

# QUALITATIVE MODELLING AND ANALYSIS OF BIOCHEMICAL PATHWAYS WITH PETRI NETS

**Ina Koch**

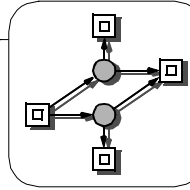
TFH BERLIN, BIOINFORMATICS

<http://www.tfh-berlin.de/bi/>

**Monika Heiner**

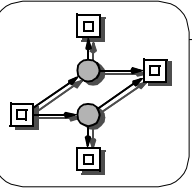
BTU COTTBUS, COMPUTER SCIENCE

<http://www-dssz.informatik.tu-cottbus.de>

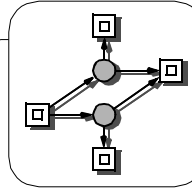


## OUTLINE

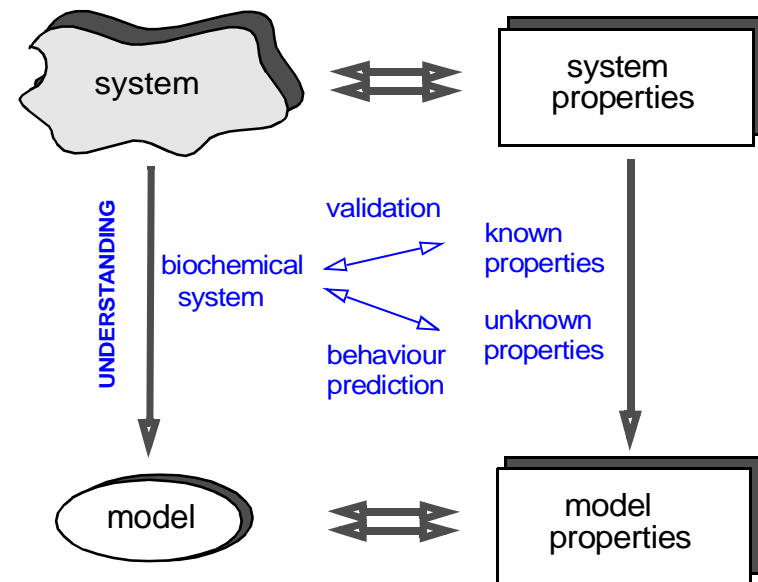
1. MOTIVATION
2. INTRODUCTION INTO QUALITATIVE MODELLING
3. INTRODUCTION INTO QUALITATIVE ANALYSIS
  - PROPERTIES
  - REACHABILITY GRAPH
  - TRANSITION / PLACE INVARIANTS
4. CASE STUDIES
  - APOPTOSIS IN MAMMALIAN CELLS
  - CENTRAL CARBON METABOLISM  
IN POTATO TUBERS
  - GLYCOLYSIS / PENTOSE PHOSPHATE  
PATHWAYS IN ERYTHROCYTES
5. SUMMARY, OUTLOOK
6. REFERENCES

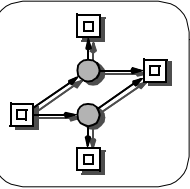


# 1. MOTIVATION



## MODEL- BASED SYSTEM ENGINEERING

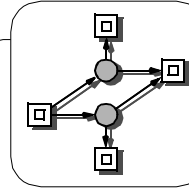




## BIOCHEMICAL SYSTEMS, EXAMPLES

- metabolic pathways / networks
  - > *stoichiometric relations known*
  - > *concentrations of metabolites often known*
- signal transduction pathways / networks
  - > *stoichiometric relations unknown*
  - > *read arcs / test arcs*
  - > *inhibitor arcs*
- gene regulatory networks
  - > *stoichiometric relations unknown*
  - > *mRNA concentrations often known*
  - > *protein concentrations are hard to be measured*
  - > *often a mixture of metabolic and signal transduction pathways*

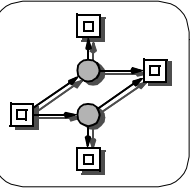
**=>> networks of elementary actions**



## BIOCHEMICAL SYSTEMS, SOME PROBLEMS

- network structure
  - > *very complex*
  - > *many reversible reactions*
  - > *dense, apparently unstructured*
  - > *hard to read*
  - > *tend to grow fast*
- knowledge
  - > *uncertain*
  - > *growing, changing*
  - > *distributed over independent data bases, papers, journals*
- representations
  - > *verbose descriptions*
  - > *contradictory and / or fuzzy statements*
  - > *diverse graphical representations*
  - > *various, mostly ambiguous*

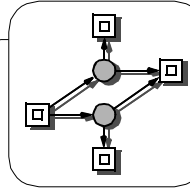
**=>> models of biochemical systems are typically full of assumptions**



## REPRESENTATIONS, OBJECTIVES

- ❑ readability  
-> *understanding*
- ❑ animation  
-> *experience*
- ❑ validation  
-> *consistency checks*
- ❑ analysis  
-> *behaviour prediction*  
(*qualitative / quantitative*)

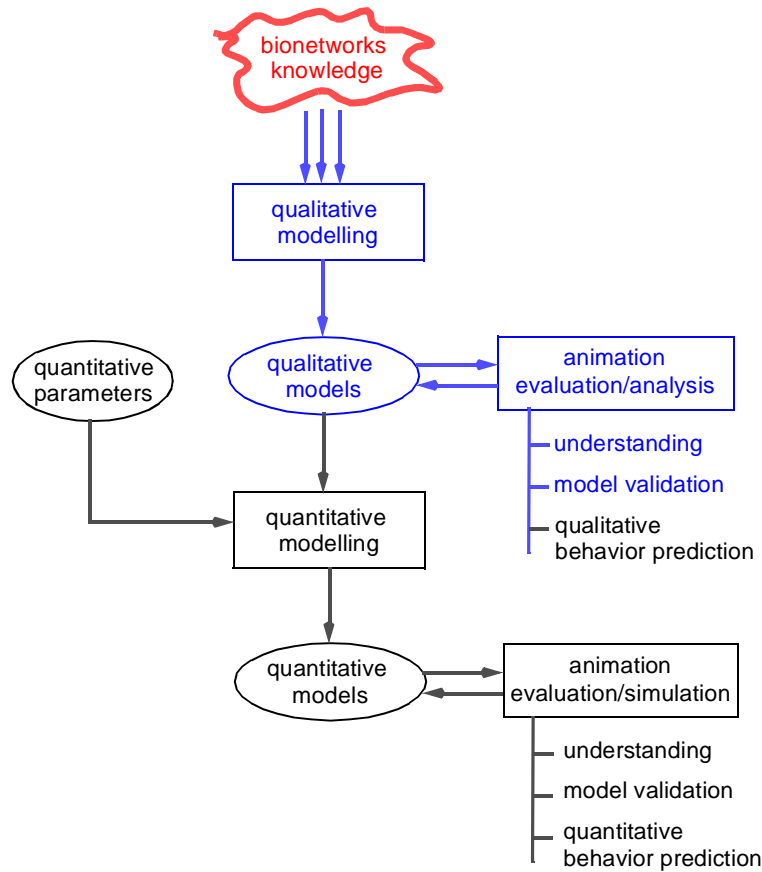
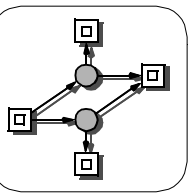
==>> *How many representations  
do we really need ?*



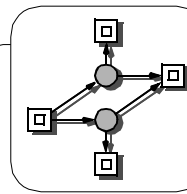
## WHY PETRI NETS ?

- ❑ a suitable intermediate representation for  
-> *different languages*  
-> *different stages of certainty*
- ❑ modelling power  
-> *partial order semantics*  
-> *applicable on any abstraction level*  
-> *specification of limited resources possible*
- ❑ analysis power  
-> *various complementary analysis methods*  
-> *reliable tools*
- ❑ *integration of qualitative and quantitative analyses*
- ❑ **BUT:**  
*modelling power <-> analysis power*

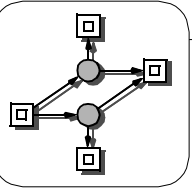
# FRAMEWORK



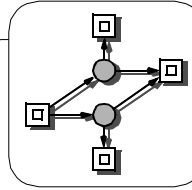
# PETRI NETS, A BIT OF HISTORY



- ❑ Carl Adam Petri, 1962,  
PhD University of Technology Darmstadt  
-> *basic ideas introduced*
- ❑ early 1970's  
-> *first papers contributing to Petri net theory*
- ❑ Petri, 1976  
-> *application to chemical networks mentioned*
- ❑ early 1980's  
-> *first monographs appear*
- ❑ Reddy, 1993  
-> *first paper on bio application*
- ❑ late 1990's  
-> *increasing interest in applying Petri nets for modelling and analysis of bio networks*



## 2. INTRODUCTION INTO QUALITATIVE MODELLING



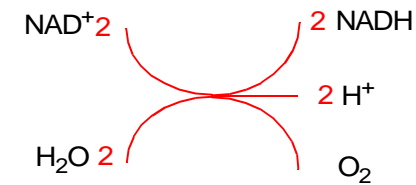
### BIOCHEMICAL SYSTEMS, BASIC COMPONENT

- chemical reactions -> atomic actions

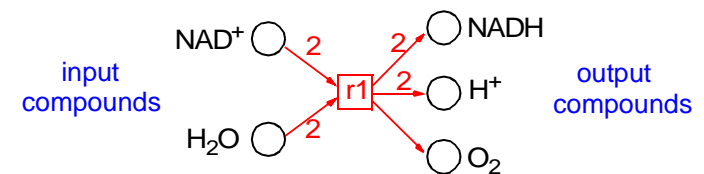
EX. LIGHT-INDUCED PHOSPHORYLATION

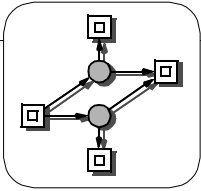


↓  
hyperarcs



↓  
Petri nets





## PETRI NETS, STRUCTURE

- two types of nodes

-> *places*

"passive elements", conditions, local states, chemical compounds

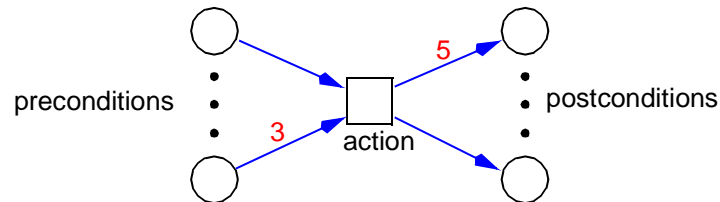


-> *transitions*

"active elements", events, actions, chemical reactions



- arcs



-> *directed*

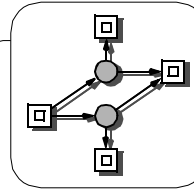
-> *never arcs between nodes of the same type*

-> *for any node, arbitrary number of pre-nodes and post-nodes*

- arc inscriptions

-> *arc weight / multiplicity*

-> *amount of units of the substances involved in the basic (re-) action*



## BIO PETRI NETS, PLACES

- involved substances / chem. compounds / complexes

- primary compounds

-> *metabolites*

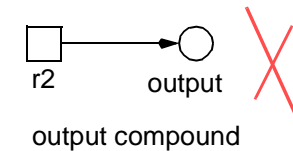
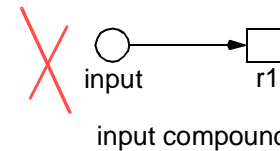
-> *complexes*

-> ...

- input / output compounds

-> *special primary compounds*

-> *boundary places*

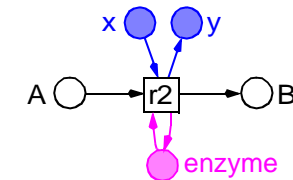


- auxiliary compounds

-> *side conditions for reactions*

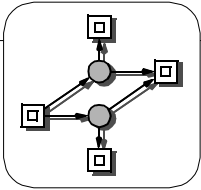
-> *ubiquitous -> fusion nodes*

e. g. *electron carrier, phosphate carrier;*



- catalysing compounds

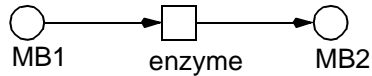
-> *enzymes, if any*



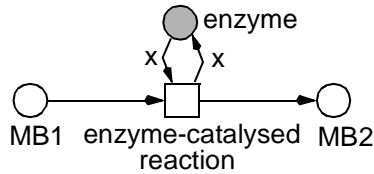
## BIO PETRI NETS, TRANSITIONS

- spontaneous reactions
- enzyme-catalysed reactions  
-> *two ways of modelling*

without enzyme concentration

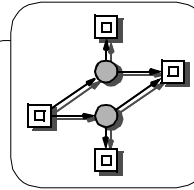


with enzyme concentration x



x - amount of enzyme units required by the reaction

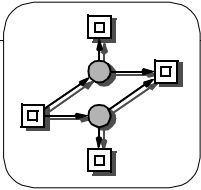
- transport steps, if any  
-> *inhomogeneous substance distribution*



## PETRI NETS, SYSTEM STATE

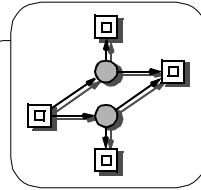
- tokens  
-> *moving objects,*  
*e. g. units of substances (e. g. Mole), ...*
- condition is not fulfilled
- condition is (one times) fulfilled
- Ⓝ condition is n times fulfilled
- > *token amount -*  
*amount of available units of a given compound*
- marking  
-> *How many tokens are on each place?*
- > *system state*  
-> *substance distribution*
- > *initial marking*  
-> *initial substance distribution*



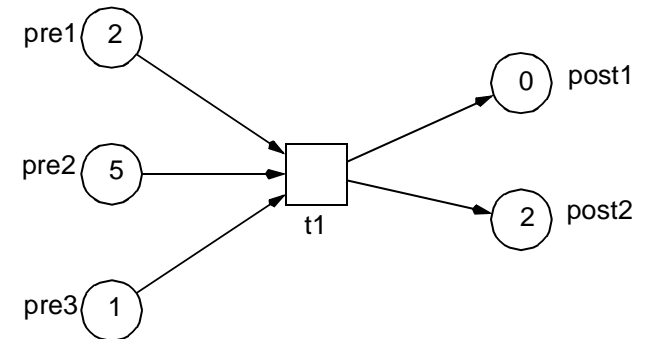


## PETRI NETS, BEHAVIOUR

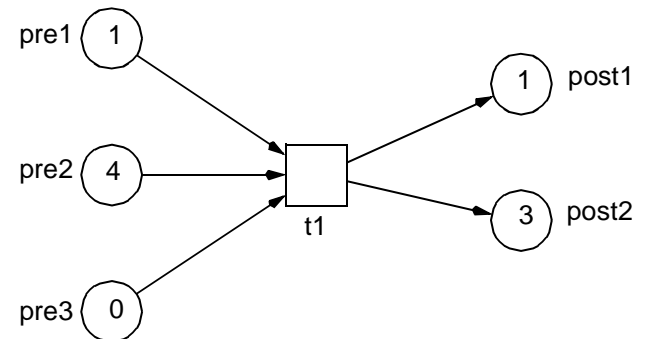
- ❑ flow of tokens  
-> *defined by firing rule*
- ❑ an action **may** happen (fire), if  
-> *all preconditions are fulfilled (corresponding to the arc weights);*
- ❑ **if** an action happens (fires), **then**  
-> *tokens are removed from all preconditions (corresponding to the arc weights), and*  
-> *tokens are added to all postconditions (corresponding to the arc weights);*
- ❑ an action happens (firing of a transition)  
-> **atomic**  
-> **time-less**



## FIRING RULE, EXAMPLE 1

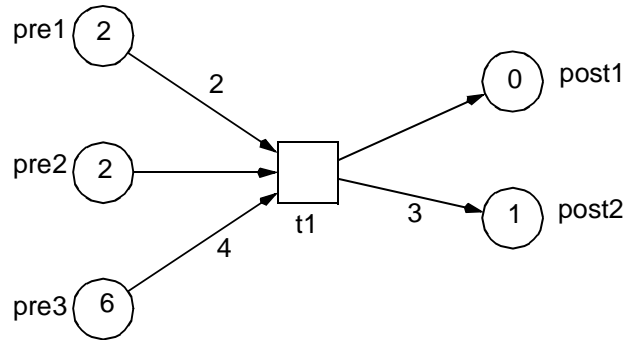


t1 fires

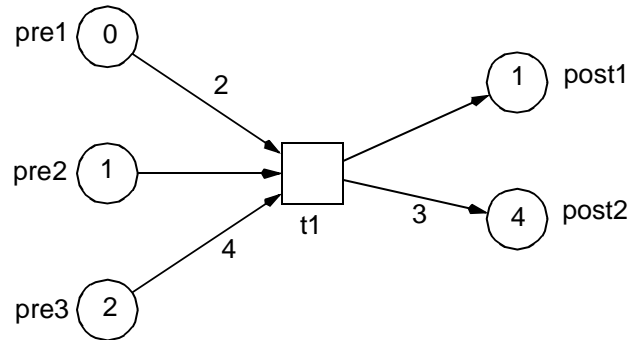


fire1.spped

### FIRING RULE, EXAMPLE 2



t1 fires



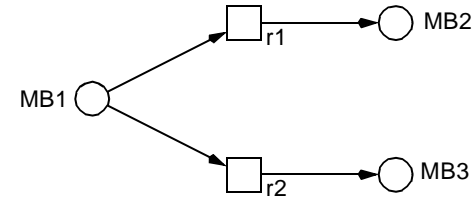
fire2.spped

### TYPICAL BASIC STRUCTURES 1

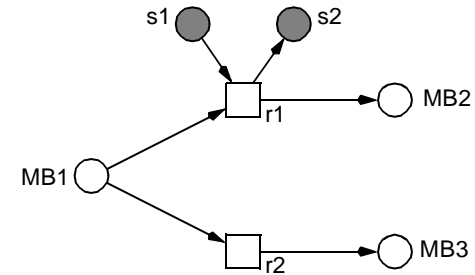
#### CHAIN OF REACTIONS



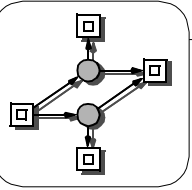
#### (FREE-CHOICE) BRANCHING / CONFLICT



#### BRANCHING WITH SIDE CONDITION

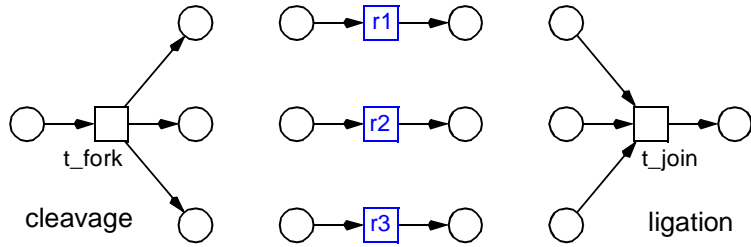


basicStructures1.spped



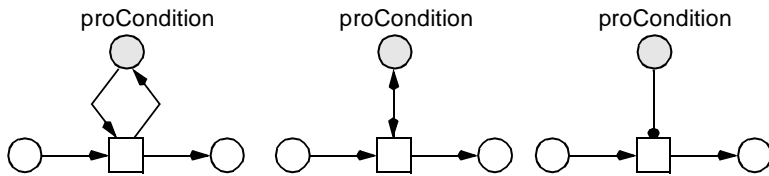
### TYPICAL BASIC STRUCTURES 2

#### CONCURRENCY



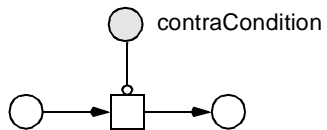
r1, r2, r3 are concurrent = independent

#### READ ARCS / TEST ARCS

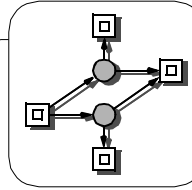


#### INHIBITOR ARCS

**BUT: CAUTION !**

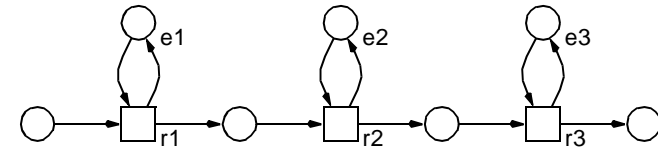


basicStructures2.spped

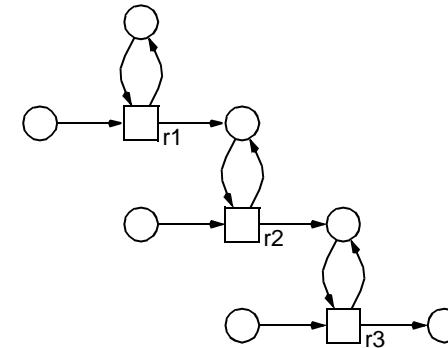


### TYPICAL BASIC STRUCTURES 3

#### METABOLIC PATHWAY

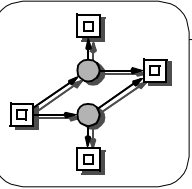


#### SIGNAL TRANSDUCTION CASCADE

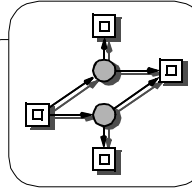
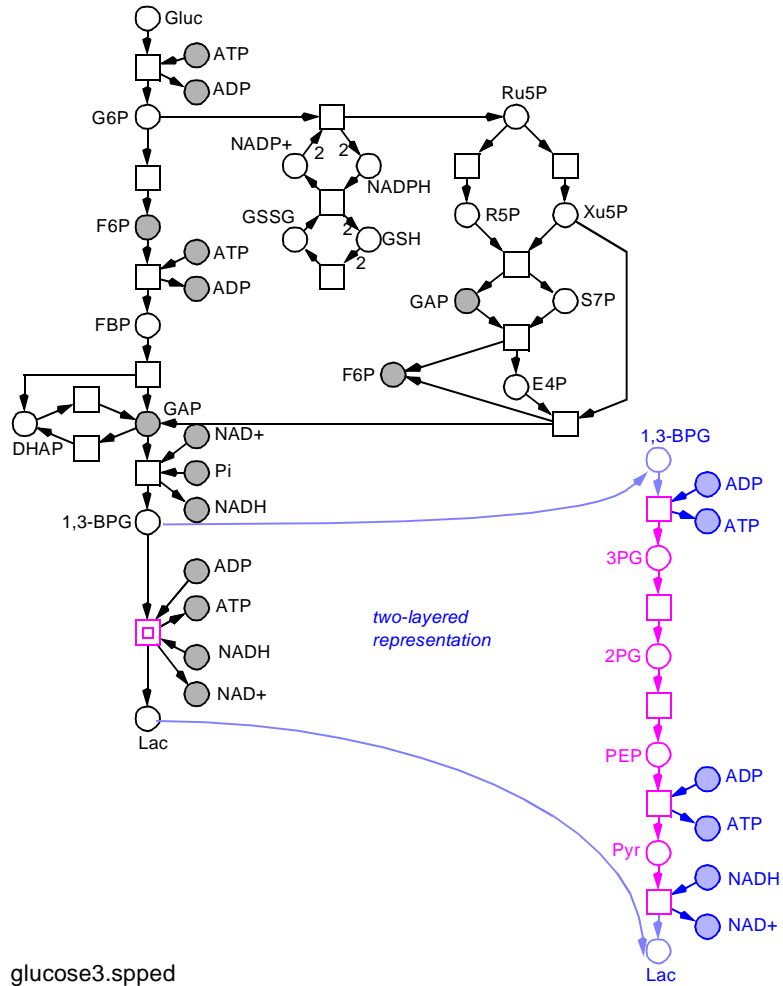


basicStructures3.spped





### EXAMPLE 1 AS PETRI NET, VERSION 3



### EXTENSIONS, SUMMARY

#### SYNTACTIC SUGAR

- logical / fusion nodes

-> connectors



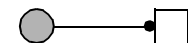
- hierarchical nodes

-> different levels of abstraction



- read arcs

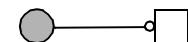
-> pro-conditions

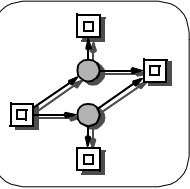


#### MODELLING POWER

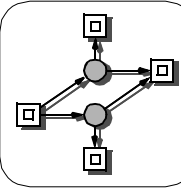
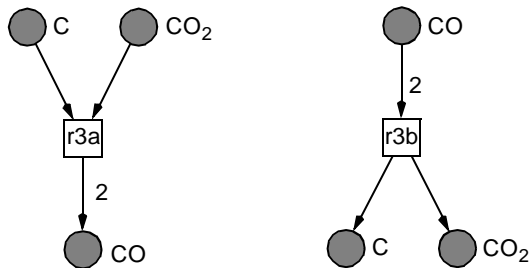
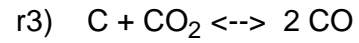
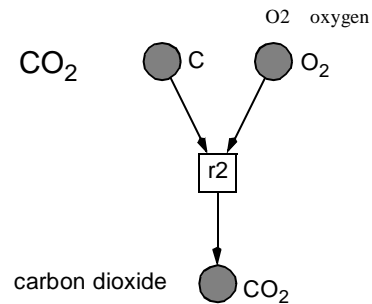
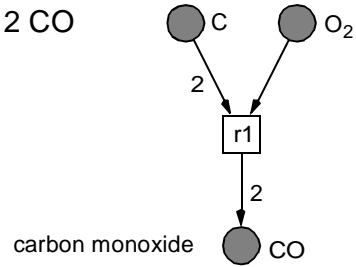
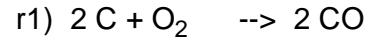
- inhibitor arcs

-> contra-conditions



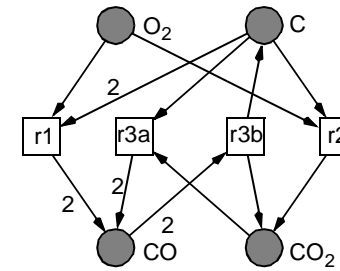


### EXAMPLE 2, CARBON OXIDATION, BASIC REACTIONS

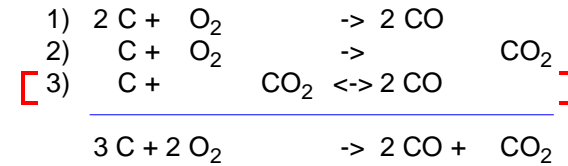


### EXAMPLE 2, COMPOSITION

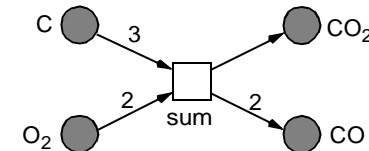
#### BASIC MODEL

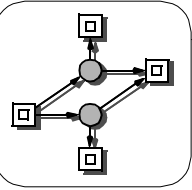


#### SYSTEM'S TOTAL EQUATION



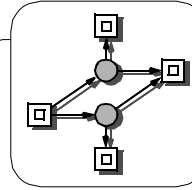
#### MODEL OF THE SYSTEM'S TOTAL EQUATION





## NETWORKS NEED ENVIRONMENT BEHAVIOUR

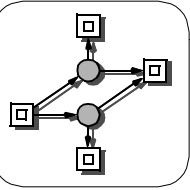
- ❑ to animate the model
  - > *infinite substance flow*
  - > *deeper insights*
- ❑ to validate the model
  - > *consistency criteria*
- ❑ steady flow
  - > *input compounds*
  - > *output compounds*
- ❑ auxiliary compounds
  - > *as much as necessary*
- ❑ minimal assumptions



## ENVIRONMENT BEHAVIOUR, THREE STYLES

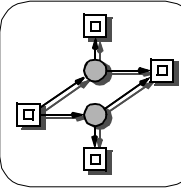
- ❑ **style 1**                   -> **validation criterion 1**
  - > *weak* assumptions
  - > *infinite flow into/out the network*
- ❑ **style 2**                   -> **validation criterion 2**
  - > *firm* assumptions
  - > *infinite many primary compounds*
  - > *finite, but sufficient reservoir of auxiliary compounds*
- ❑ **style 3**                   -> **validation criterion 3**
  - > *strong* assumptions
  - > *finite, but sufficient reservoir of auxiliary compounds*
  - > *quantitative relations of input / output compounds*
  - > *finite reservoir of primary compounds*

INCREASING STRENGTH

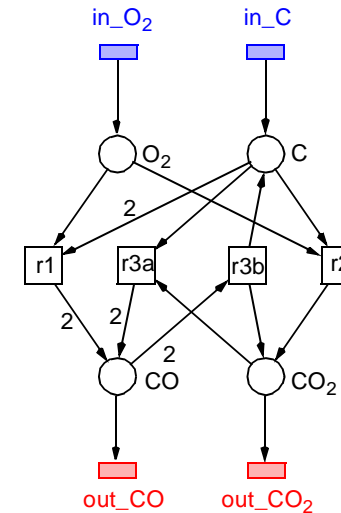


### ENVIRONMENT BEHAVIOUR, STYLE 1

- ❑ no assumptions about quantitative relations of input / output compounds
- ❑ input *compounds*  
-> *generating pre-transitions*
- ❑ output *compounds*  
-> *consuming post-transitions*
- ❑ auxiliary *compounds*  
-> *generating pre-transitions*  
-> *consuming post-transitions*  
-> *infinite reservoir*
- ❑ no boundary places, but boundary transitions

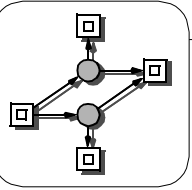


### CARBON OXIDATION, SYSTEM MODEL, STYLE 1



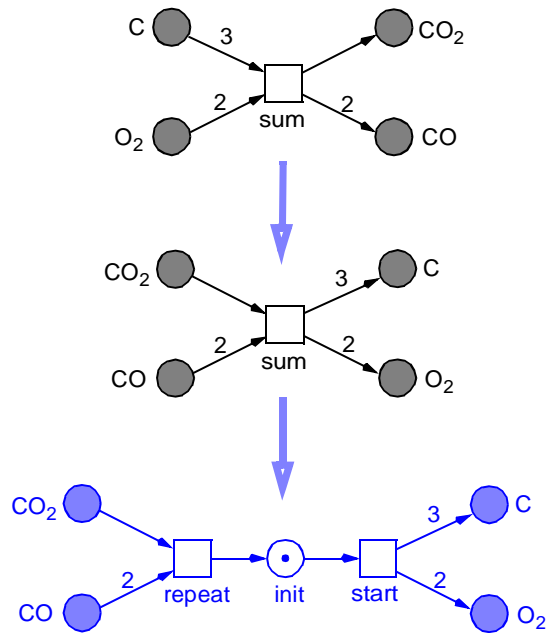
carbon1.sped



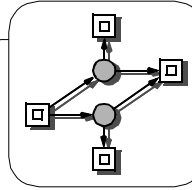


### ENVIRONMENT BEHAVIOUR, STYLE 3

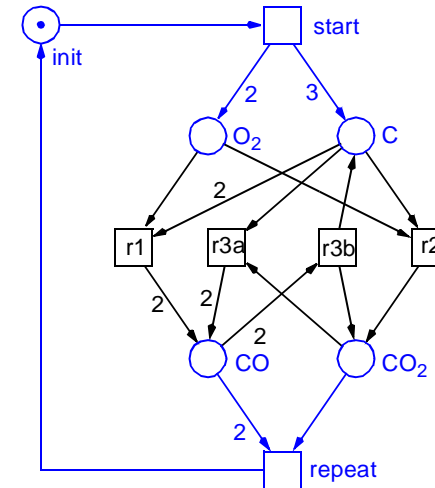
- strong assumptions about quantitative relations of input / output compounds
- 'inverse' total equation



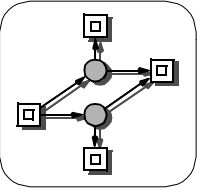
- there are no boundary nodes



### CARBON OXIDATION, SYSTEM MODEL, STYLE 3

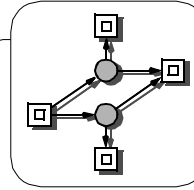


carbon2.spped



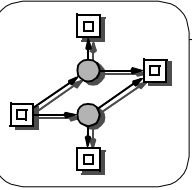
## ENVIRONMENT BEHAVIOUR, STYLE 2

- ❑ mixture of style 1 and style 3
- ❑ no assumptions about quantitative relations of input / output compounds
- ❑ input *compounds*  
-> *generating pre-transitions*
- ❑ output *compounds*  
-> *consuming post-transitions*
- ❑ auxiliary *compounds*  
-> *finite, but sufficient reservoir*  
-> *no boundary pre- / post-transitions*
- ❑ boundary transitions only for input / output compounds

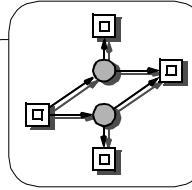


## BIOCHEMICAL PETRI NETS, SUMMARY

- ❑ biochemical networks  
-> *networks of (abstract) chemical reactions*
- ❑ **biochemically interpreted Petri net**  
  
-> *partial order sequences of chemical reactions*  
- *transforming input into output compounds*  
- *respecting the given stoichiometric relations*  
  
-> *set of all pathways*  
*from the input to the output compounds*  
*respecting the stoichiometric relations*
- ❑ pathway  
-> *self-contained partial order sequence*  
*of elementary (re-) actions*
- ❑ basic assumption  
-> ***steady state behaviour***



### 3. INTRODUCTION INTO QUALITATIVE ANALYSIS



### QUALITATIVE PROPERTIES



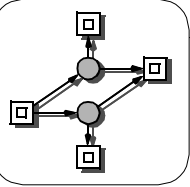
#### behavioural properties

- general semantic properties
  - boundedness*
  - liveness*
  - reversibility*
- special semantic properties
  - safety properties*
  - progress properties*



#### structural properties

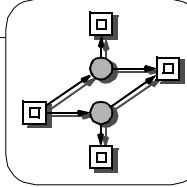
- especially valuable:
  - local(ly decidable) structural properties
- certain combinations of structural properties allow conclusions on behavioural properties



## PETRI NET PROPERTIES, OVERVIEW / INA

### 1. SIMPLE STRUCTURAL PROPERTIES

- ORD** ordinary (*1-multiplicity of all arcs*)
- HOM** homogeneous (*all output arcs of a given place have the same multiplicity*)
- NBM** non-blocking multiplicity (*for each place applies: MIN multiplicity of input arcs  $\geq$  MAX multiplicity of output arcs*)
- PUR** pure (*no side conditions*)
- CSV** conservative (*any firing preserves token amount*)
- SCF** static conflict free
- CON** connected
- SC** strongly connected
- Ft0** there is a transition without pre-place
- tF0** there is a transition without post-place
- Fp0** there is a place without pre-transition
- pF0** there is a place without post-transition
- MG** marked graph (*synchronization graph*)
- SM** state machine
- FC** free choice net
- EFC** extended free choice net
- ES** extended simple net



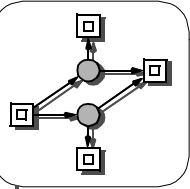
## PETRI NET PROPERTIES, OVERVIEW / INA

### 2. MORE EXPENSIVE STRUCTURAL PROPERTIES

- DTP** deadlock trap property
- SMC** state machine coverable (*covered with SM components*)
- SMD** state machine decomposable (*covered with SCSM components*)
- SMA** state machine allocatable
- CPI** covered with place invariants
- CTI** covered with transition invariants
- SB** structurally bounded

### 3. BEHAVIOURAL PROPERTIES

- B** bounded
- REV** reversible (*the initial state  $m_0$  can be reached again from all reachable states: home state*)
- DSt** dead states (*a state where no transition is enabled*)
- BSt** bad states (*a state where a fact is enabled*)
- DTr** dead transitions (*at the initial state*)
- DCF** dynamically conflict free
- L** live
- LV** live, excepted transitions dead at the initial marking
- L&S** live & safe (*1-bounded*)



## BEHAVIOURAL NET PROPERTIES, OVERVIEW

### MARKABILITY of places

- markable (*place liveness*)
- **k-bounded** (*1-bounded / safe*)
- unbounded

### LIVENESS of transitions

- zero times firing ( *$m_0$ -dead*)
- finite times firing (*dead, non-live*)
- **infinite times firing, probably** (*live*)
- infinite times firing, definitely (*livelock free*)

general semantic properties

### REACHABILITY of states

- dead states
- reproducibility
- **reversibility** ( *$m_0$  - home state*)
- bad states (*facts*)
- user-specified states

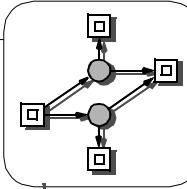
### NET INVARIANTS

- **transition invariants**
- **place invariants**

special semantic properties

### temporal relationship of logic formulae

- safety properties
- progress properties



## QUALITATIVE ANALYSIS METHODS, OVERVIEW

### NET REDUCTION

### STRUCTURAL PROPERTIES

### LINEAR PROGRAMMING

- place / transition invariants
- state equation
- trap equation

static analysis

### STATE SPACE ANALYSIS

- **(complete) reachability graph**

### compressed state spaces

- BDDs, NDDs, ..., xDDs
- Kronecker products

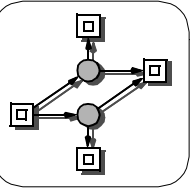
dynamic analysis

### reduced state spaces

- coverability graph
- symmetry
- stubborn sets

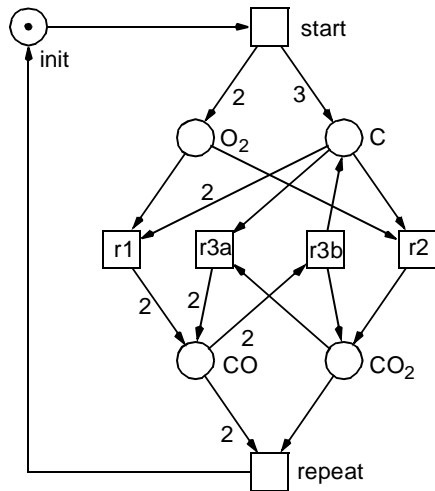
(model checking)

### branching process

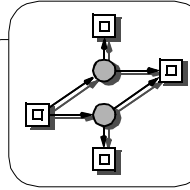


## REACHABILITY GRAPH (RG)

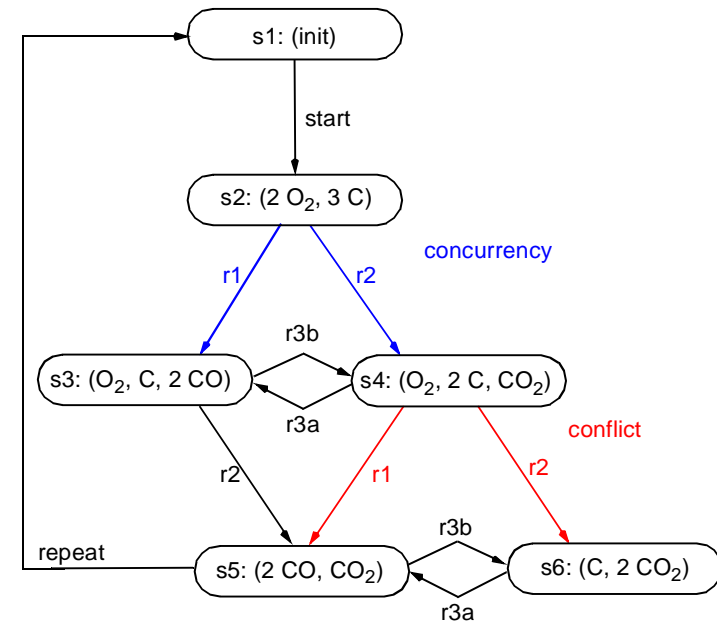
- nodes - system states
- arcs - the (single) firing transition
- example - carbon oxidation, environment style 3



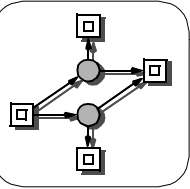
carbon2.ssped



## RG (CARBON OXIDATION)

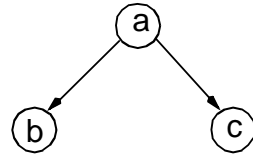


-> interleaving description  
of the whole system behaviour

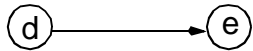


### STRONGLY CONNECTED GRAPH

- basic graph properties  
-> applies also for general (monochromatic) graphs

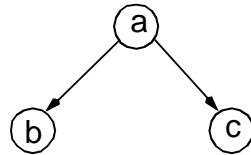


- needs directed graphs  
undirected graphs: connected = strongly connected



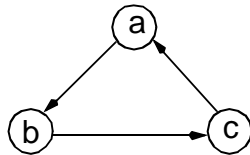
not connected

- for each pair of nodes x, y holds:  
there exists a path from x to y  
-> path(x, y);



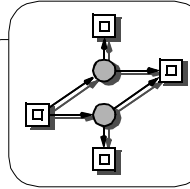
connected

- path(x, y):  
sequence of arcs starting at x and ending at y;



strongly connected

- general importance  
ex: system of one-way streets;  
question: is every place (intersection) from any place reachable?



### EXAMPLE: RG AND THREE BASIC PN PROPERTIES

- no concurrency  
->  $rg(pn) == pn$

- rg - finite  
-> bounded pn

- rg - not sc  
-> pn not reversible

- no dead states, but liveness?

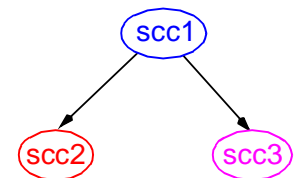
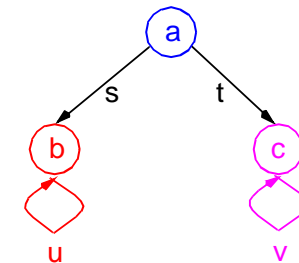
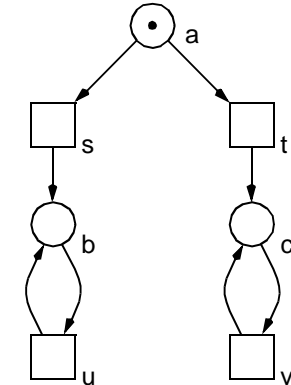
- condensed rg  
node - sc component (scc)

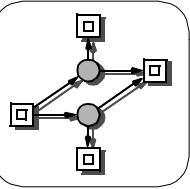
scc: maximal set of sc nodes;

a terminal scc  
-> possible terminal system behaviour

-> must contain all transitions in a live pn

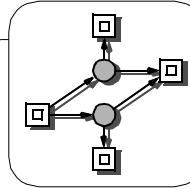
- not all terminal scc contain all transitions  
-> the pn is not live



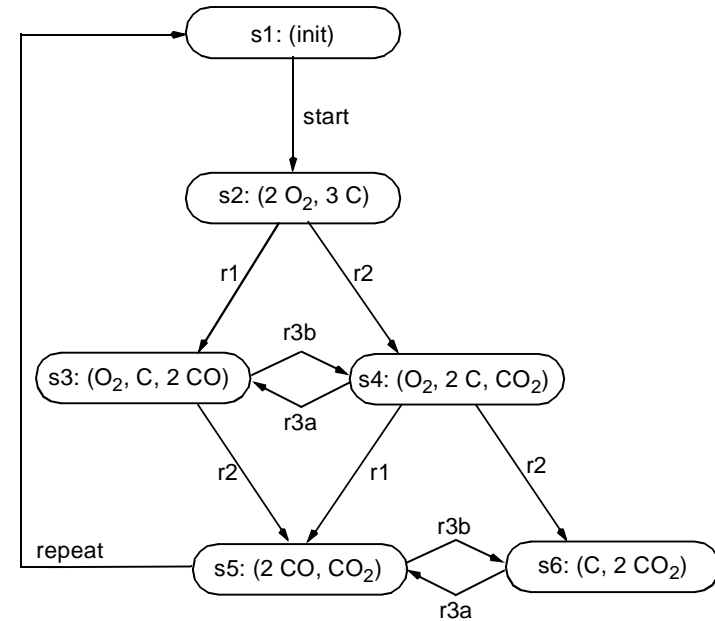


## BASIC PROPERTIES & RG, SUMMARY

- How many tokens may reside at most in a given place . . .
  - > (0, 1, k, oo) ?
  - > boundedness
  - > rg is finite
  
- How often may a transition fire . . .
  - > (0-times, n-times, oo-times) ?
  - > liveness
  - > every terminal scc contains all transitions
  
- Is the initial system state . . .
  - > always reachable again ?
  - > reversibility
  - > rg is sc (consists of one scc)

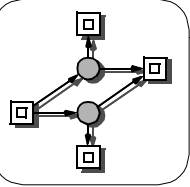


## RG(CARBON OXIDATION), EVALUATION



- RG is finite
  - > *BND*
  
- 1 Strongly Connected Component (SCC)
  - > *REV*
  
- the only SCC contains all transitions
  - > *LIVE*





## REACHABILITY GRAPH, CONSTRUCTION ALGORITHM

PROCEDURE rg (IN Net  $pn$ , IN Marking  $m_0$ ,  
OUT MSet  $nodes$ , OUT ArcSet  $arcs$ );

MSet  $U = \{m_0\}$ , // unprocessed markings  
 $N = \emptyset$ ; // rg nodes  
 ArcSet  $E = \emptyset$ ; // rg arcs (pre, post, t)  
 Marking  $m'$ ; // successor marking  
 Transition  $t$ ;

WHILE  $U \neq \emptyset$  DO

choose one  $m \in U$ ;  
 $U = U - \{m\}$ ;  $N = N \cup \{m\}$ ;

FOR ALL  $t$  enabled at  $m$  DO

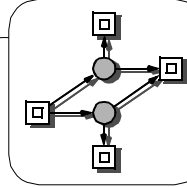
$m' = m + \Delta t$ ;  
 IF  $m' \notin N \cup U$  // new marking  
 THEN  $U = U \cup \{m'\}$   
 ENDIF;  
 $E = E \cup \{(m, m', t)\}$

ENDFOR

ENDWHILE;

$nodes = N$ ;  $arcs = E$ ;

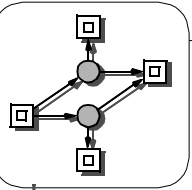
ENDPROC rg.



## REACHABILITY GRAPH, OBSERVATIONS

- ❑ **unbounded** Petri net  
 -> the RG is **infinite**
- ❑ **bounded** Petri net  
 -> the RG is **finite**
- ❑ simple construction algorithm  
 -> *single step firing rule*
- ❑ concurrency  
 -> *enumeration of all interleaving sequences*
- ❑ **branching arcs in the RG**  
 -> **conflict** **OR**  
 -> **concurrency**
- ❑ RG tend to be very large  
 -> *automatic evaluation necessary*
- ❑ **worst case: over-exponential growth**  
 -> *alternative analyses techniques ?*

## QUALITATIVE ANALYSIS METHODS, OVERVIEW

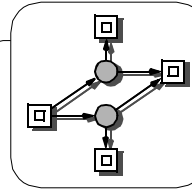


- NET REDUCTION
- STRUCTURAL PROPERTIES
- LINEAR PROGRAMMING**
  - place / transition invariants
  - state equation
  - trap equation
- STATE SPACE ANALYSIS
  - (complete) reachability graph
  - compressed state spaces
    - BDDs, NDDs, ..., xDDs
    - Kronecker products
  - reduced state spaces
    - coverability graph
    - symmetry
    - stubborn sets
  - branching process

static analysis

dynamic analysis  
(model checking)

## INCIDENCE MATRIX C - A REPRESENTATION OF THE NET STRUCTURE



	P+T				
P+T	p1	card(P)	t1	card(T)	
p1		$\phi$		- PRE	
card(P)					
t1		+ POST		$\phi$	
card(T)					

↓ POST<sup>T</sup> - PRE

C =

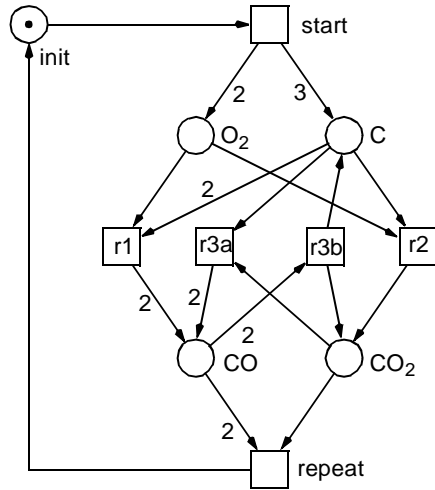
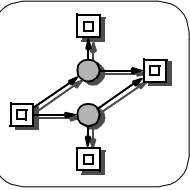
P \ T	t1	...	tj	...	tm
p1					
⋮					
pi			cij		
⋮					
pn					

$$cij = (pi, tj) = F(tj, pi) - F(pi, tj) = \Delta tj(pi)$$

-> token change  
in place pi by firing of transition tj

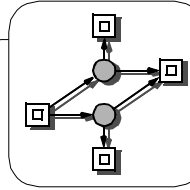
-> stoichiometric matrix

### CARBON/BND, INCIDENCE MATRIX



P \ T	r1	r2	r3a	r3b	start	repeat
O <sub>2</sub>	-1	-1	0	0	+2	0
C	-2	-1	-1	+1	+3	0
CO	+2	0	2	-2	0	-2
CO <sub>2</sub>	0	+1	-1	+1	0	-1
init	0	0	0	0	-1	+1

### STATE EQUATION



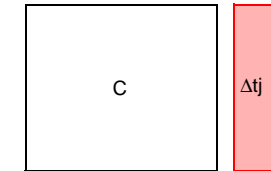
incidence matrix C

P \ T	t1	tj	tm
p1			
pi		$\Delta t_j$	
pn			

PARIKH VECTOR  
 $\text{parikh}(t_j)$

:
0
:
1
:
0
:

$t_j$



$\Delta t_j$  - vector describing the change of the whole marking by firing of  $t_j$

Let the word  $w \in T^*$  be a sequence of firing transitions.

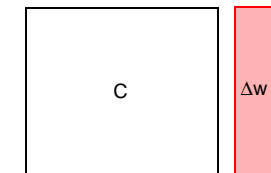
$\text{parikh}(w)$  - transition vector, whereby the position  $j$  gives the amount of  $t_j$  in  $w$

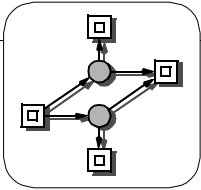
The change on the marking  $\Delta w$  by firing a transition sequence  $w$  can be computed by multiplying the incidence matrix C with the Parikh vector  $\text{parikh}(w)$  of that transition sequence.

PARIKH VECTOR  
 $\text{parikh}(w)$

1
2
:
:
0
:
:
3

$t_0$   
 $t_1$



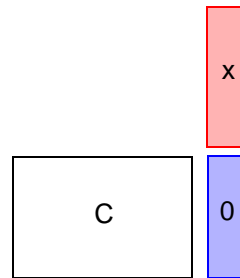


### T-INVARIANTS

- Lautenbach, 1973

- T-invariants

- > *integer solutions  $x$  of  $Cx = 0, x \neq 0, x \geq 0$*
- > *Parikh vector*
- > *exponential complexity*



- **minimal** T-invariants

- > *there is no T-invariant with a smaller support*
- > *greatest common divisor (gcd) of all entries is 1*

- support

- > *set of transitions belonging to the T-invariant*

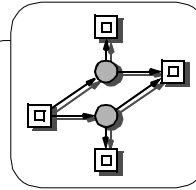
- any T-invariant is a non-negative linear combination of minimal ones

- > *multiplication with a positive integer*
- > *addition*
- > *division by gcd*

$$kx = \sum_i a_i x_i$$

- Covered by T-Invariants (CTI)

- > *each transition belongs to a T-invariant*
- > *if a bounded net is live, then it is CTI*



### T-INVARIANTS, INTERPRETATION

- T-invariants = (multi-) sets of transitions

- > *zero effect on marking*
- > *reproducing a marking / system state*
- > *steady state substance flows*
- > *elementary modes, Schuster 1993*

- the T-invariant corresponds to cycles in the RG, if the T-invariant is realizable

- in the RG, concurrency of transitions is described by all transitions' interleaving sequences

- if there are concurrent transitions in a realizable T-invariant, then there is a RG cycle for each interleaving sequence

- > *T-inv3, T-inv4*

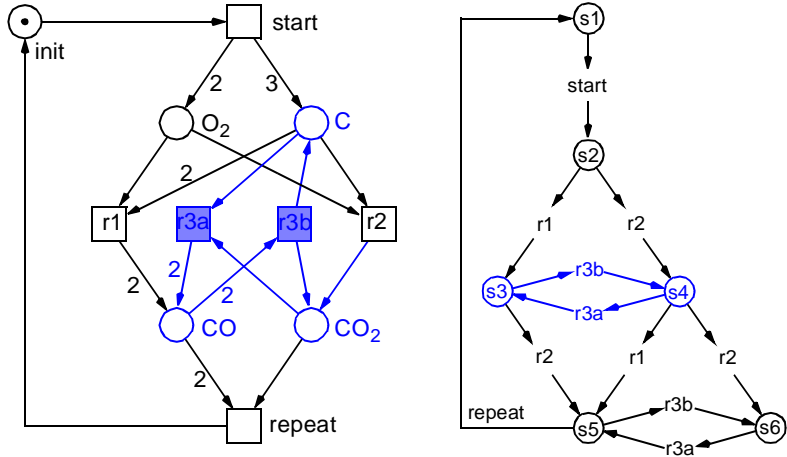
- pre-sets of supports = post-sets of supports

- a T-invariant defines a subnet

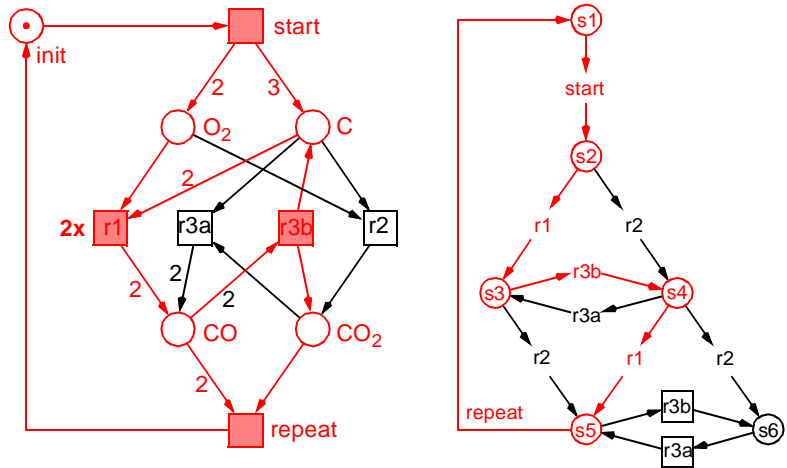
- > *the T-invariant's transitions (the support),*  
*+ all their pre- and post-places*  
*+ the arcs in between*

### CARBON/BND, T-INVARIANTS 1, 2

T-inv1 = (r3a, r3b) -> inner cycle

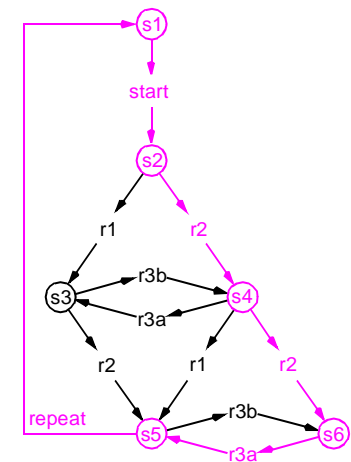
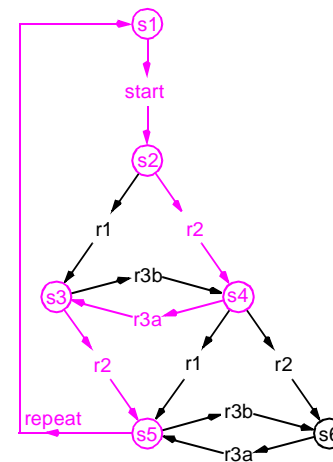
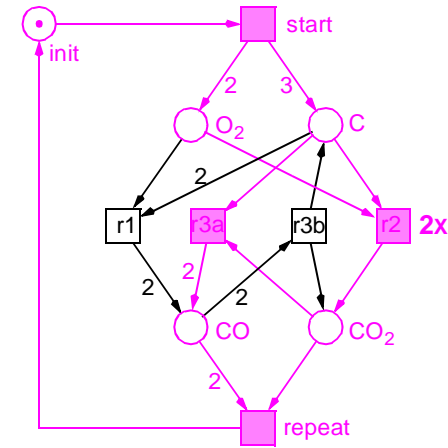


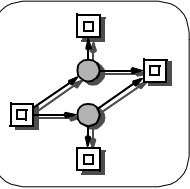
T-inv2 = (start, 2 r1, r3b, repeat) -> input/output cycle



### CARBON/BND, T-INVARIANTS 3

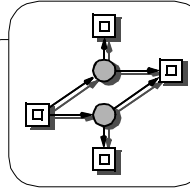
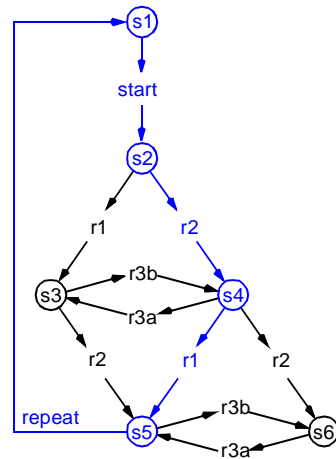
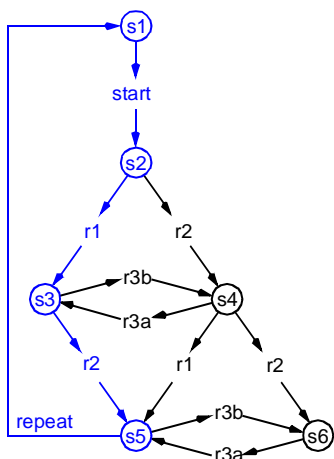
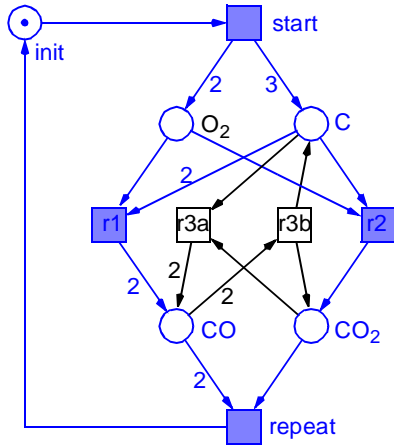
T-inv3 = (start, 2 r2, r3a, repeat) -> start - r2 < r3a > repeat





### CARBON/BND, T-INVARIANTS 4

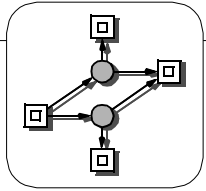
T-inv4 = (start, r1, r2, repeat) → start  $\langle \begin{matrix} r1 \\ r2 \end{matrix} \rangle$  repeat



### CARBON/UNBOUNDED, T-INVARIANTS, INTERPRETATION

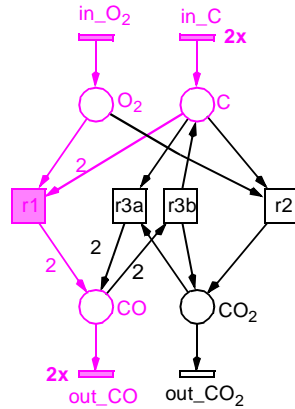
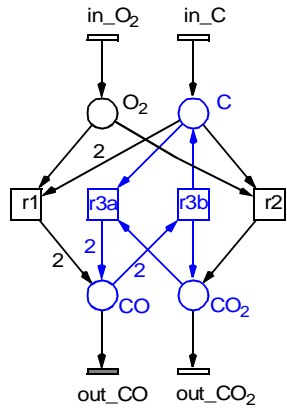
- steady state = constant token distribution
  - preservation of a given system state under continuous firing requires
    - > relative transition firing rates = T-invariant's entries
    - > ex T-inv2: a given state is preserved, if in\_C and out\_CO fire twice as often as in\_O2 and r1;
  - the in- / out-components of the T-invariant
    - > sum equation of the T-invariants remaining transitions
- T-inv1: --  
-> inner cycle
- T-inv2:  $O_2 + 2 C \rightarrow 2 CO$   
-> stoichiometric equation of r1
- T-inv3:  $C + O_2 \rightarrow CO_2$   
-> stoichiometric equation of r2
- T-inv4:  $2 C + O_2 \rightarrow 2 CO$   
-> sum of the stoichiometric equations of r2, r3a
- T-inv5:  $C + O_2 \rightarrow CO_2$   
-> sum of the stoichiometric equations of r1, r3b

### CARBON/UNBOUNDED, T-INVARIANTS 1 - 3

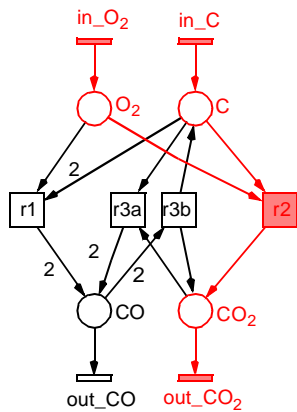


T-inv1 = (r3a, r3b)

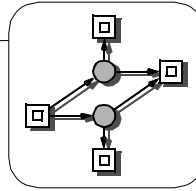
T-inv2 = (in\_O<sub>2</sub>, 2 in\_C, r1, 2 out\_CO)



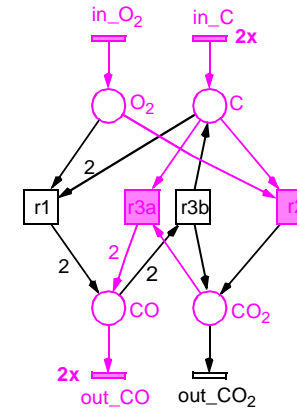
T-inv3 = (in\_O<sub>2</sub>, in\_C, r2, out\_CO<sub>2</sub>)



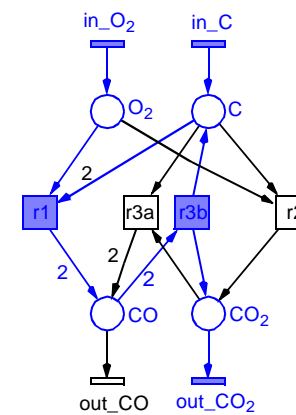
### CARBON/UNBOUNDED, T-INVARIANTS 4, 5

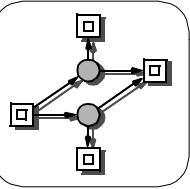


T-inv4 = (in\_O<sub>2</sub>, 2 in\_C, r2, r3a, 2 out\_CO)



T-inv5 = (in\_O<sub>2</sub>, in\_C, r1, r3b, out\_CO<sub>2</sub>)





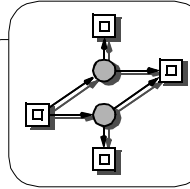
## T-INVARIANTS, SUMMARY

### TWO INTERPRETATIONS

- state-reproducing transition sequence (partial order) of transitions occurring one after the other
- relative transition firing rates of transitions occurring permanently & concurrently

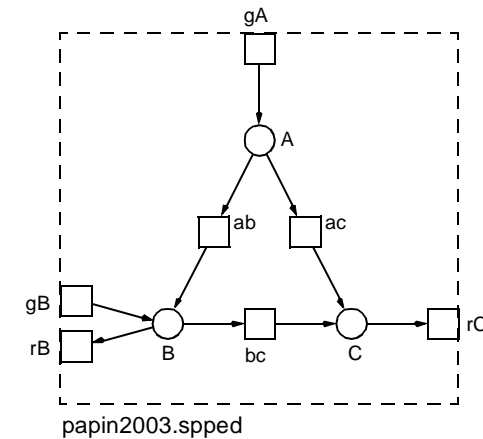
### BASIC TYPES IN BIO NETWORKS

- trivial minimal T-invariants
  - > *boundary transitions of auxiliary compounds*
  - > *reversible reactions*
- non-trivial minimal T-invariants
  - > *i/o-T-invariants*  
*covering boundary transitions of input / output compounds*
  - > *inner cycles*



## PATHWAY ANALYSIS

- substances involved
  - > *input substance A*
  - > *output substance C*
  - > *auxiliary substance B*

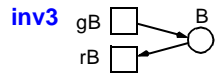
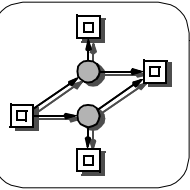


- steady state substance flows
  - > *T-invariants*
- all flow behaviour under the steady state assumption
  - > *non-negative linear combination of minimal T-invariants*

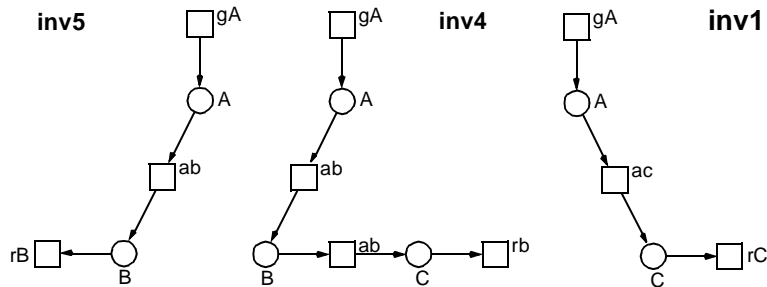
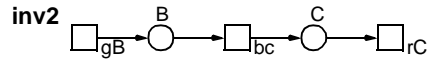


## T-INVARIANTS AND EXTREME PATHWAYS

Schilling, 2000

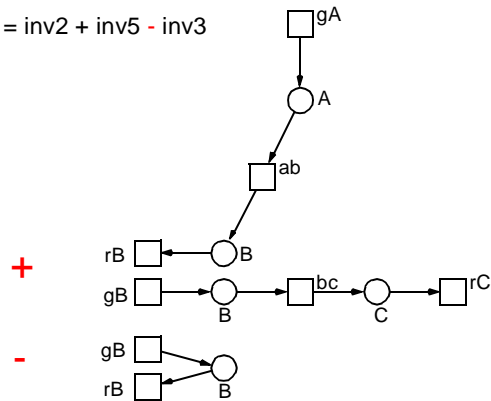


no elementary mode

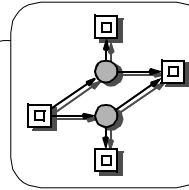


no extreme pathway

$inv4 = inv2 + inv5 - inv3$



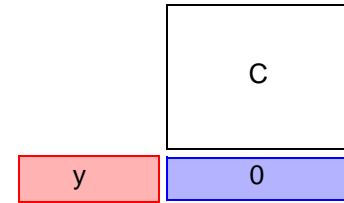
## P-INVARIANTS



□ Lautenbach, 1973

□ P-invariants

-> integer solutions  $y$  of  $yC = 0, y \neq 0, y \geq 0$



-> exponential complexity

□ minimal P-invariants

-> there is no P-invariant with a smaller support

-> gcd of all entries is 1

□ support

-> set of places belonging to the P-invariant

□ any P-invariant is a non-negative linear combination of minimal ones

-> multiplication with a positive integer

-> addition

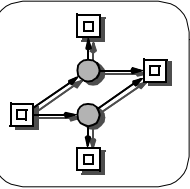
-> division by gcd

$$ky = \sum_i a_i y_i$$

□ Covered by P-Invariants (CPI)

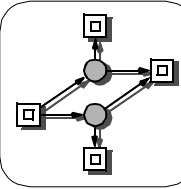
-> each place belongs to a P-invariant

-> sufficient condition for BND



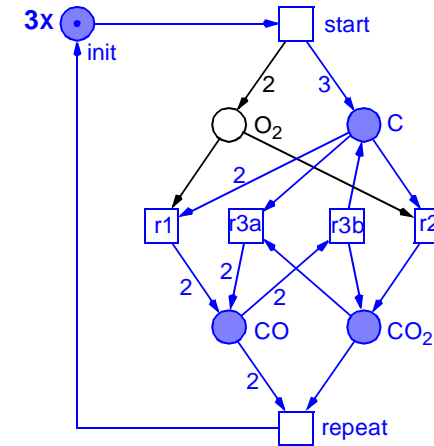
### P-INVARIANTS, INTERPRETATION

- set of places with
  - > a constant weighted sum of tokens
  - $ym = ym_0$  for all reachable markings
  - > token / compound preservation
- a place belonging to a P-invariant is bounded
  - > CPI - sufficient condition for BND
- the firing of any transition has no influence on the weighted sum of tokens on the P-invariant's places
  - > for all transition  $t$ :
    - the effect of the arcs,
    - removing tokens from a P-invariant's place
    - is equal to the effect of the arcs
    - adding tokens to a P-invariant's place
- pre-sets of supports = post-sets of supports
- a P-invariant defines a subnet
  - > the P-invariant's places (the support),
  - + all their pre- and post-transitions
  - + the arcs in between

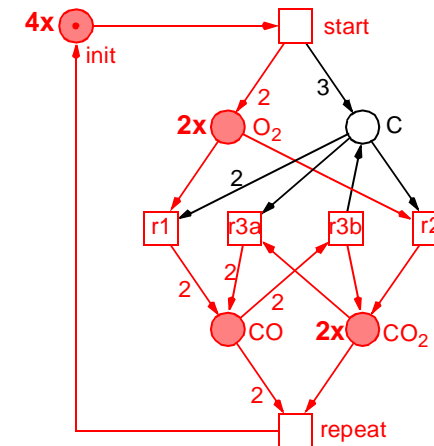


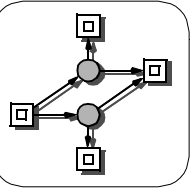
### CARBON/BND, P-INVARIANTS

P-inv1 = (3 init, C, CO, CO<sub>2</sub>) -> carbon preservation



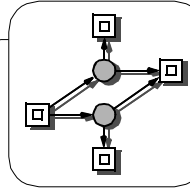
P-inv2 = (4 init, 2 O<sub>2</sub>, CO, 2 CO<sub>2</sub>) -> oxygen preservation





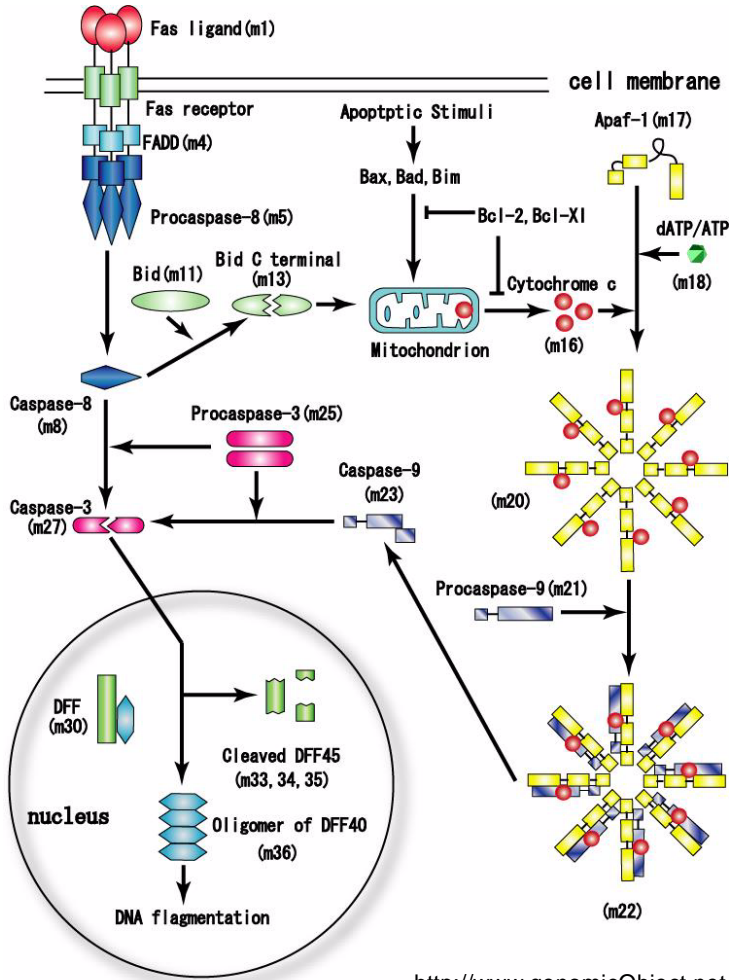
## ANALYSIS, SUMMARY

- ❑ validation criterion 1
  - > CTI,  
stronger - covered by i/o T-invariants
  - > no minimal T-invariant  
without biological interpretation
  - > no known biological behaviour  
without corresponding T-invariant
- ❑ validation criterion 2
  - > no minimal P-invariant  
without biological interpretation (?)
- ❑ validation criterion 3
  - > CPI
  - > all expected temporal-logic properties -> TRUE  
(not discussed here)



## 4. CASE STUDIES

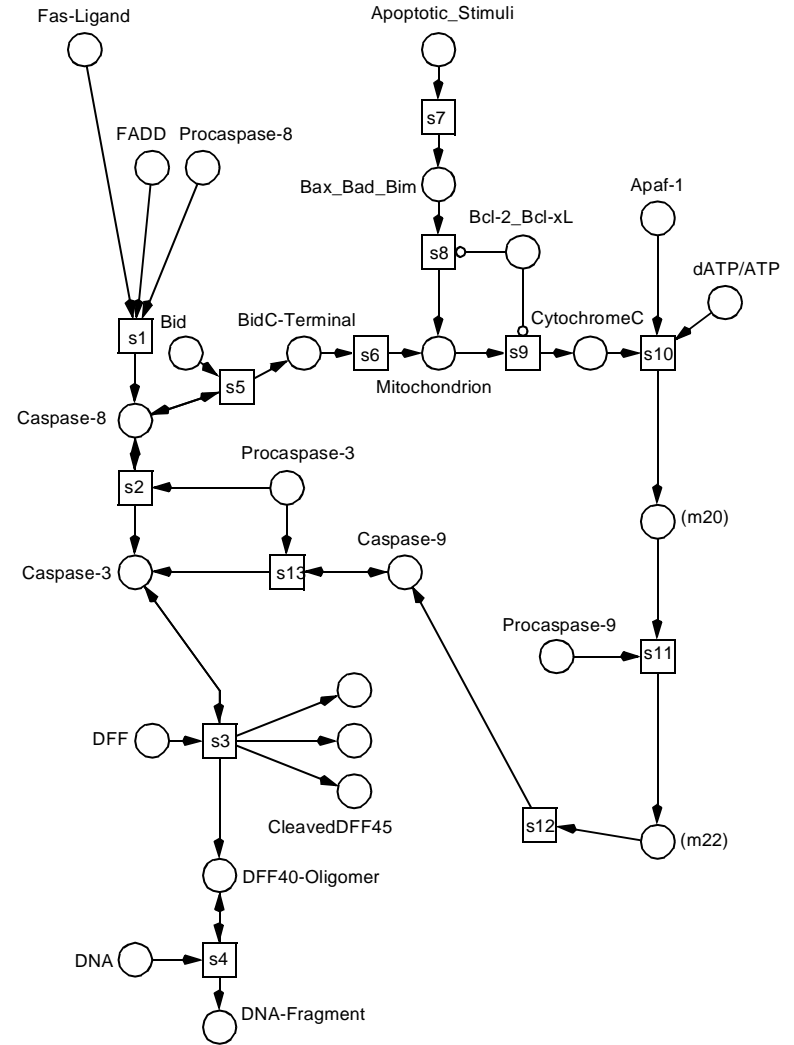
### CASE STUDY 1 - APOPTOSIS, THREE BASIC PATHWAYS



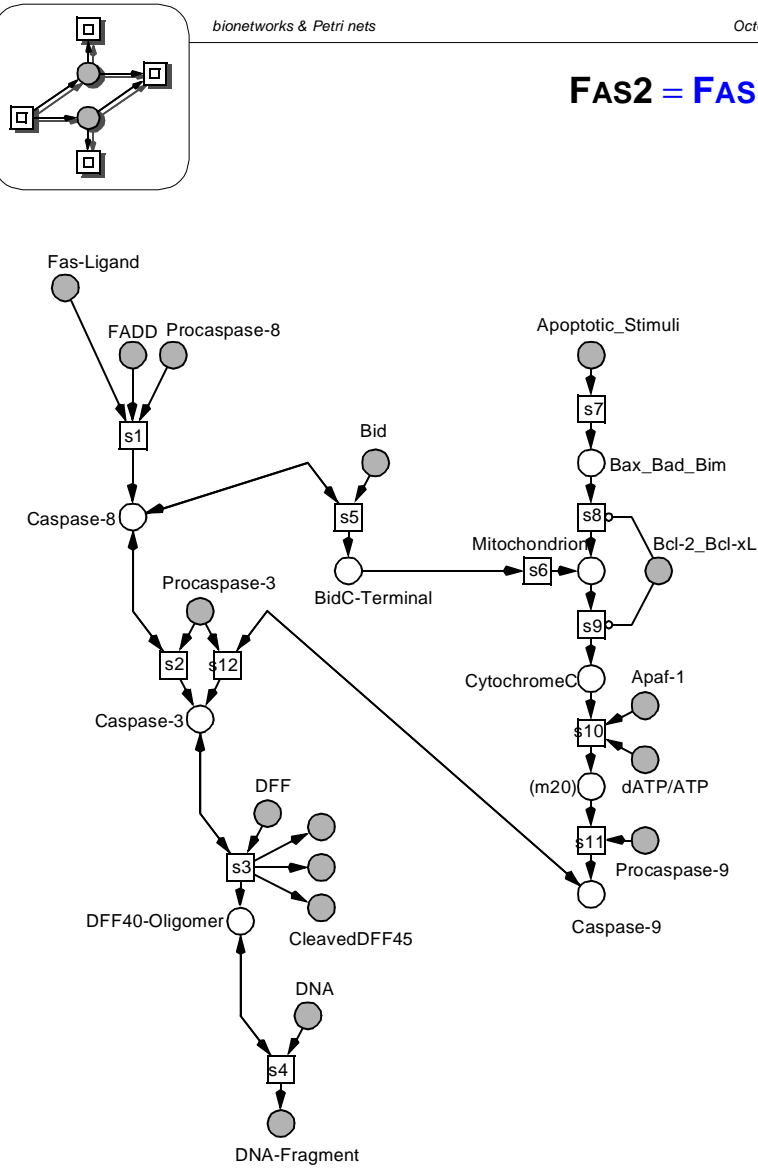
<http://www.genomicObject.net>

### APOPTOSIS IN MAMMALIAN CELLS

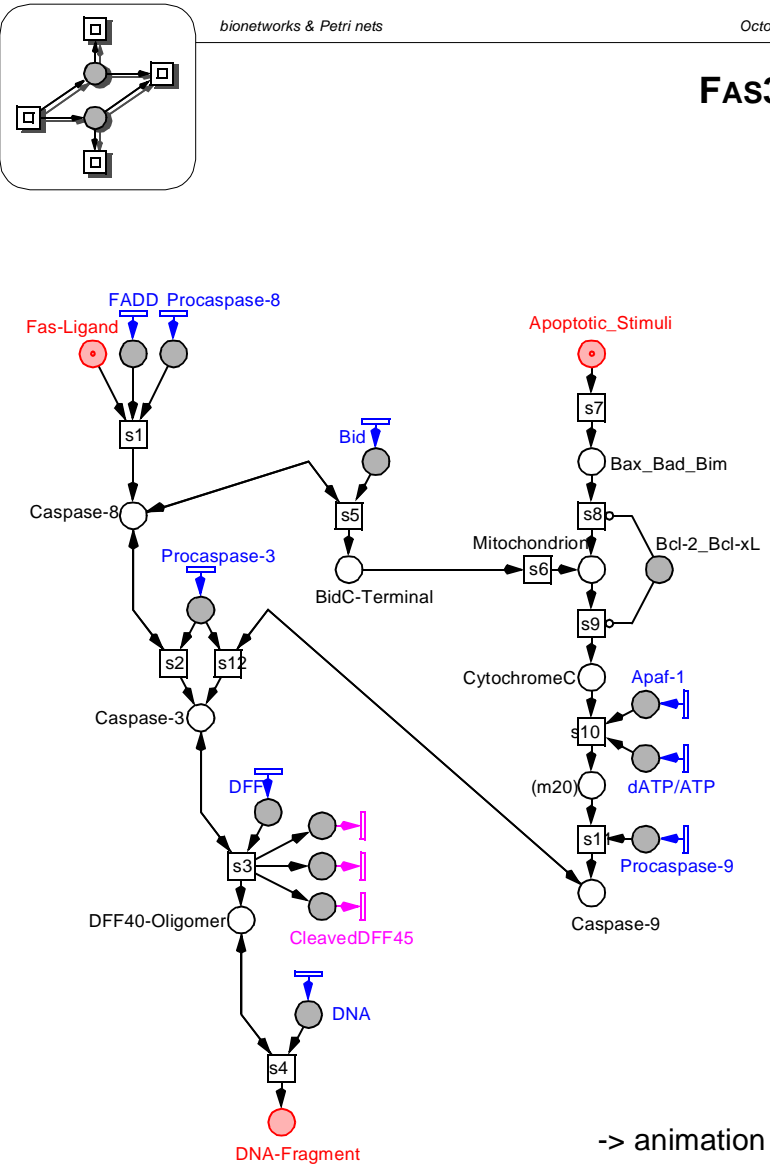
### FAS1



# FAS2 = FAS1



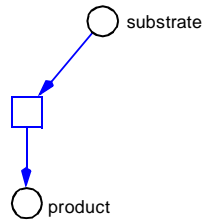
# FAS3



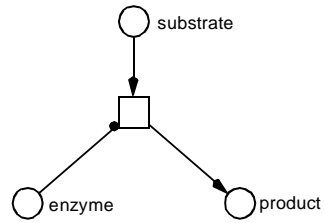
-> animation

### REFINEMENT: AUTOCATALYSIS

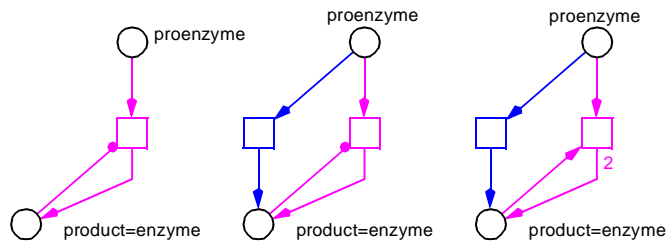
REACTION



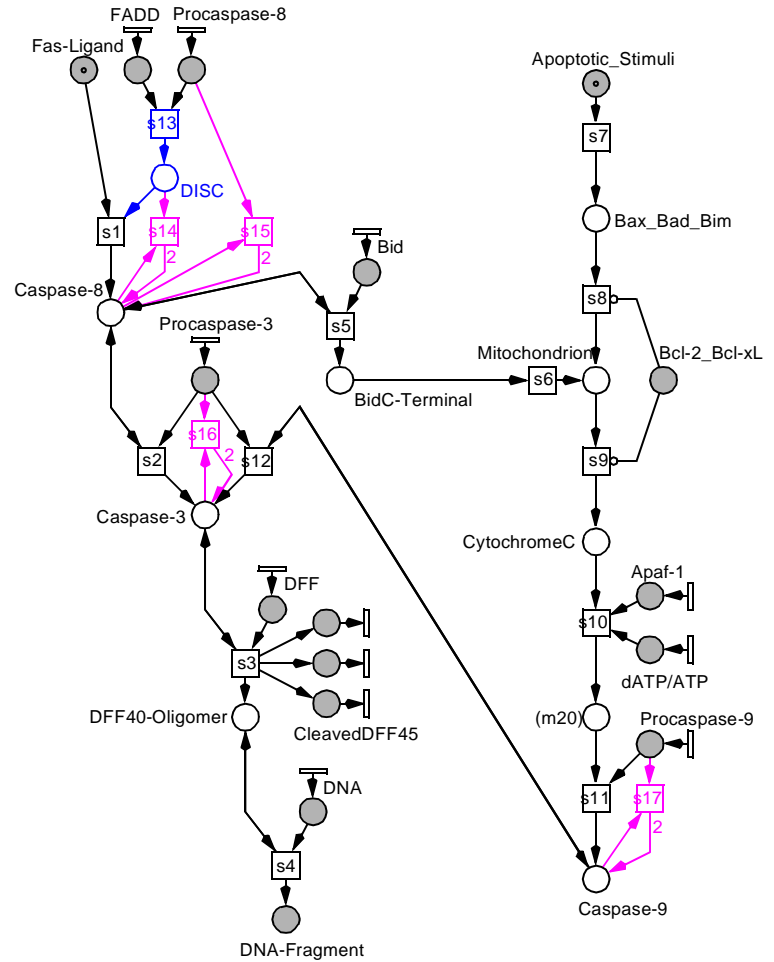
CATALYSIS



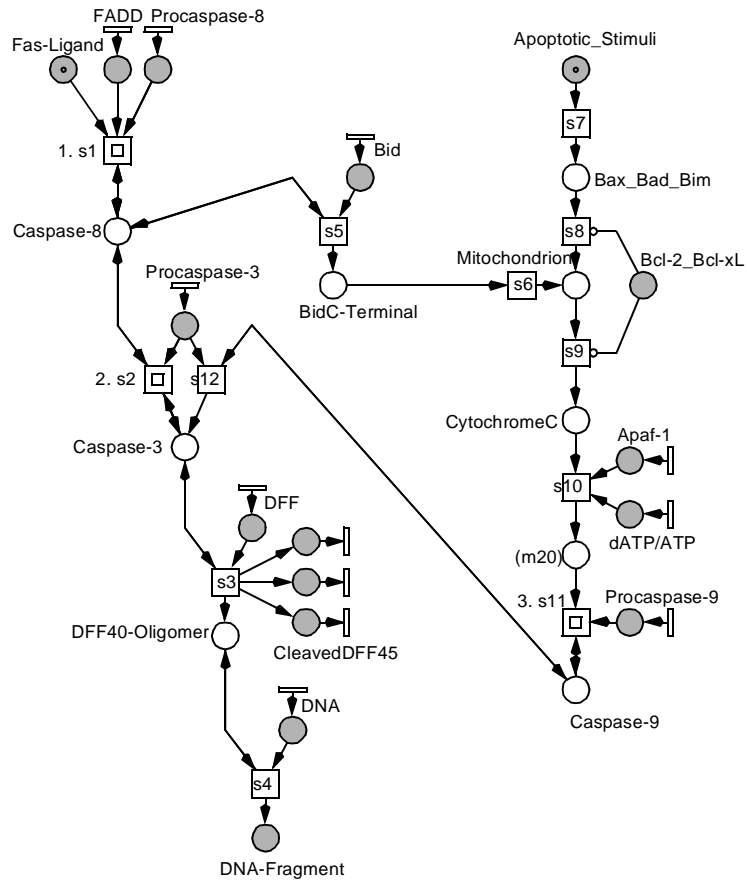
AUTOCATALYSIS



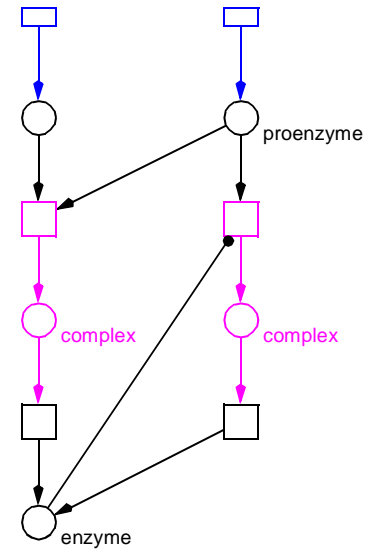
### FAS4A



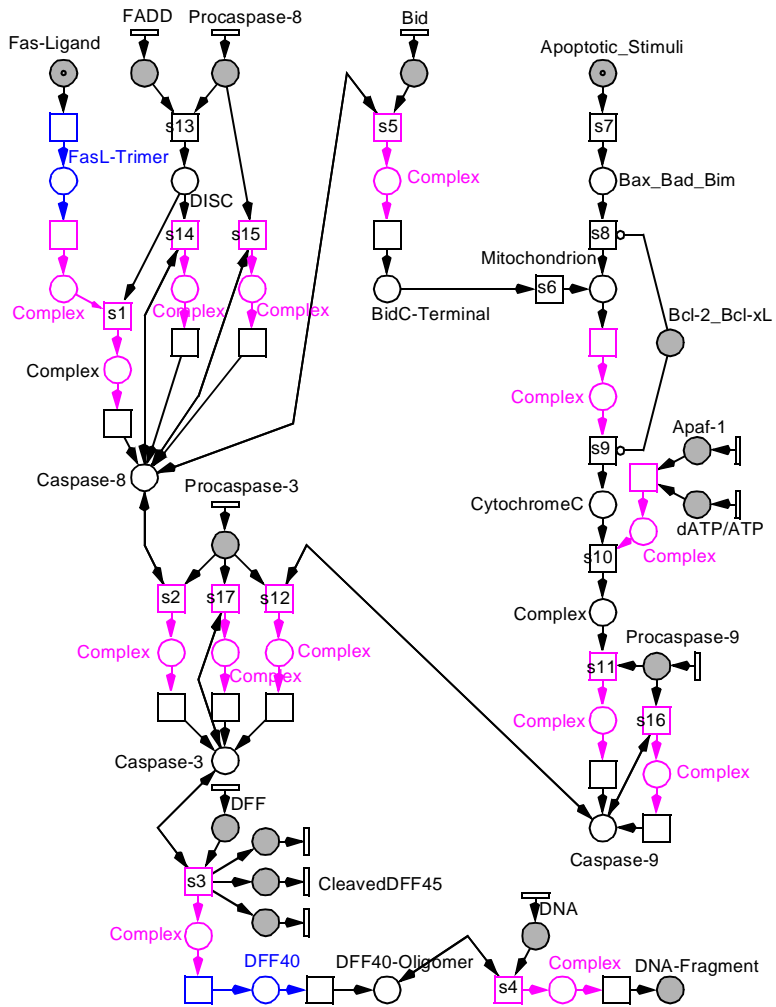
# FAS4B ≈ FAS3



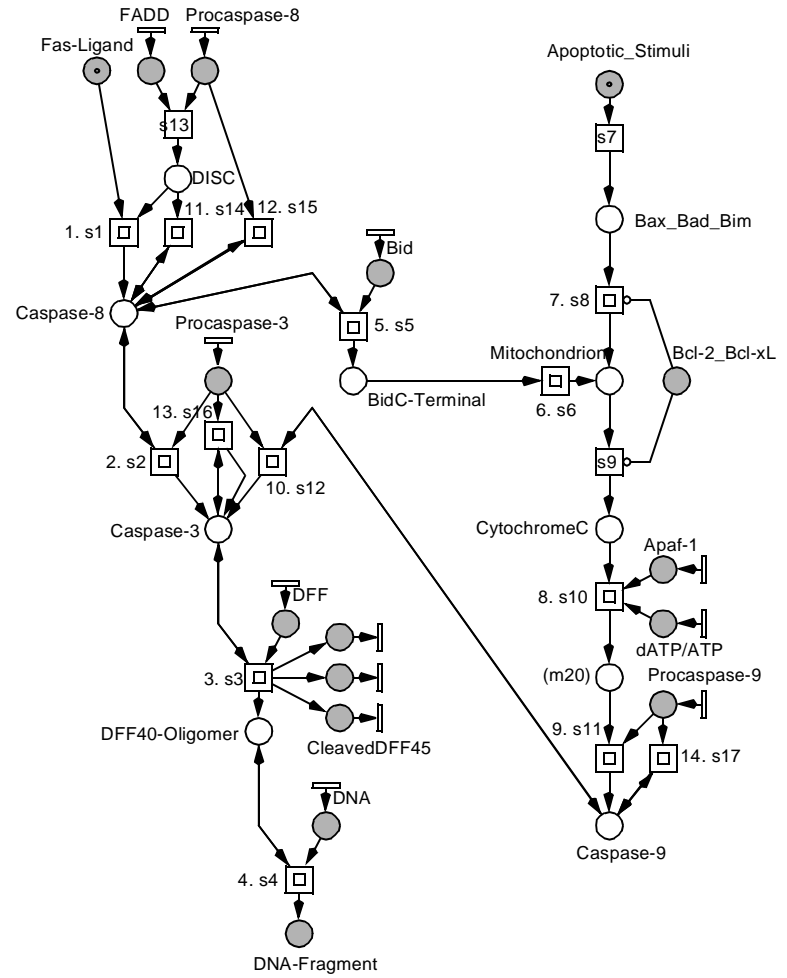
# REFINEMENT: INTERMEDIATE COMPLEXES



### FAS5A

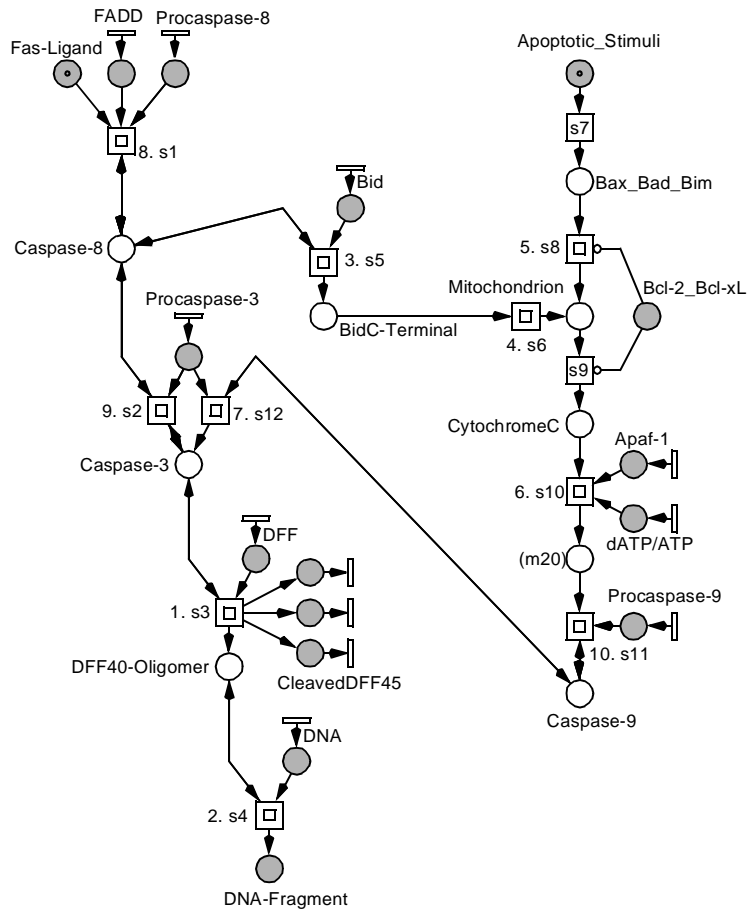


### FAS5B ≈ FAS4A

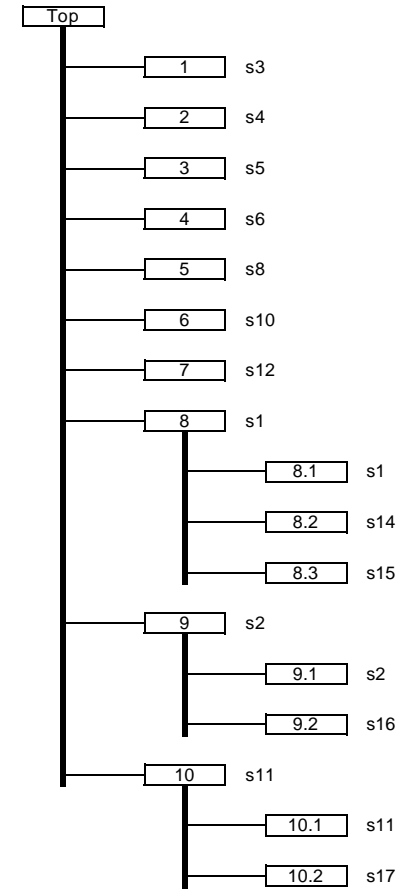


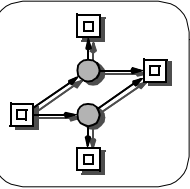


### FAS5C ≈ FAS3



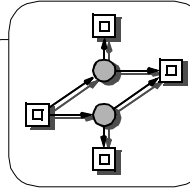
### HIERARCHY TREE FAS5C



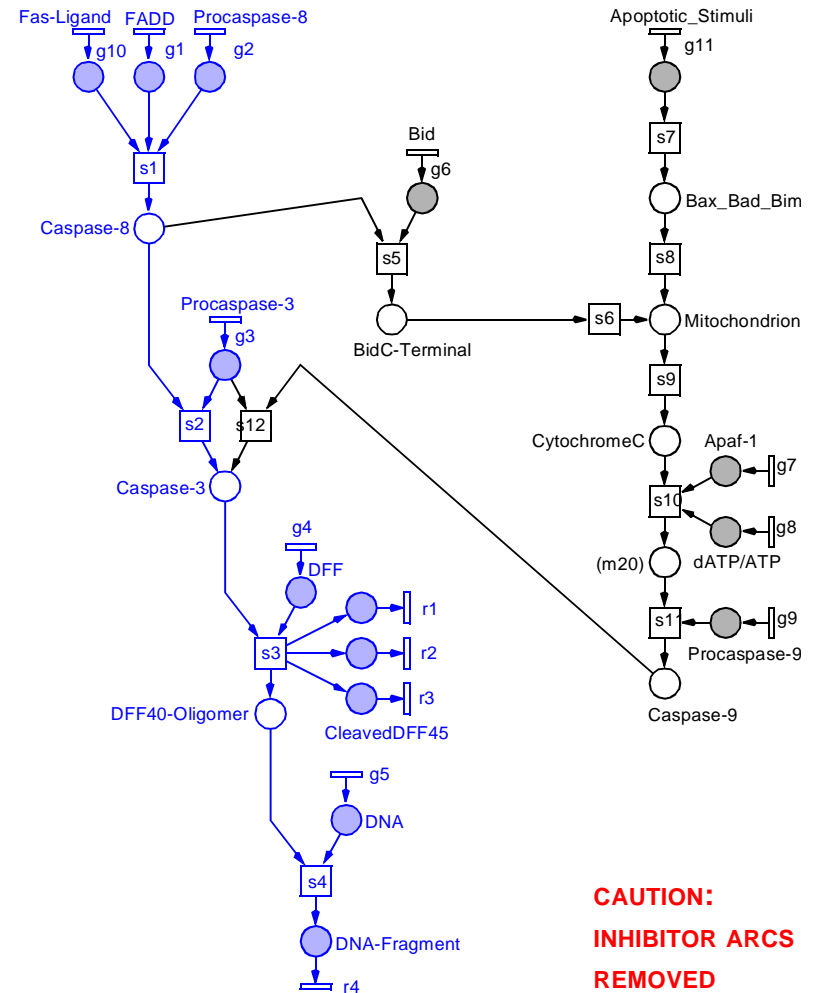


### STEP-WISE MODELLING

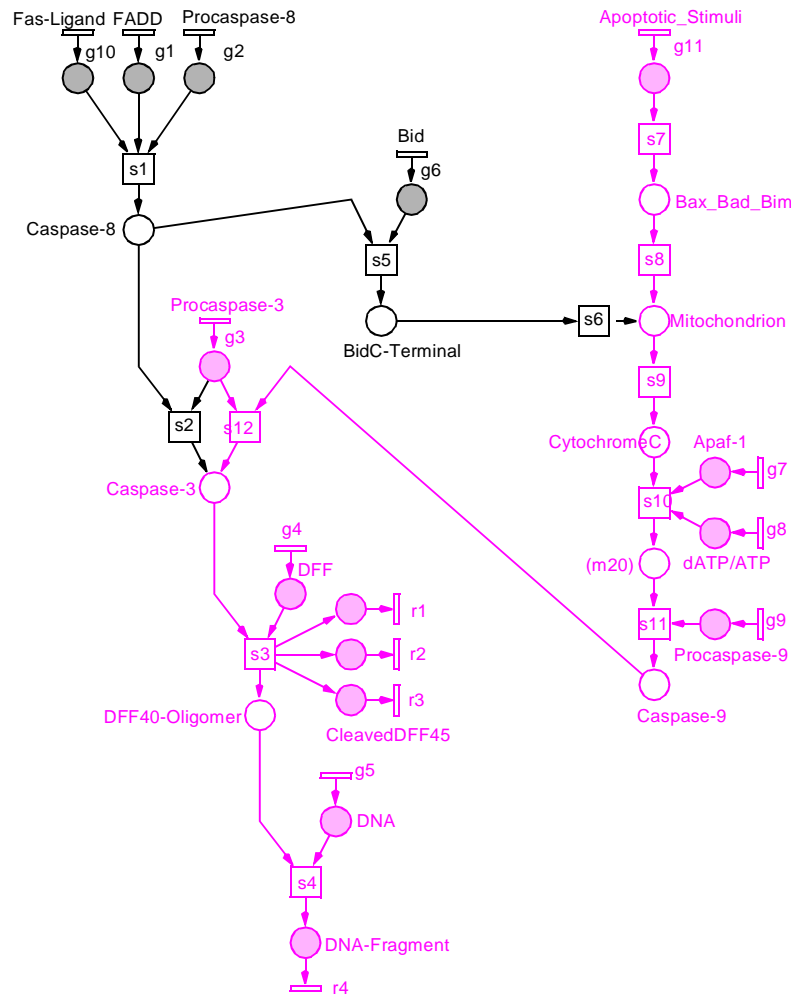
- 1. literal scheme transformation Fas1
- 2. layout improvement Fas2
  - > use of syntactic sugar
- 3. adding environment behaviour Fas3
  - > animation
- 4. adding autocatalysis Fas4a
  - > hierarchic Petri net Fas4b
- 5. adding intermediate complexes Fas5a
  - > refined hierarchies Fas5b
  - Fas5c
- >> EXPLOIDING SYNTACTIC SUGAR**
- 6. abstraction for analysis Fas6



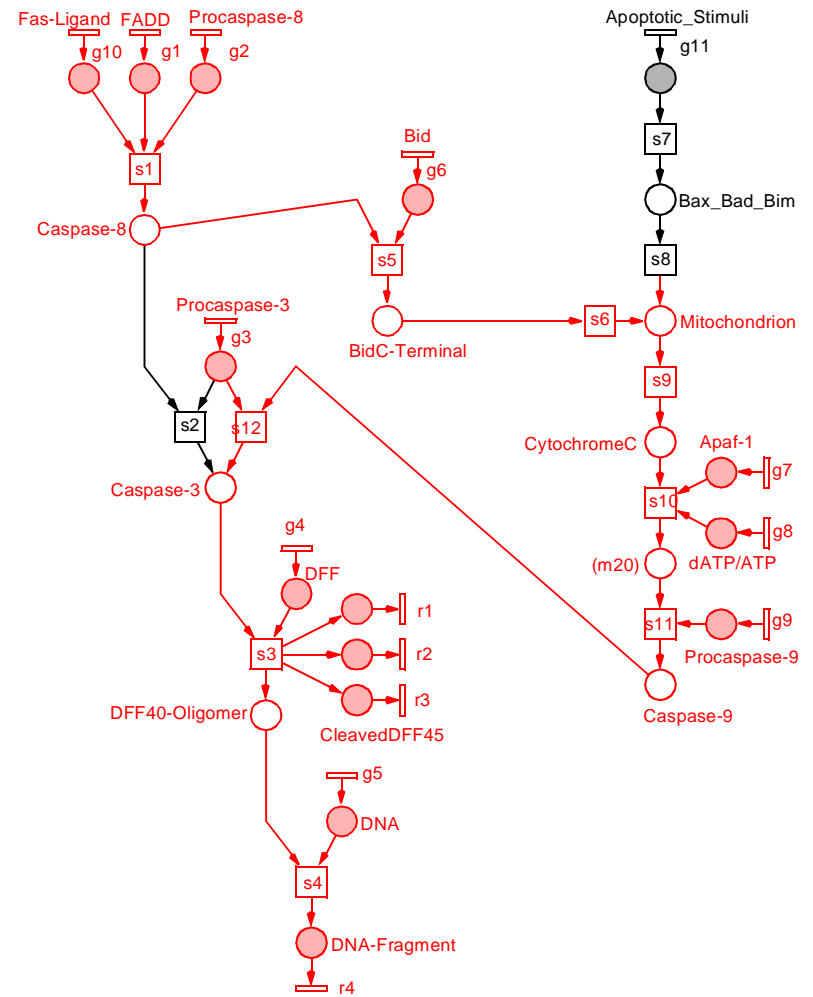
### ANALYSIS OF FAS6, T- INVARIANT 1, DEATH-RECEPTOR PATHWAY

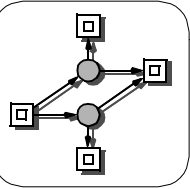


## ANALYSIS OF FAS6, T- INVARIANT 2, MITOCHONDRIAL PATHWAY



## ANALYSIS OF FAS6, T- INVARIANT 3, CROSS-TALK BY BID

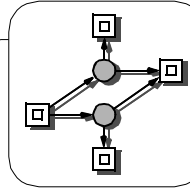




### APOPTOSIS, ANALYSES SUMMARY

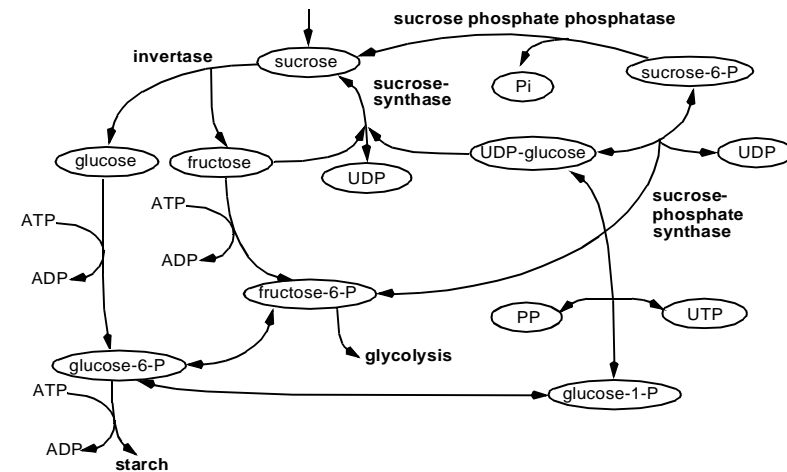
- ❑ environment style 1
- ❑ size of the net - 24 places / 27 transitions
- ❑ many read arcs, resolved for analysis
- ❑ no P-invariants
- ❑ three minimal positive i/o T-invariants
  - > *three basic behaviours*
  - > *any net behaviour = non-negative linear combination of them*
- ❑ covered by i/o T-invariants
  - > *no idle parts*
- ❑ reproducible empty marking (guess)
  - > *cyclic behaviour possible (reversibility)*
- ❑ INA result vector

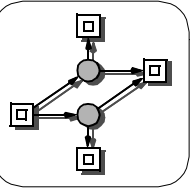
ORD	HOM	NBM	PUR	CSV	SCF	CON	SC	Ft0	tF0	Fp0	pF0	MG	SM	FC	EFC	ES
Y	Y	Y	Y	N	N	Y	N	Y	Y	N	N	N	N	N	N	N
DTP	SMC	SMD	SMA	CPI	CTI	B	SB	REV	DSt	BSt	DTr	DCF	L	LV	L&S	
Y	N	N	N	N	Y	N	N	?	N	?	N	?	Y	Y	N	



### CASE STUDY 2 - POTATO TUBER

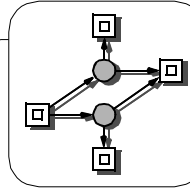
- ❑ central carbon metabolism in potato tubers
  - > *stoichiometric relations known*
  - > *non-ordinary place/transition net*
  - > *many reversible reactions*
- ❑ schematic overview



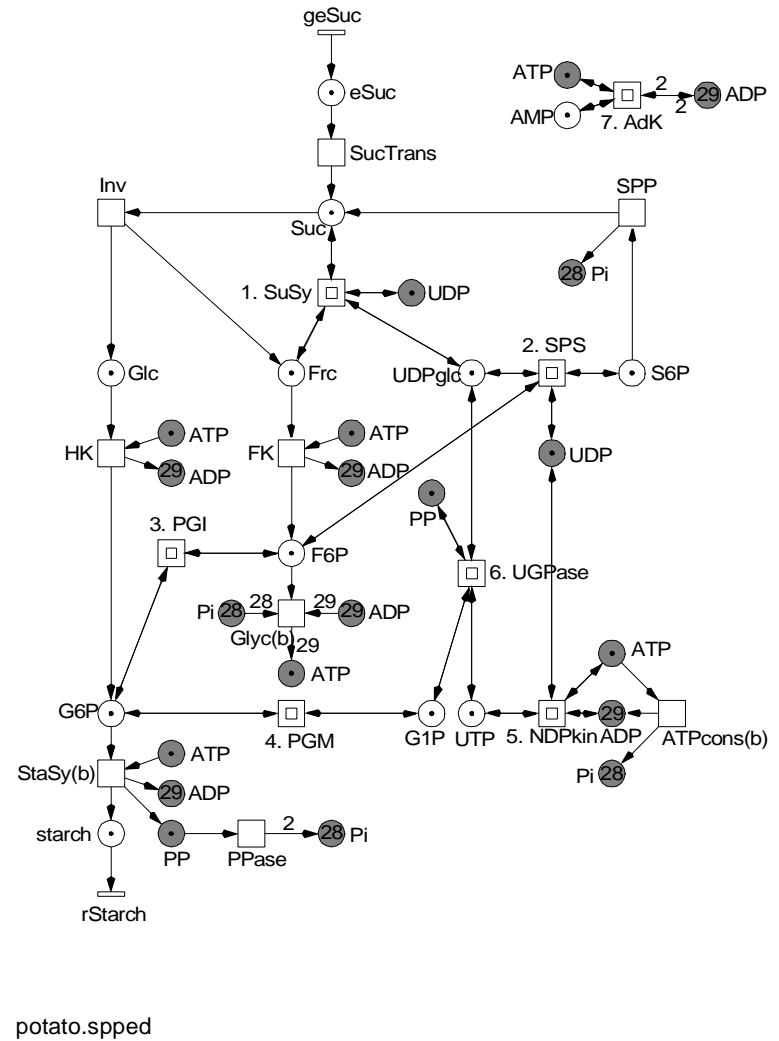


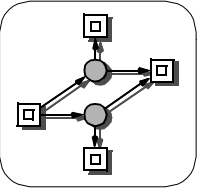
### STOICHIOMETRIC EQUATIONS

- R1. SuSy: *sucrose synthase*  
 $Suc + UDP \leftrightarrow UDPglc + Frc$
- R2. UGPase: *UDPglucose pyrophosphorylase*  
 $UDPglc + PP \leftrightarrow G1P + UTP$
- R3. PGM: *phosphoglucomutase*  
 $G6P \leftrightarrow G1P$
- R4. FK: *fructokinase*  
 $Frc + ATP \rightarrow F6P + ADP$
- R5. PGI: *phosphoglucose isomerase*  
 $G6P \leftrightarrow F6P$
- R6. HK: *hexokinase*  
 $Glc + ATP \rightarrow G6P + ADP$
- R7. Inv: *invertase*  
 $Suc \rightarrow Glc + Frc$
- R8. Glyc(b): *glycolysis*  
 $F6P + 29 ADP + 28 P_i \rightarrow 29 ATP$
- R9. SPS: *sucrose phosphatase*  
 $F6P + UDPglc \leftrightarrow S6P + UDP$
- R10. SPP: *sucrose phosphate phosphatase*  
 $S6P \rightarrow Suc + P_i$
- R11. NDPkin: *NDP kinase*  
 $UDP + ATP \leftrightarrow UTP + ADP$
- R12. SucTrans: *sucrose transporter*  
 $eSuc \rightarrow Suc$
- R13. ATPcons(b): *ATP consumption*  
 $ATP \rightarrow ADP + P_i$
- R14. StaSy(b): *starch synthesis*  
 $G6P + ATP \rightarrow starch + ADP + PP$
- R15. AdK: *adenylate kinase*  
 $ATP + AMP \leftrightarrow 2 ADP$
- R16. PPase: *pyrophosphatase*  
 $PP \rightarrow 2 P_i$



### POTATO TUBER, PETRI NET

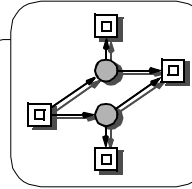




## POTATO TUBER, ANALYSES SUMMARY 1

- ❑ environment style 2
- ❑ size of the net - 17 places / 25 transitions
- ❑ three P-invariants, but not CPI
  - > *PI-1: UDPglc, UTP, UDP.*  
*uridine preservation*
  - > *PI-2: ATP, AMP, ADP.*  
*adenosine preservation*
  - > *PI-3: G6P, F6P, G1P, UTP, ATP(2),*  
*ADP, S6P, Pi, PP(2).*  
*preservation of phosphorylated metabolites*
- ❑ P-invariants need sufficient tokens at initial marking to make the net live
  - > *how to calculate ?*
- ❑ INA result vector

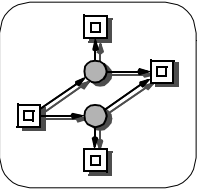
ORD	HOM	NBM	PUR	CSV	SCF	CON	SC	Ft0	tF0	Fp0	pF0	MG	SM	FC	EFC	ES
N	N	N	Y	N	N	Y	N	Y	Y	N	N	N	N	N	N	N
DTP	CPI	CTI	B	SB	REV	DSt	BSt	DTr	DCF	L	LV	L&S				
?	N	Y	N	N	?	?	?	?	?	?	?	N				



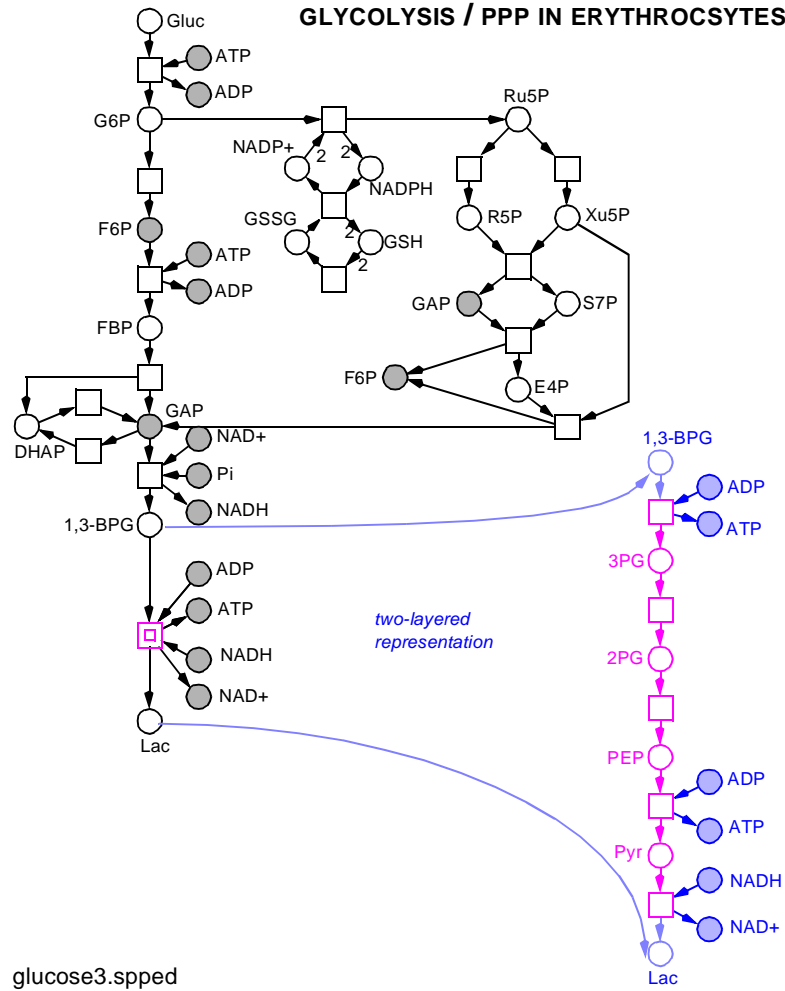
## POTATO TUBER, ANALYSES SUMMARY 2

- ❑ covered by i/o T-invariants
- ❑ 19 T-invariants
  - > *7 trivial ones (reversible reactions)*
  - > *12 i/o invariants*
  - > *no inner cycles*
- ❑ 3 i/o T-invariants with sucrose cleavage by sucrose synthase
  - > *e.g. TI-8: geSuc, SucTrans, SuSy(29), UGPase, PGM\_rev, FK(29), Glyc(b), StaSy(b), rStarch, SPS(28), SPP(28), NDPkin\_rev.*
- ❑ 9 i/o T-Invariants with sucrose cleavage by invertase
  - > *e.g. TI-11: geSuc, SucTrans, Inv(14), UGPase\_rev(13), PGM(13), HK(14), FK, Glyc(b), StaSy(b), rStarch, SuSy\_rev(13), NDPkin(13), PPase(14).*
- ❑ AdK and SPS\_rev
  - > *do not occur in a no-trivial T-invariant*
  - > *removing AdK and / or SPS\_rev has no influence on the system behaviour*

### CASE STUDY 3 - GLYCOLYSIS, VERSION 3

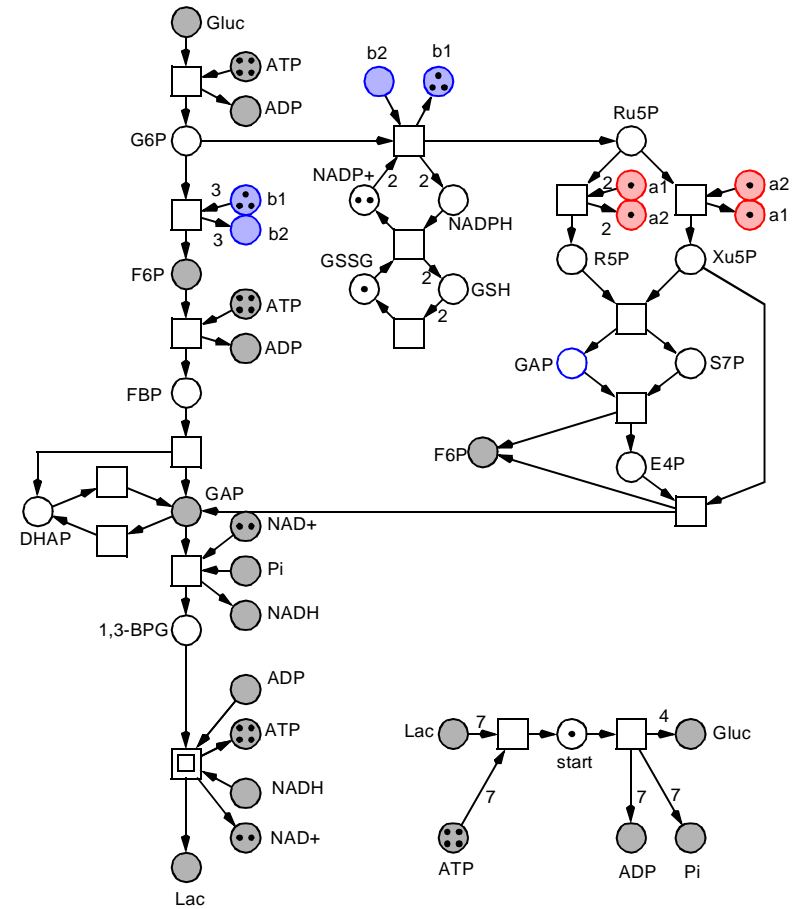
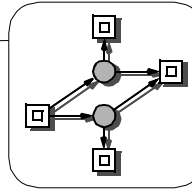


#### GLYCOLYSIS / PPP IN ERYTHROCYTES

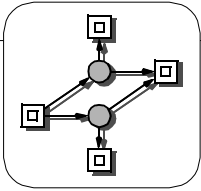


glucose3.spped

### GLYCOLYSIS, VERSION 4

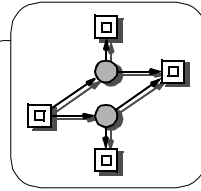


glukose4.spped



### ASSUMPTIONS IN VERSION 4

- ❑ the two appearances of GAP can be separated (no logical / fusion nodes)
  - ❑ the branching probabilities at the conflicts of G6P and Ru5P are known and may be characterized by the relations
    - G6P - 3 :1
    - Ru5P - 2 :1
- > *STEADY STATE:*  
*all intermediates have to be balanced with respect to inputs and outputs*

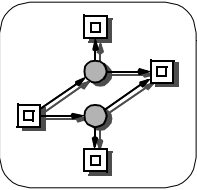


### GLYCOLYSIS, ANALYSES SUMMARY

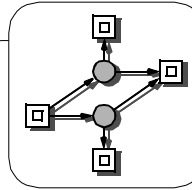
- ❑ environment style 3
- ❑ size of the net - 32 places, 22 transitions
- ❑ CPI
  - > BND
  - > 39 P-invariants
  - > interpretation ?
- ❑ CTI
  - > 1 trivial T-invariant
  - > 1 i/o invariant
- ❑ size of the reachability graph - 14.862
  - > live
  - > reversible
- ❑ INA result vector

ORD	HOM	NBM	PUR	CSV	SCF	CON	SC	Ft0	tF0	Fp0	pF0	MG	SM	FC	EFC	ES
N	Y	N	Y	N	N	Y	Y	N	N	N	N	N	N	N	N	Y
DTP	CPI	CTI	B	SB	REV	DSt	BSt	DTr	DCF	L	LV	L&S				
?	Y	Y	Y	Y	Y	N	?	N	N	Y	Y	N				





# 5. SUMMARY AND OUTLOOK



## BIO PETRI NETS, WHAT FOR ?

- **unifying view**
  - > *different biochemical systems*
  
- **step-wise model construction**  
of graphical (= visual) models for
  - > *animation*
  - > *validation*
  - > *(qualitative) analysis*
  - > *(quantitative) simulation*
  
- **integration of**
  - > *model validation*
  - > *behaviour prediction*
  
- **one all-purpose model**
  - > *animation model*
  - > *"qualitative model = animation model"*
  - > *"quantitative model =  
qualitative model  
+ quantitative parameter"*

### MODEL CLASSES

#### PETRI NETS

PLACE/TRANSITION  
PETRI NET  
(COLOURED PN)

validation by  
Petri net theory

validation/prediction  
by model checking

#### TIME-DEPENDENT PN

DISCRETE<sup>\*)</sup>  
PETRI NET

worst-case  
evaluation

CONTINUOUS  
PETRI NET

behaviour  
prediction

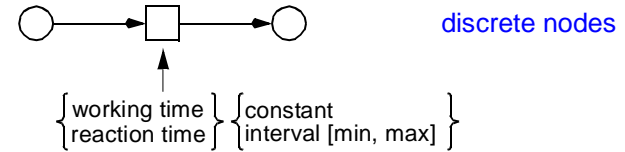
STOCHASTIC  
PETRI NET

reliability  
prediction

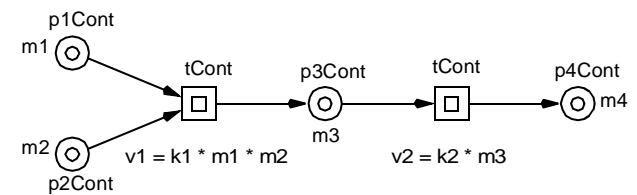
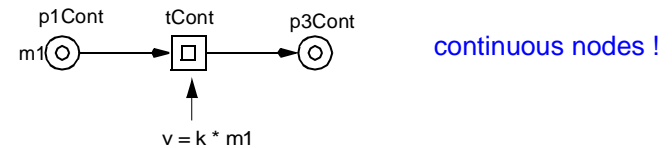
\*) DISCRETELY TREATABLE

### INTEGRATION OF QUANTITATIVE ANALYSES

- CONTINUOUS,  
BUT DISCRETELY TREATABLE TIME



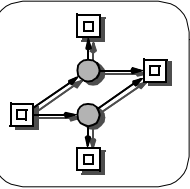
- CONTINUOUS TIME



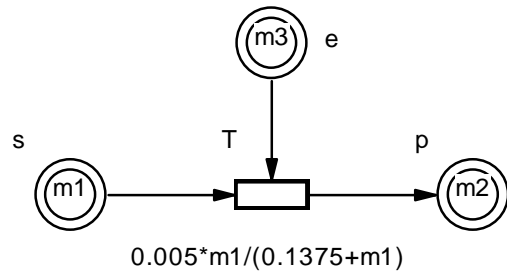
system of differential equations

$$\begin{cases} d [p1Cont] / dt = d [p2Cont] / dt = - v1 \\ d [p4Cont] / dt = v2 \\ d [p3Cont] / dt = v1 - v2 \end{cases}$$

-> SELF-MODIFYING PETRI NETS



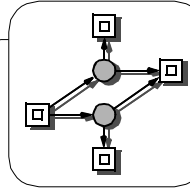
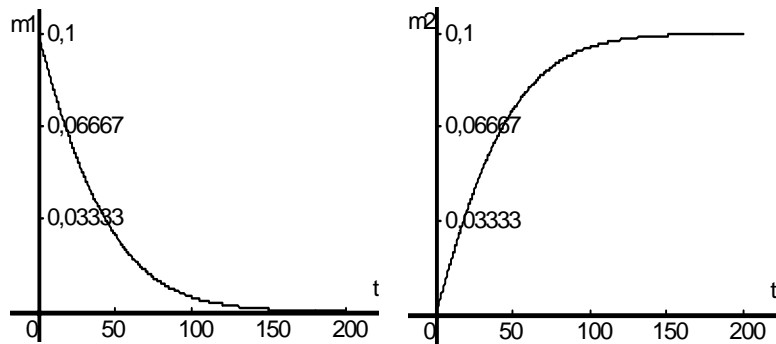
## MICHAELIS-MENTEN REACTION [GENOMIC OBJECT NET]



$V_{max} = 0.005$  (maximal reaction rate)

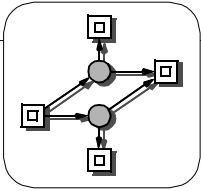
$K_m = 0.1375$  (Michaelis constant)

$$d[s]/dt = d[p]/dt = V_{max} * [s] / (K_m + [s])$$



## CASE STUDIES

- apoptosis in mammalian cells
- detailed glycolysis/pentose phosphate pathways in all human cells
- blood clotting in human (hemostasis versus fibrinolysis)
- lipoprotein metabolism (liver) in human
- G1/S - phase in mammalian cells
- detailed central carbon metabolism in potato tubers
- central carbon metabolism in *E.coli*



## REFERENCES I

### PETRI NETS, THEORY

**Baumgarten, B.:**

Petri Nets, Principles and Applications (in German);  
Wissenschaftsverlag 1990.

**Lautenbach, K.:**

Exact Liveness Conditions of a Petri Net Class (in German);  
Berichte der GMD 82, Bonn 1973.

**Murata, T.:**

Petri Nets: Properties, Analysis and Applications;  
Proc. of the IEEE 77(1989)4, pages 541-580.

**Petri, C. A.:**

Communication with Automata (in German);  
PhD Thesis, Schriften des Rheinisch-Westfälischen Instituts für Instrumentelle  
Mathematik an der Univ. Bonn, No. 2, 1962.

**Petri, C. A.:**

Interpretations of Net Theory;  
GMD, Interner Bericht 75-07, 2nd improved edition December 1976.

**Reisig, W.:**

Petri Nets - An Introduction;  
Springer 1982.

**Starke, P. H.:**

Analysis of Petri Net Models (in German);  
B. G. Teubner 1990.

**for more related books and papers see**

<http://www.daimi.au.dk/PetriNets/> -> Petri Net Bibliography

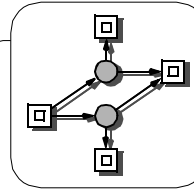
### PETRI NETS, TOOLS

**Fieber, M.:**

Design and Implementation of a Generic, Adaptive Tool to Handle Graphs (in  
German);  
Master Thesis, BTU Cottbus, CS Dep., Juli 2004.

**Starke, P. H.:**

INA - Integrated Net Analyzer, Manual;  
<http://www.informatik.hu-berlin.de/lehrestuehle/automaten/ina/#manual>  
Humboldt University Berlin, CS Dep., 1998.



## REFERENCES II

### PETRI NETS, BIO APPLICATIONS

**Heiner, M.; Koch, I.:**

Petri Net Based System validation in Systems Biology;  
Proc. ICATPN 2004, LNCS 3099, Springer 2004, pp. 216 - 237.

**Heiner, M.; Koch, I.; Will, J.:**

Model Validation of Biological Pathways Using Petri Nets - Demonstrated for  
Apoptosis;  
Journal BioSystems 2004, Vol 75/1-3, pp. 15-28.

**Koch, I.; Junker, B. H.; Heiner, M.:**

Application Petri Net Theory for Model Validation of the Sucrose-to-starch Pathway in  
Potato Tuber; submitted.

**Popova, L.; Heiner, M.; Koch, I.:**

Modelling and Analysis of Biochemical Networks with Time Petri Nets;  
Proc. CSP 2004, Caputh 2004.

**Voss, K.; Heiner, M.; Koch, I.:**

Steady State Analysis of Metabolic Pathways Using Petri Nets;  
Silico Biol. 3, 0031 (2003), <http://www.bioinfo.de/isb/2003/03/0031/>

**Reddy, V. N.; Mavrovouniotis, M. L.; Liebman, M. N.:**

Petri Net Representation in Metabolic Pathways.  
Proc. First Int. Conf. on Intelligent Systems for Molecular Biology, AAAI Press, Menlo  
Park, 1993, pp. 328-336.

**Reddy, V. N.; Liebman, M. N.; Mavrovouniotis, M. L.:**

Qualitative Analysis of Biochemical Reaction Systems;  
Comput. Biol. Med. 26(1996)1, pp. 9-24.

**for more related application papers see****Will, J.; Heiner, M.:**

Petri Nets in Biology, Chemistry, and Medicine - Bibliography;  
Computer Science Reports 04/02, BTU Cottbus, November 2002, 36 p.

### ELEMENTARY MODES / EXTREME PATHWAYS

**Schuster, S.; Hilgetag, C.; Schuster, R.:**

Determining Elementary Modes of Functioning in Biochemical Reaction Networks at  
Steady State. Proc. Second Gauss Symposium (1993) pp. 101-114

**Schilling, C. H.; Letscher, D.; Palsson, B. O.:**

Theory for the Systemic Definition of Metabolic Pathways and their Use in Interpreting  
Metabolic Function from a Pathway-Oriented Perspective;  
J. theor. Biol. (2000) 203, pp. 229-248.