures ?**
 - > *token conservation*
 - > **P - INVARIANTS**
 - > *token distribution reproduction*
 - > **T - INVARIANTS**

- ❑ **... and many more -> temporal logics (CTL, LTL)**

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**GENERAL
BEHAVIOURAL
PROPERTIES**

**SPECIAL
BEHAVIOURAL
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Petri nets, typical analysis techniques

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MODEL ANIMATION (?)

Dynamic Analyses

- ❑ **reachability / occurrence graph,**
 - > *(labelled) state transition graph*
 - > *Kripke structure, CTMC, . . .*

- ❑ **nodes**
 - > *system states / markings*

- ❑ **arcs**
 - > *the (single) firing transition*
 - > *single step firing*

- ❑ **interleaving semantics**
 - > *(sequential) finite automaton*
 - > *concurrency == enumerating all interleaving sequences*

- ❑ **reachability graph construction - simple algorithm**

- ❑ **boundedness**
-> *finite graph*

- ❑ **reversibility**
-> *one Strongly Connected Component (SCC)*

- ❑ **liveness**
-> *every transition contained in all terminal SCC*

- ❑ **dead states**
-> *terminal nodes*

- ❑ **boundedness**
-> *finite graph*

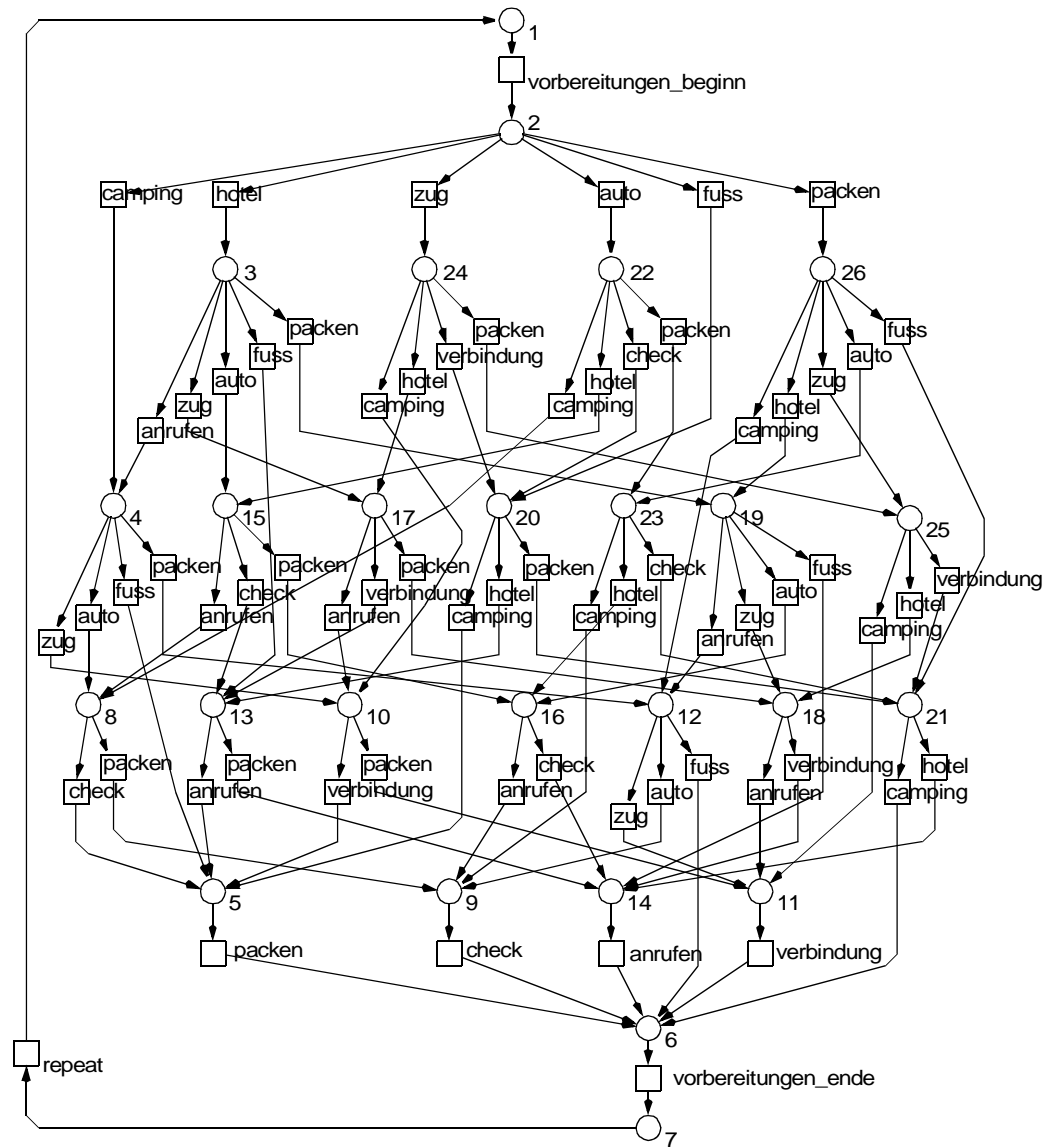
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-> reachability graphs tend to be huge <-

REACHABILITY GRAPH, Example 2



- ❑ **infinite for unbounded nets**

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- ❑ **worst-case for finite state spaces**

[Priese, Wimmel 2003]

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... cannot be bounded by a primitive recursive function ...

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- worst-case for finite state spaces [Priese, Wimmel 2003]
... cannot be bounded by a primitive recursive function ...
- proof -> Petri net computer for a function $f: \mathbb{N}_0^m \rightarrow \mathbb{N}_0$

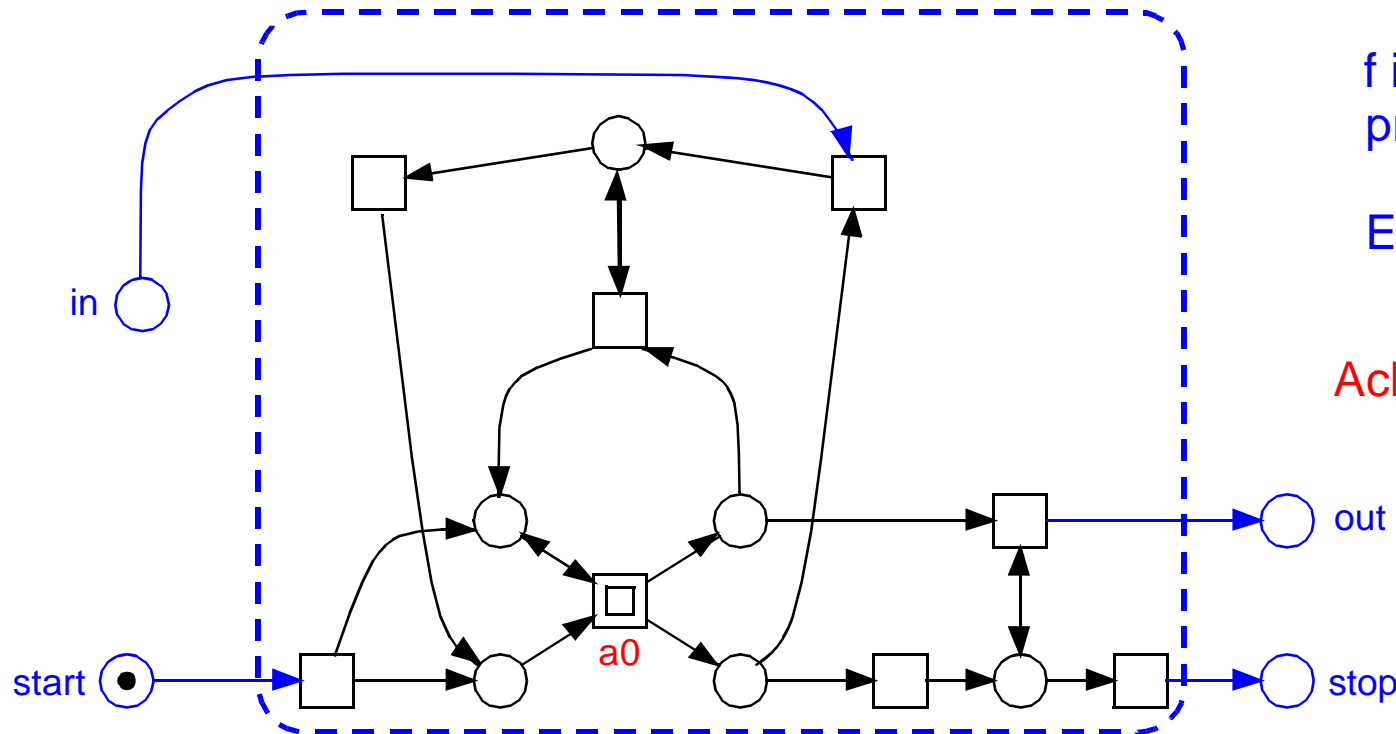


f is weakly
pn-computable:

EF(out = f(in) &
stop = 1)

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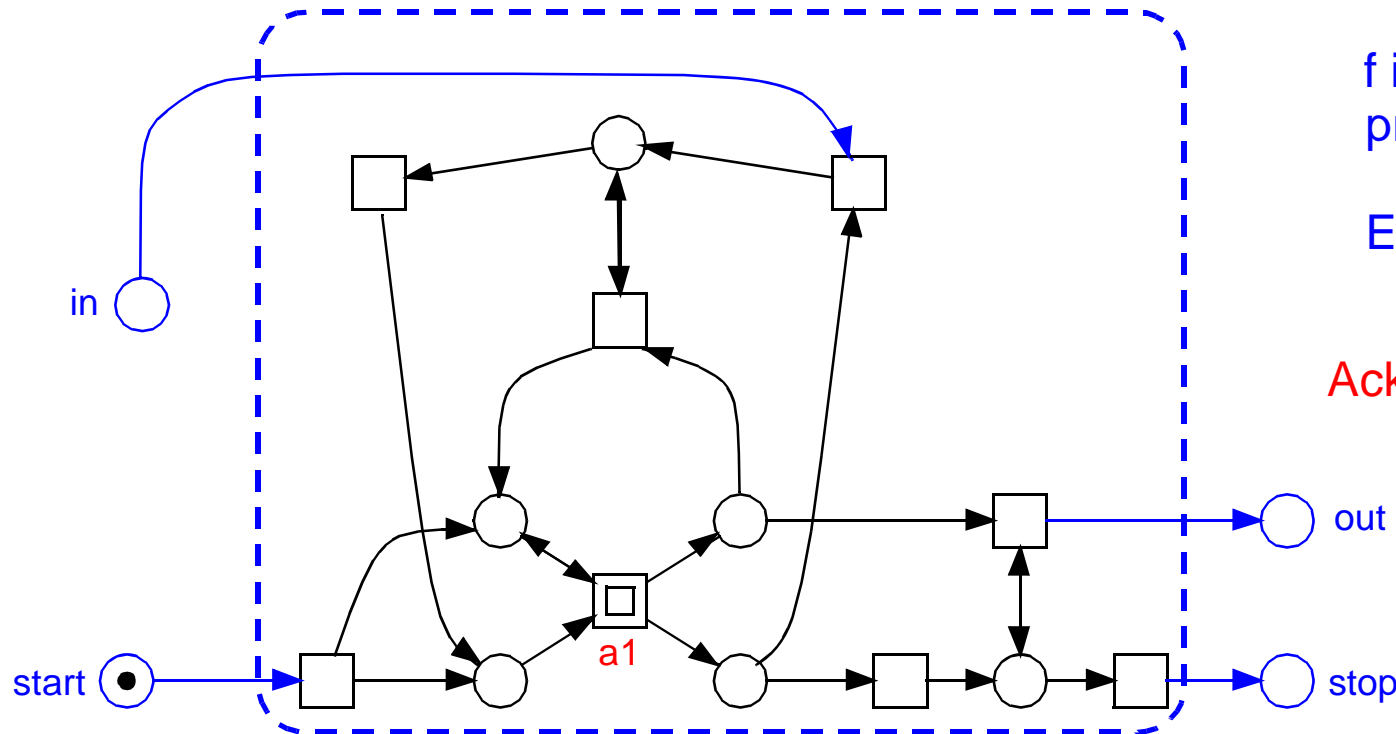
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Ackermann function a1

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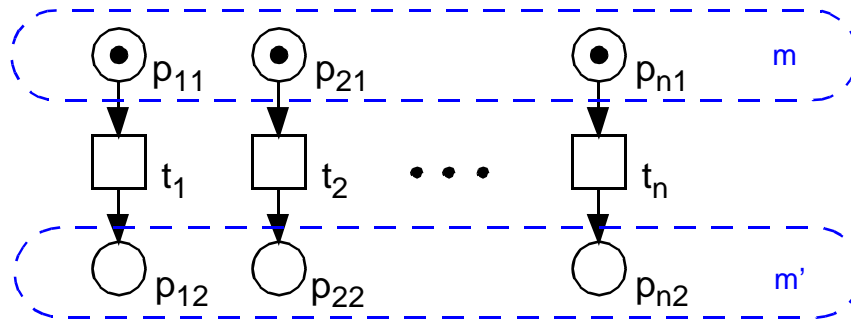
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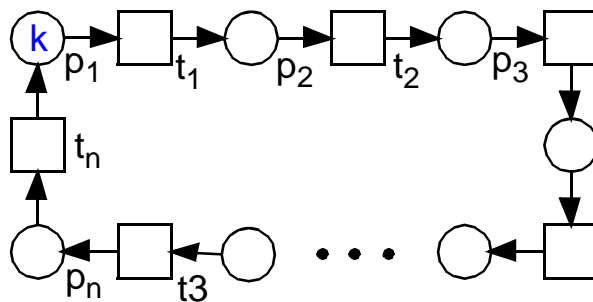
Ackermann function a2



$n!$ interleaving sequences

$m \rightarrow m'$

$2^n - 2$ intermediate states



$\frac{(n + k - 1)!}{(n - 1)! k!}$ states

(combination with repetition)

❑ **static analyses**

-> no state space construction

❑ **dynamic analyses**

-> total / partial state space construction

- **static analyses** **-> no state space construction**
 - > *structural properties (graph theory)*
 - > *P / T - invariants (linear algebra)*
 - > *state equation*

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- **dynamic analyses** **-> total / partial state space construction**
 - > *analysis of **general** behavioural system properties,
i.e. boundedness, liveness, reversibility*

 - > *model checking of **special** behavioural system properties,
e.g. reachability of a given (sub-) system state (with constraints),
reproducibility of a given (sub-) system state (with constraints)*

 - => expressed in temporal logics (CTL / LTL),
as very flexible & powerful query language**

Static Analyses

❑ **boundary nodes**

- > *input transitions* -> *not BND*
- > *input places* -> *not LIVE*
- > *LIVE & BND* -> *no boundary nodes*

❑ **conservative -> BND**

❑ ...

❑ **Deadlock-Trap Property (DTP)**

- > *no structural deadlock* -> *live*
- > *ORD & DTP* -> *no dead states*
- > *ORD & ES & DTP* -> *LIVE*
- > *ORD & EFC & DTP* <-> *LIVE*

- a representation of the net structure

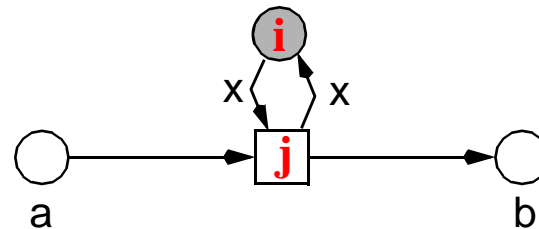
$$C =$$

| P \ T | t1 | ... | tj | ... | tm |
|-------|----|-----|--------------|-----|----|
| p1 | | | | | |
| pi | | | c_{ij} | | |
| ⋮ | | | Δt_j | | |
| pn | | | | | |

$$c_{ij} = (p_i, t_j) = F(t_j, p_i) - F(p_i, t_j) = \Delta t_j(p_i)$$

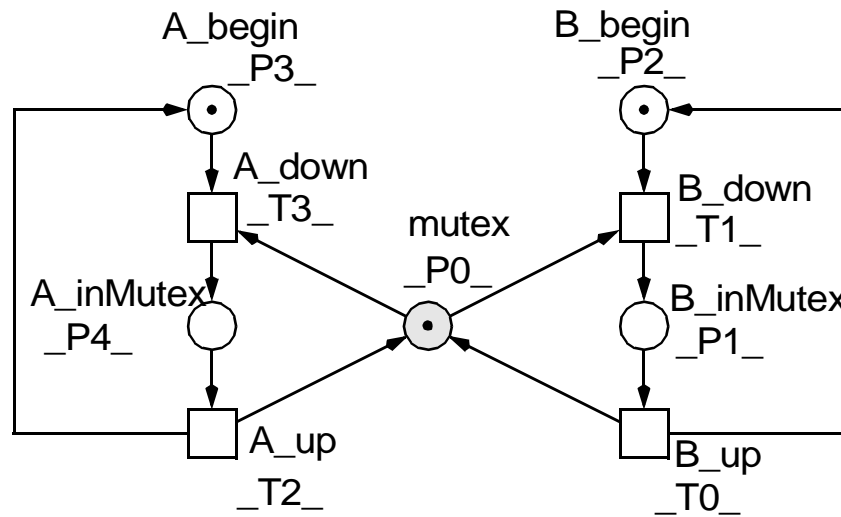
$$\Delta t_j = \Delta t_j^*$$

- matrix entry c_{ij} :
token change in place p_i by firing of transition t_j
- matrix column Δt_j :
vector describing the change of the whole marking by firing of t_j
- side-conditions are neglected

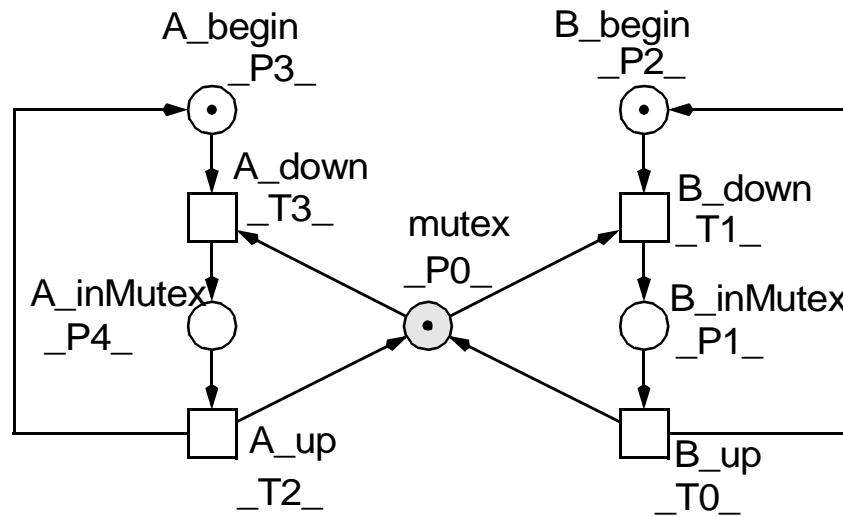


$$c_{ij} = 0$$

INCIDENCE MATRIX, EXAMPLE

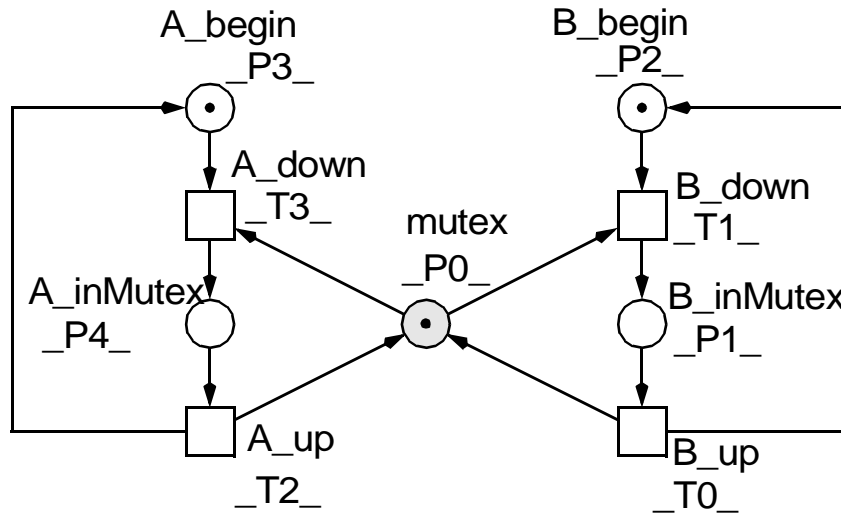


INCIDENCE MATRIX, EXAMPLE



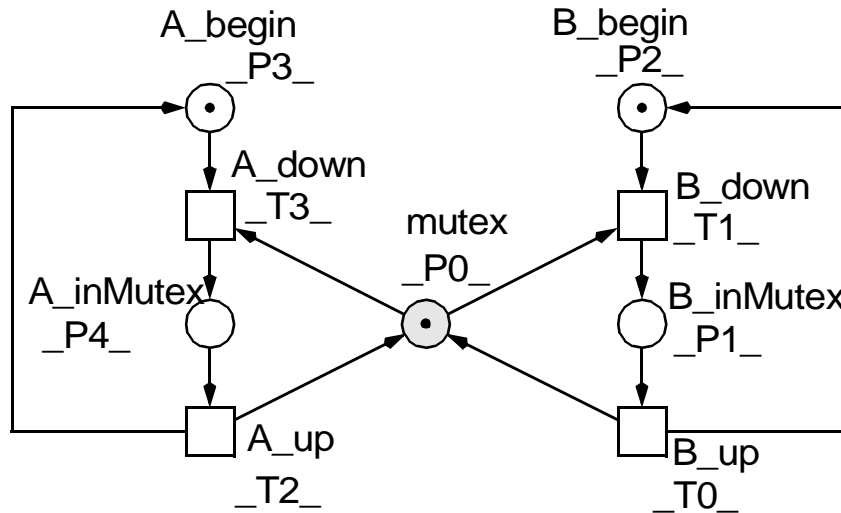
| P \ T | B_up | B_down | A_up | A_down |
|-----------|------|--------|------|--------|
| mutex | +1 | -1 | +1 | -1 |
| B_inMutex | -1 | +1 | 0 | 0 |
| B_begin | +1 | -1 | 0 | 0 |
| A_begin | 0 | 0 | +1 | -1 |
| A_inMutex | 0 | 0 | -1 | +1 |

INCIDENCE MATRIX, EXAMPLE



| P \ T | B_up | B_down | A_up | A_down | m0 |
|-----------|------|--------|------|--------|----|
| mutex | +1 | -1 | +1 | -1 | 1 |
| B_inMutex | -1 | +1 | 0 | 0 | 0 |
| B_begin | +1 | -1 | 0 | 0 | 1 |
| A_begin | 0 | 0 | +1 | -1 | 1 |
| A_inMutex | 0 | 0 | -1 | +1 | 0 |

INCIDENCE MATRIX, EXAMPLE



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|-----------|------|--------|------|--------|
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| B_inMutex | -1 | +1 | 0 | 0 |
| B_begin | +1 | -1 | 0 | 0 |
| A_begin | 0 | 0 | +1 | -1 |
| A_inMutex | 0 | 0 | -1 | +1 |

A_down

m0 →

1

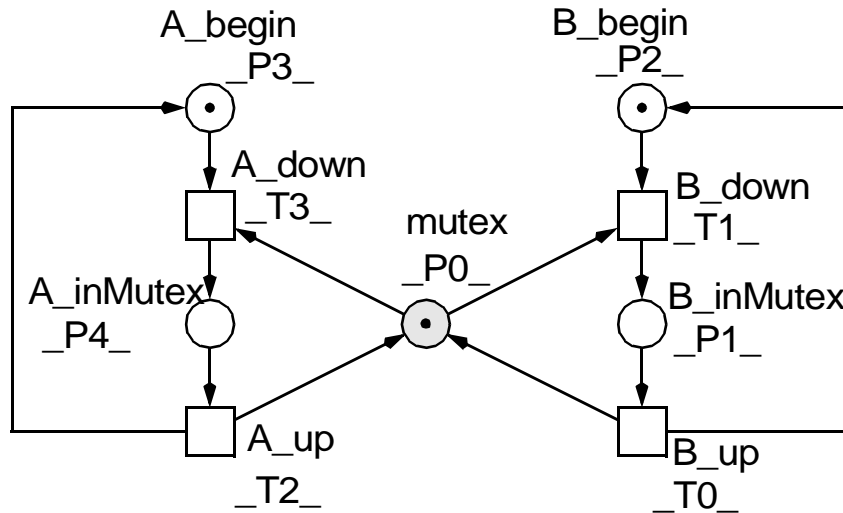
0

1

1

0

INCIDENCE MATRIX, EXAMPLE



| P \ T | B_up | B_down | A_up | A_down |
|-----------|------|--------|------|--------|
| mutex | +1 | -1 | +1 | -1 |
| B_inMutex | -1 | +1 | 0 | 0 |
| B_begin | +1 | -1 | 0 | 0 |
| A_begin | 0 | 0 | +1 | -1 |
| A_inMutex | 0 | 0 | -1 | +1 |

A_down

m0 →

1

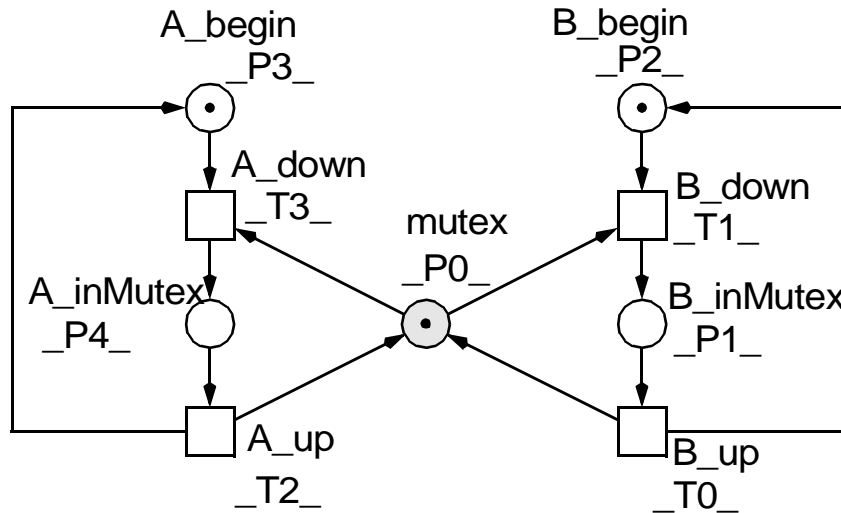
0

1

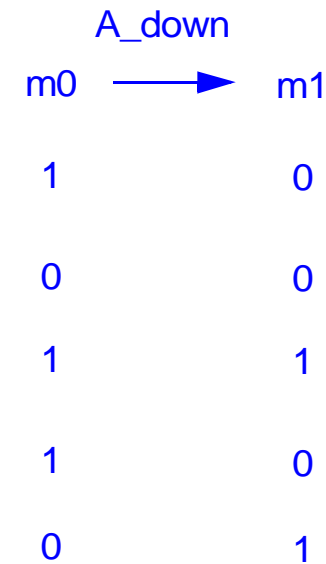
1

0

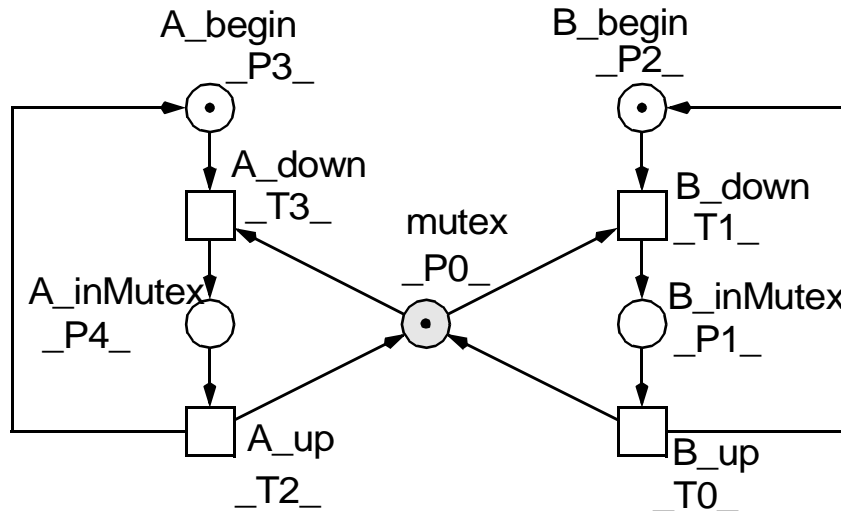
INCIDENCE MATRIX, EXAMPLE



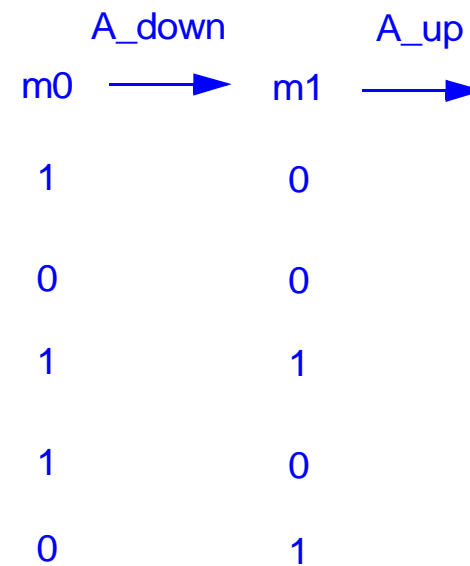
| P \ T | B_up | B_down | A_up | A_down |
|-----------|------|--------|------|--------|
| mutex | +1 | -1 | +1 | -1 |
| B_inMutex | -1 | +1 | 0 | 0 |
| B_begin | +1 | -1 | 0 | 0 |
| A_begin | 0 | 0 | +1 | -1 |
| A_inMutex | 0 | 0 | -1 | +1 |



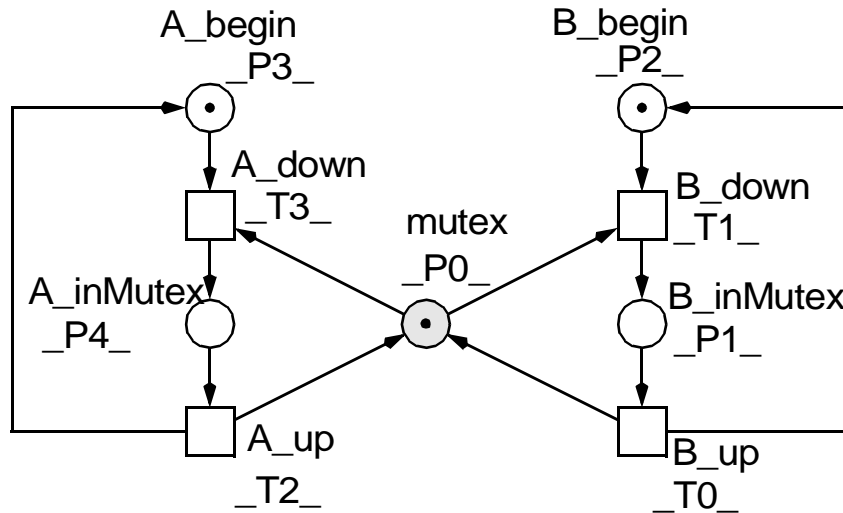
INCIDENCE MATRIX, EXAMPLE



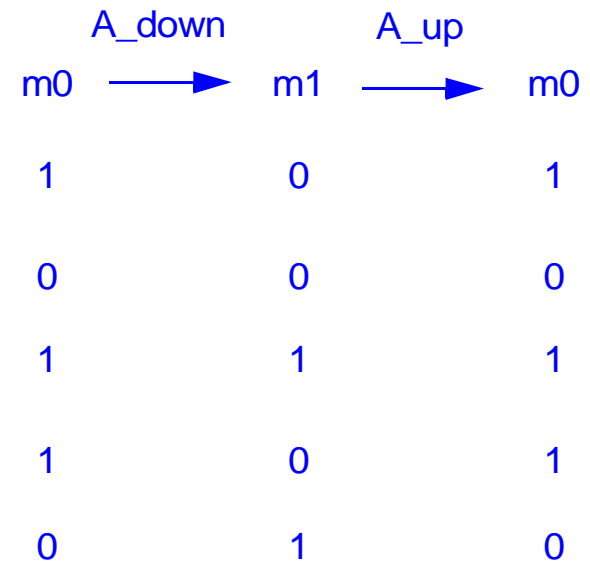
| P \ T | B_up | B_down | A_up | A_down |
|-----------|------|--------|------|--------|
| mutex | +1 | -1 | +1 | -1 |
| B_inMutex | -1 | +1 | 0 | 0 |
| B_begin | +1 | -1 | 0 | 0 |
| A_begin | 0 | 0 | +1 | -1 |
| A_inMutex | 0 | 0 | -1 | +1 |



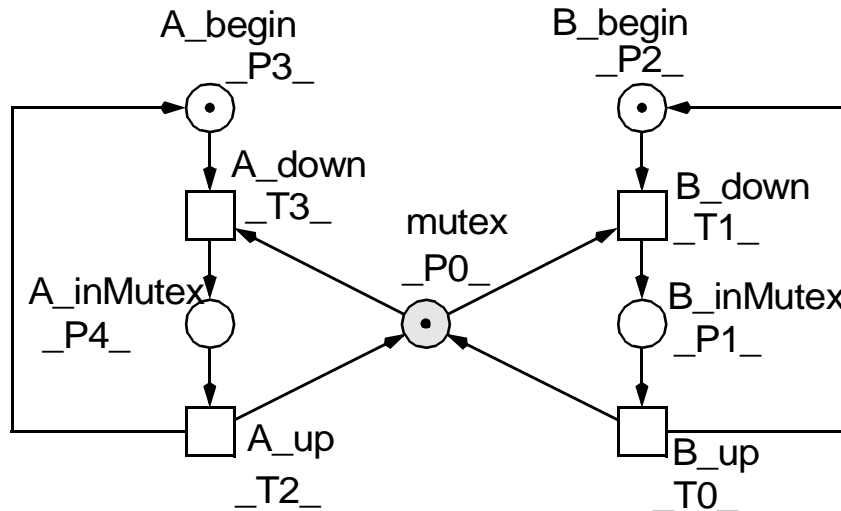
INCIDENCE MATRIX, EXAMPLE



| P \ T | B_up | B_down | A_up | A_down |
|-----------|------|--------|------|--------|
| mutex | +1 | -1 | +1 | -1 |
| B_inMutex | -1 | +1 | 0 | 0 |
| B_begin | +1 | -1 | 0 | 0 |
| A_begin | 0 | 0 | +1 | -1 |
| A_inMutex | 0 | 0 | -1 | +1 |



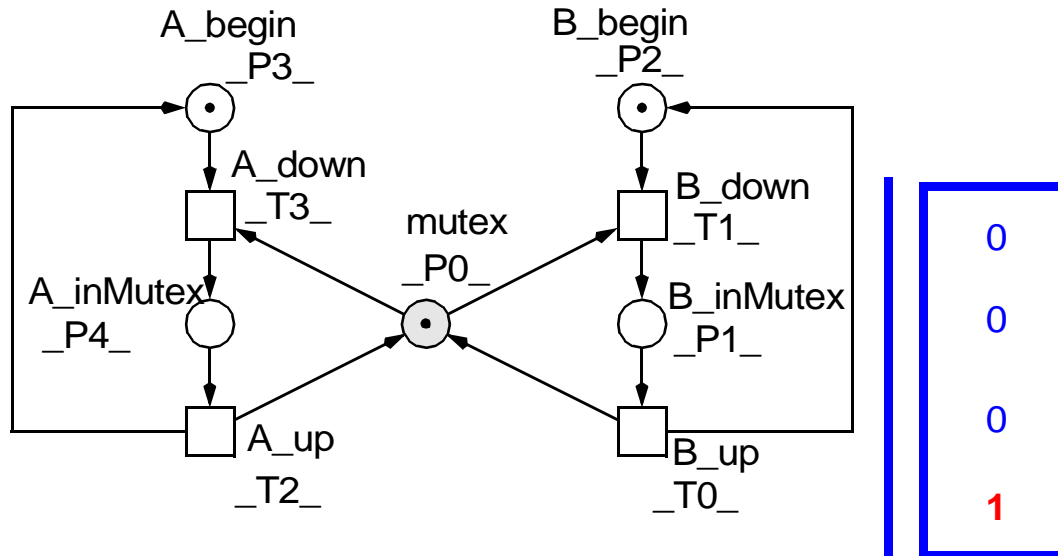
INCIDENCE MATRIX, EXAMPLE



0
 0
 0
 1

| P \ T | B_up | B_down | A_up | A_down |
|-----------|------|--------|------|--------|
| mutex | +1 | -1 | +1 | -1 |
| B_inMutex | -1 | +1 | 0 | 0 |
| B_begin | +1 | -1 | 0 | 0 |
| A_begin | 0 | 0 | +1 | -1 |
| A_inMutex | 0 | 0 | -1 | +1 |

INCIDENCE MATRIX, EXAMPLE

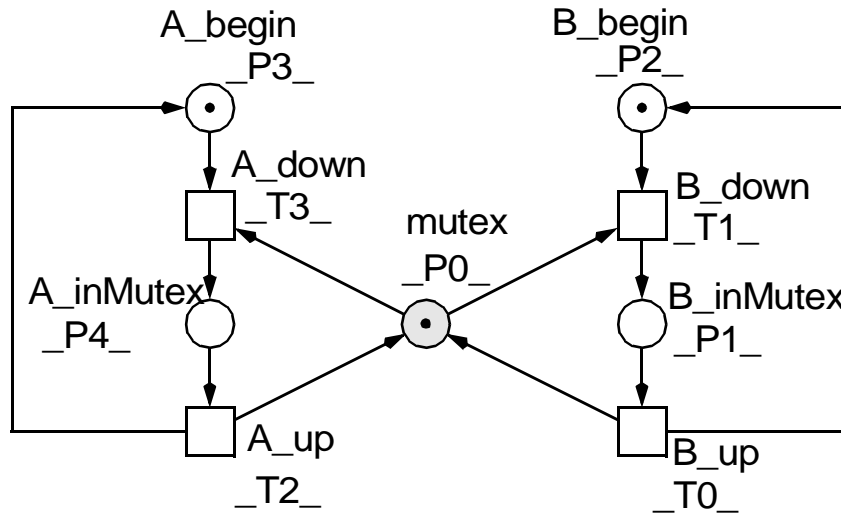


| P \ T | B_up | B_down | A_up | A_down |
|-----------|------|--------|------|--------|
| mutex | +1 | -1 | +1 | -1 |
| B_inMutex | -1 | +1 | 0 | 0 |
| B_begin | +1 | -1 | 0 | 0 |
| A_begin | 0 | 0 | +1 | -1 |
| A_inMutex | 0 | 0 | -1 | +1 |

0
0
0
1

-1
0
0
-1
1

INCIDENCE MATRIX, EXAMPLE



| |
|---|
| 0 |
| 0 |
| 0 |
| 1 |

| P \ T | B_up | B_down | A_up | A_down |
|-----------|------|--------|------|--------|
| mutex | +1 | -1 | +1 | -1 |
| B_inMutex | -1 | +1 | 0 | 0 |
| B_begin | +1 | -1 | 0 | 0 |
| A_begin | 0 | 0 | +1 | -1 |
| A_inMutex | 0 | 0 | -1 | +1 |

| |
|----|
| -1 |
| 0 |
| 0 |
| -1 |
| 1 |

| |
|---|
| 1 |
| 0 |
| 1 |
| 1 |
| 0 |

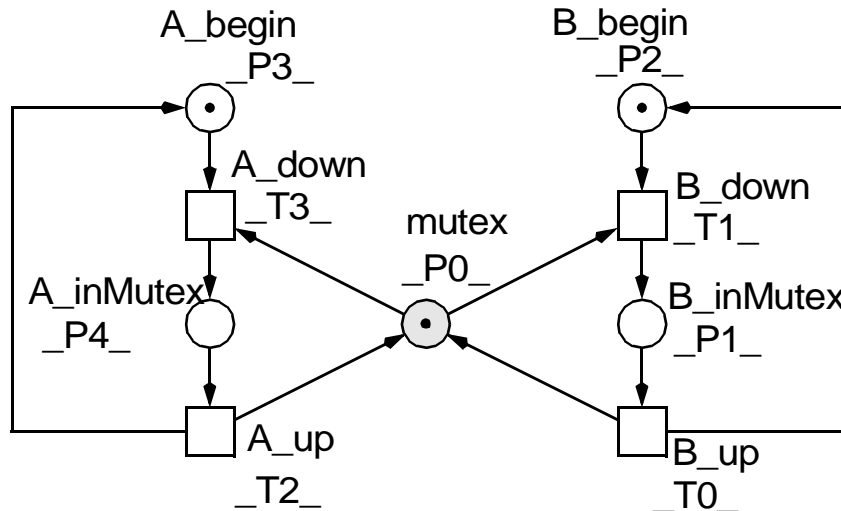
m0 → m1

+

=

| |
|---|
| 0 |
| 0 |
| 1 |
| 0 |
| 1 |

INCIDENCE MATRIX, EXAMPLE



$$\begin{bmatrix} 0 \\ 0 \\ 1 \\ 0 \end{bmatrix}$$

| P \ T | B_up | B_down | A_up | A_down |
|-----------|------|--------|------|--------|
| mutex | +1 | -1 | +1 | -1 |
| B_inMutex | -1 | +1 | 0 | 0 |
| B_begin | +1 | -1 | 0 | 0 |
| A_begin | 0 | 0 | +1 | -1 |
| A_inMutex | 0 | 0 | -1 | +1 |

$$\begin{bmatrix} 1 \\ 0 \\ 0 \\ 1 \\ -1 \end{bmatrix}$$

$$+$$

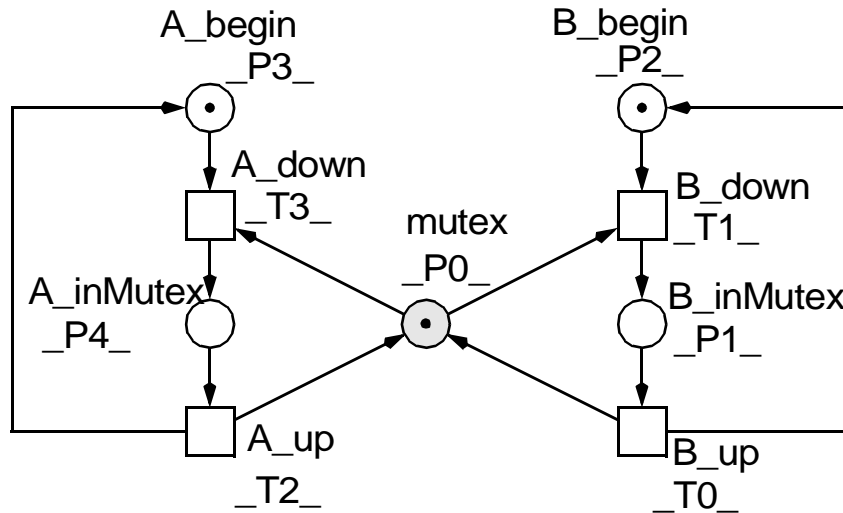
$$\begin{bmatrix} 0 \\ 0 \\ 1 \\ 0 \\ 1 \end{bmatrix}$$

=

$$\begin{bmatrix} 1 \\ 0 \\ 1 \\ 1 \\ 0 \end{bmatrix}$$

m1 → m0

INCIDENCE MATRIX, EXAMPLE



| |
|---|
| 0 |
| 0 |
| 1 |
| 1 |

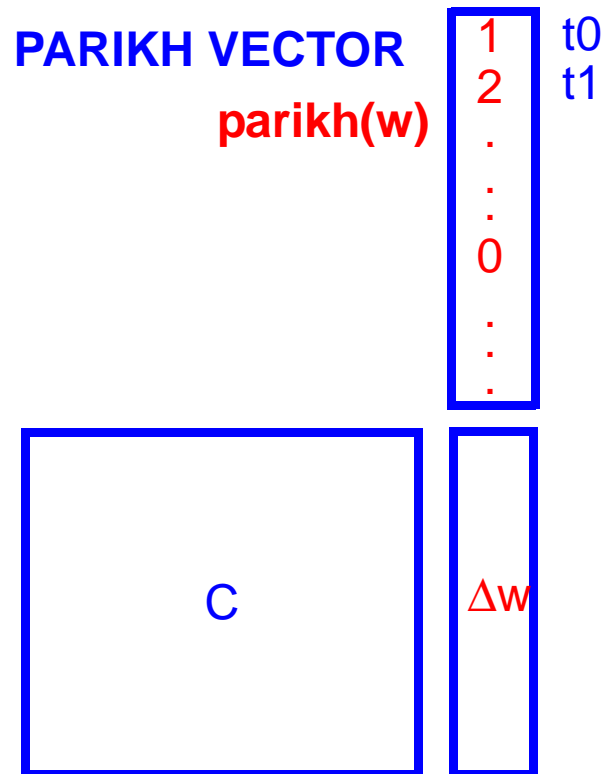
| P \ T | B_up | B_down | A_up | A_down |
|-----------|------|--------|------|--------|
| mutex | +1 | -1 | +1 | -1 |
| B_inMutex | -1 | +1 | 0 | 0 |
| B_begin | +1 | -1 | 0 | 0 |
| A_begin | 0 | 0 | +1 | -1 |
| A_inMutex | 0 | 0 | -1 | +1 |

| |
|---|
| 0 |
| 0 |
| 0 |
| 0 |
| 0 |

$$\begin{matrix} m0 \\ \rightarrow \\ m0 \end{matrix}$$

| |
|---|
| 1 |
| 0 |
| 1 |
| 1 |
| 0 |

$$= \begin{matrix} m0 \\ \rightarrow \\ m0 \end{matrix}$$



Let the word $w = t_1-t_0-t_1-\dots$
be a sequence of firing transitions.

The change of the marking

Δw

by firing that sequence
can be computed

by multiplying

the incidence matrix C

with the

Parikh vector $\text{parikh}(w)$

of that transition sequence.

PARIKH VECTOR
parikh(w)

$$\begin{bmatrix} 1 \\ 2 \\ \vdots \\ 0 \\ \vdots \\ \vdots \end{bmatrix} \begin{matrix} t_0 \\ t_1 \end{matrix}$$
$$C \begin{bmatrix} \Delta w \end{bmatrix} + m_0 = m$$

- linear programming problem

$$\begin{aligned} m &= m_0 + C x, \\ x &\geq 0 \end{aligned}$$

x - T-vector

- There exists an integer solution for every reachable marking m (the Parikh vector of the transition sequence going to m).
- **the integer solvability is a necessary condition for the reachability of a marking**
- **NON-REACHABILITY CHECK**
 - > *if there is no integer solution, then the marking is not reachable*
- **each reachable marking is a linear combination of**
 - > *initial marking*
 - > *columns of incidence matrix*

□ Lautenbach, 1973

□ P-invariants

-> integer solutions y of

$$yC = 0, y \neq 0, y \geq 0$$

-> multisets of places

□ minimal P-invariants

-> there is no P-invariant with a smaller support

-> sets of places

-> gcd of all entries is 1

□ any P-invariant is a non-negative linear combination of minimal ones

-> multiplication with a positive integer

-> addition

-> Division by gcd

$$ky = \sum_i a_i y_i$$

□ Covered by P-Invariants (CPI)

-> each place belongs to a P-invariant

-> CPI => BND (sufficient condition)

□ Lautenbach, 1973

-> Schuster, 1993

□ T-invariants

-> *multisets of transitions*

-> integer solutions x of

$$Cx = 0, x \neq 0, x \geq 0$$

-> Parikh vector

□ minimal T-invariants

-> there is no T-invariant with a smaller *support*

-> *sets of transitions*

-> gcd of all entries is 1

□ any T-invariant is a non-negative linear combination of minimal ones

-> multiplication with a positive integer

-> addition

-> Division by gcd

$$kx = \sum_i a_i x_i$$

□ Covered by T-Invariants (CTI)

-> each transition belongs to a T-invariant

-> *BND & LIVE => CTI (necessary condition)*

Petri nets, an application

-> error-correcting Petri nets